Enterprise Multilayer and Routed Access Campus Design

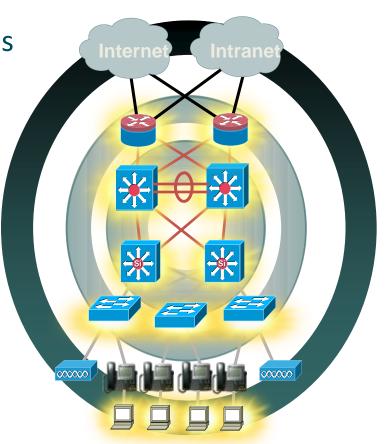


Yaman Hakmi Systems Engineer





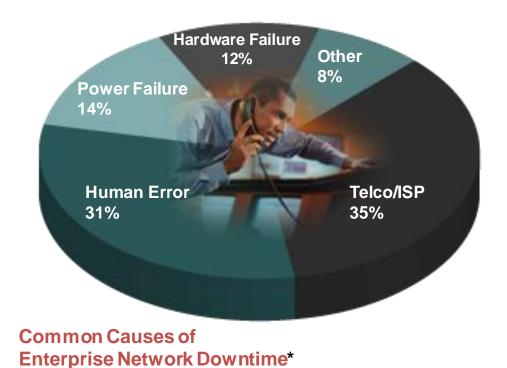
- Multilayer Campus Design Principles
- Latest Cisco Campus Networking Portfolio
 - Catalyst 6500
 - Nexus 7000
- Routed Access Campus Design
- Summary





Mitigating the Exposure

Most Common Causes of Downtime



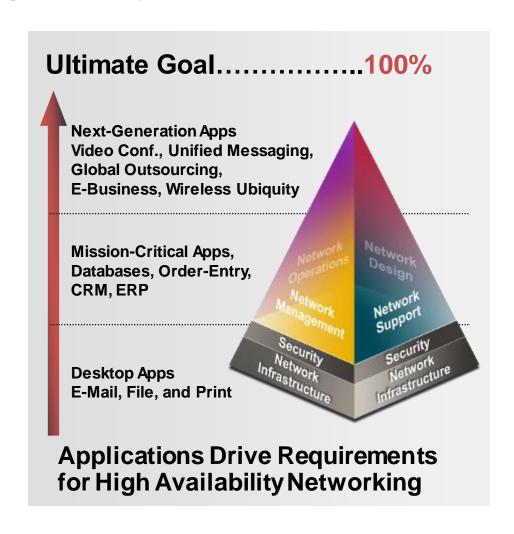
*Source: Yankee Group, The Road to Five-Nines Network

Enterprise Class Availability

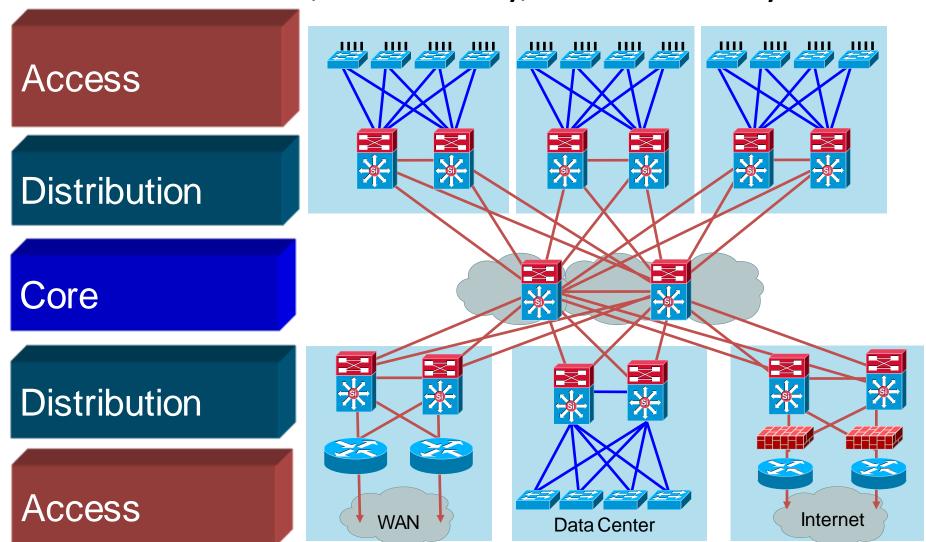
Resilient Campus Communication Fabric

Systems Design Approach to High Availability

- VOIP availability is the baseline for the enterprise networks
 - Human ear notices the difference in voice within 150–200 msec, which translates only ten consecutive packet loss with G711 codec
- Video loss is even more noticeable and it is rapidly becoming new frontier for jitter and delay requirements
- 200 msec end-to-end campus convergence is the design goal

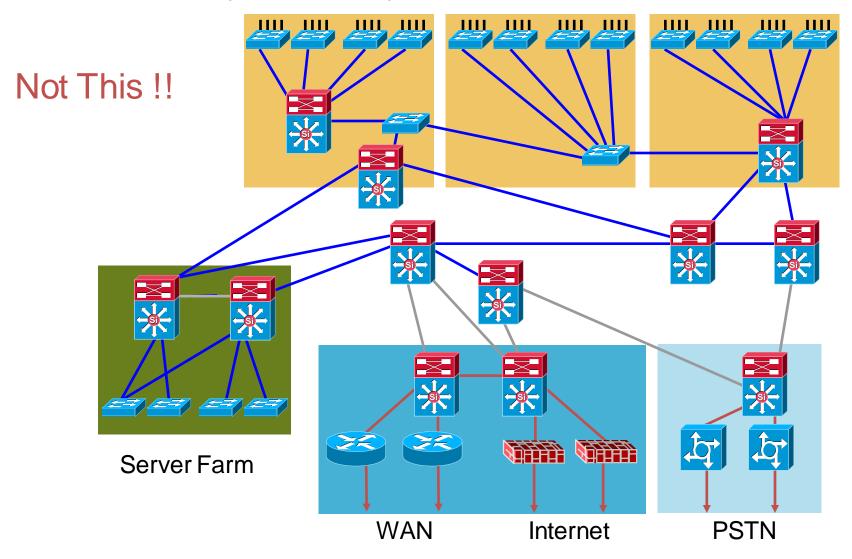


High Availability Campus Design Structure, Modularity, and Hierarchy



Hierarchical Campus Network

Structure, Modularity and Hierarchy



Hierarchical Network Design

Without a Rock Solid Foundation, the Rest Doesn't Matter

Access

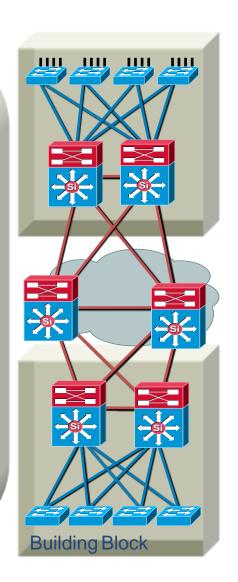
Distribution

Core

Distribution

Access

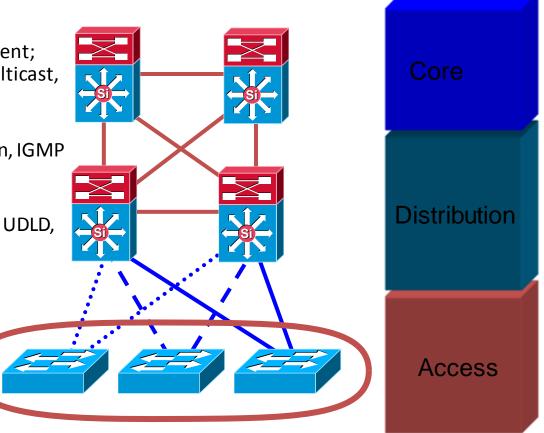
- Offers hierarchy—each layer has specific role
- Modular topology—building blocks
- Easy to grow, understand, and troubleshoot
- Creates small fault domains—
 Clear demarcations and isolation
- Promotes load balancing and redundancy
- Promotes deterministic traffic patterns
- Incorporates balance of both Layer 2 and Layer 3 technology, leveraging the strength of both
- Utilizes Layer 3 Routing for load balancing, fast convergence, scalability, and control



Access Layer

Feature Rich Environment

- It's not just about connectivity
- Layer 2/Layer 3 feature rich environment; convergence, HA, security, QoS, IP multicast, etc.
- Intelligent network services: QoS, trust boundary, broadcast suppression, IGMP snooping
- Intelligent network services: PVST+, Rapid PVST+, EIGRP, OSPF, PAgP/LACP, UDLD, FlexLink, etc
- Cisco Catalyst integrated security features (802.1x, CISF): port security, DHCP snooping, DAI, IPSG, etc.
- Automatic phone discovery, conditional trust boundary, power over Ethernet, auxiliary VLAN, etc.
- Spanning tree toolkit: Portfast, UplinkFast, BackboneFast, LoopGuard, BPDUGuard, RootGuard, etc.



Distribution Layer

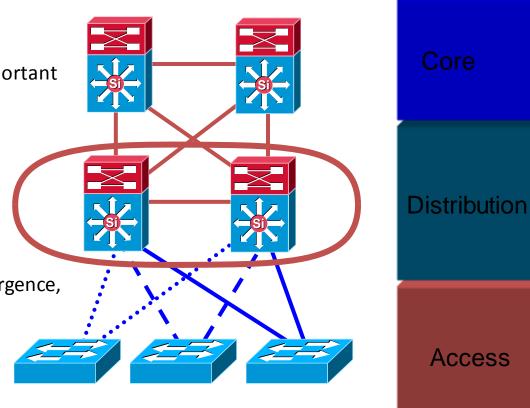
Policy, Convergence, QoS, and High Availability

Availability, load balancing,
 QoS and provisioning are the important
 considerations at
 this layer

- Aggregates wiring closets (access layer) and uplinks to core
- Protects core from high density peering and problems in access layer
- Route summarization, fast convergence, redundant

path load sharing

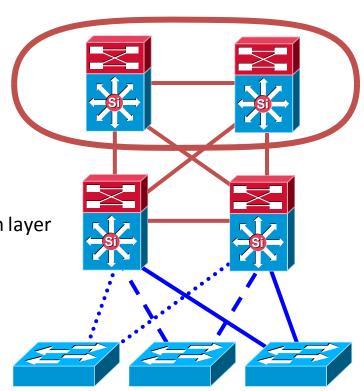
 HSRP or GLBP to provide first hop redundancy

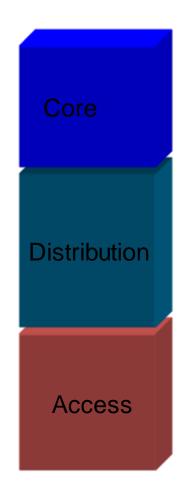


Core Layer

Scalability, High Availability, and Fast Convergence

- Backbone for the network - connects network building blocks
- Performance and stability vs. complexity - less is more in the core
- Aggregation point for distribution layer
- Separate core layer helps in scalability during future growth
- Keep the design technology-independent

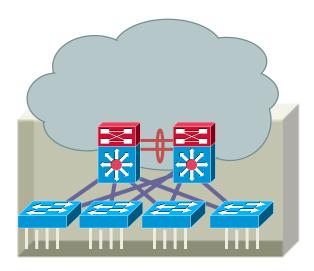




It's really a question of: Scale, Complexity, and Convergence

No Core

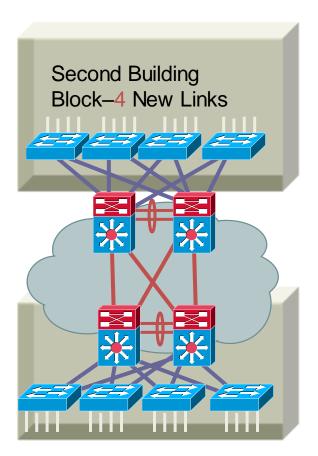
- Fully meshed distribution layers
- Physical cabling requirement
- Routing complexity



It's Really a Question of Scale, Complexity, and Convergence

No Core

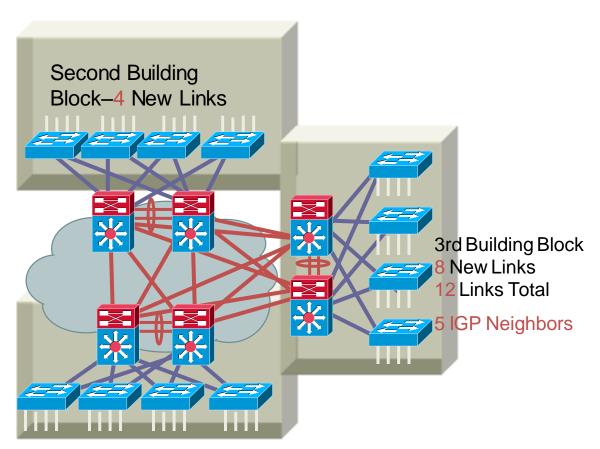
- Fully meshed distribution layers
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It's Really a Question of Scale, Complexity, and Convergence

No Core

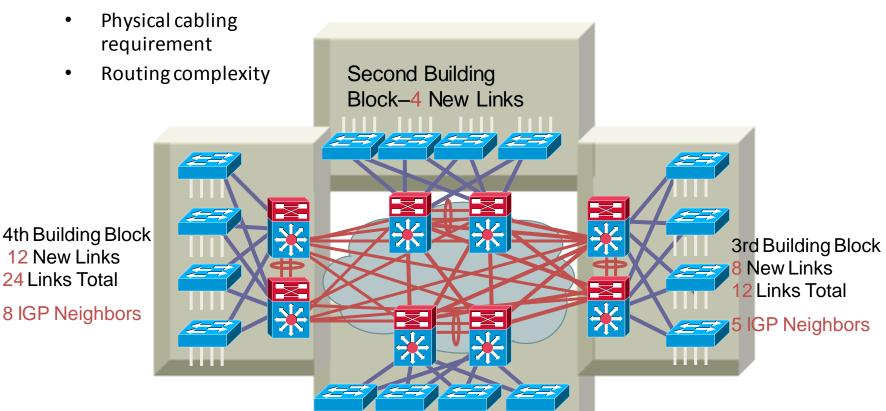
- Fully meshed distribution layers
- Physical cabling requirement
- Routing complexity



It's Really a Question of Scale, Complexity, and Convergence

No Core

Fully meshed distribution layers



It's Really a Question of Scale, Complexity, and Convergence

Dedicated Core Switches

Easier to add a module

Fewer links in the core

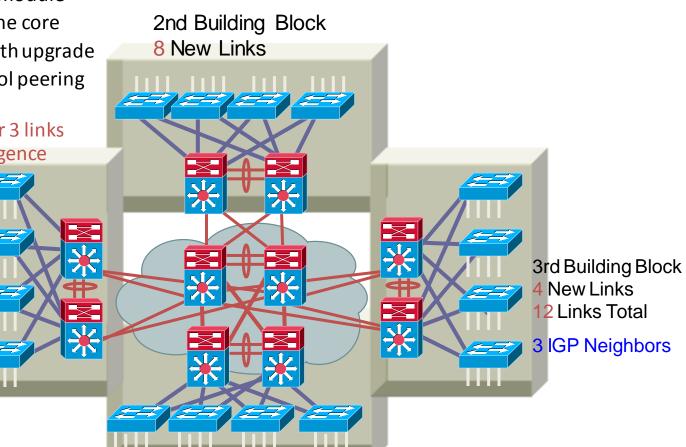
Easier bandwidth upgrade

 Routing protocol peering reduced

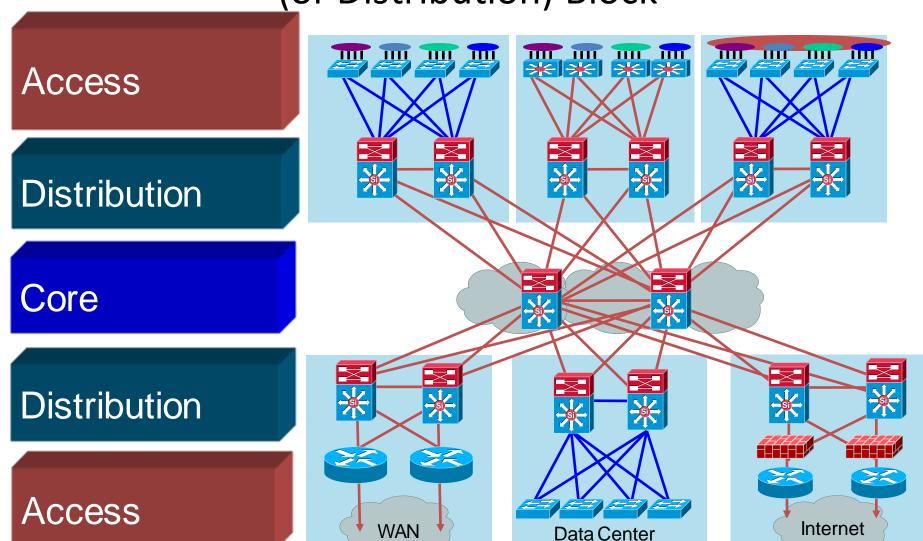
 Equal cost Layer 3 links for best convergence

4th Building Block 4 New Links 16 Links Total

3 IGP Neighbors



Design Alternatives come within a Building (or Distribution) Block



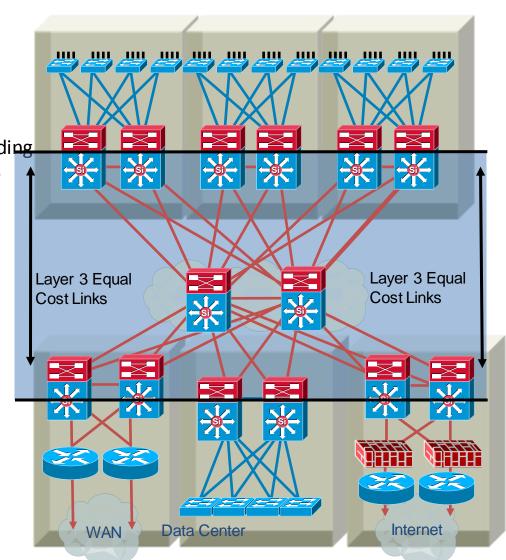
Core Layer: L3 Routing

 Typically deployed in distribution to core, and core to core interconnections

 Used to quickly re-route around failed node/links while providing load balancing over redundant paths

 Build triangles not squares for deterministic convergence

- Only peer on links that you intend to use as transit
- Insure redundant L3 paths to avoid black holes
- Summarize distribution to core to limit EIGRP query diameter or OSPF LSA propagation
- Tune CEF L3/L4 load balancing hash to achieve maximum utilization of equal cost paths (CEF polarization)



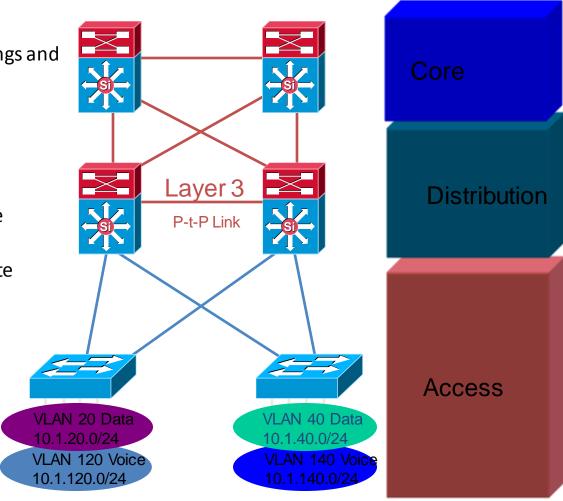
Layer 3 Distribution Interconnection

Reference Design—No VLANs Span Access Layer

- Tune CEF load balancing
- Match IOS EtherChannel settings and tune load balancing
- Summarize routes towards core
- Limit redundant IGP peering
- STP Root and HSRP primary tuning or GLBP to load balance on uplinks
- Set trunk mode on/nonegotiate
- Disable EtherChannel unless needed
- Set Port Host on access layer ports:

Disable Trunking
Disable EtherChannel
Enable PortFast

- RootGuard or BPDU-Guard
- Use security features



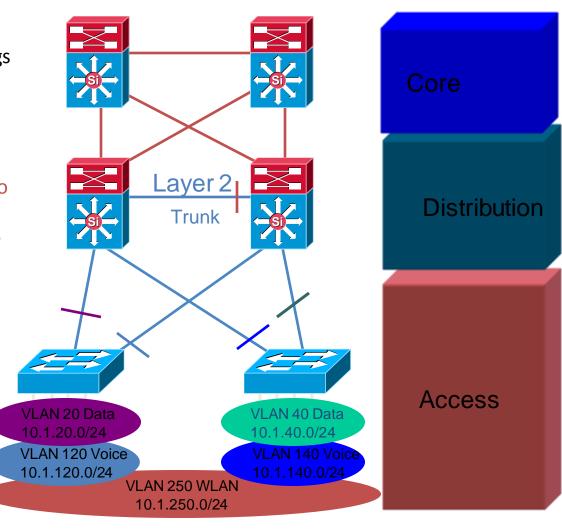
Layer 2 Distribution Interconnection

Some VLANs Span Access Layer

- Tune CEF load balancing
- Match IOS EtherChannel settings and tune load balancing
- Summarize routes towards core
- Limit redundant IGP peering
- STP Root and HSRP primary or GLBP and STP port cost tuning to load balance on uplinks
- Set trunk mode on/nonegotiate
- Disable EtherChannel unless needed
- RootGuard on downlinks
- LoopGuard on uplinks
- Set port host on access Layer ports:

Disable trunking
Disable EtherChannel
Enable PortFast

- RootGuard or BPDU-Guard
- Use security features



Multilayer Network Design Many Moving Parts

Evolved due to historical pressures

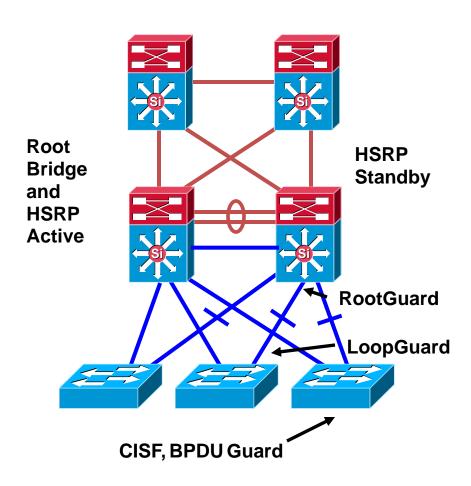
Speed of routing vs. switching Flexibility with spanning VLANs Nonroutable protocols

 Well understood optimization of interaction between the various control protocols and the topology

STP root and HSRP primary tuning to load balance on uplinks

Spanning tree toolkit (RootGuard, LoopGuard...)

 Many moving parts increase a chance of instability



Multilayer Network Design

Good Solid Design Option, But...

Utilizes multiple Control Protocols

Spanning Tree (802.1w, ...), FHRP (HSRP, ...), Routing Protocol (EIGRP, ...)

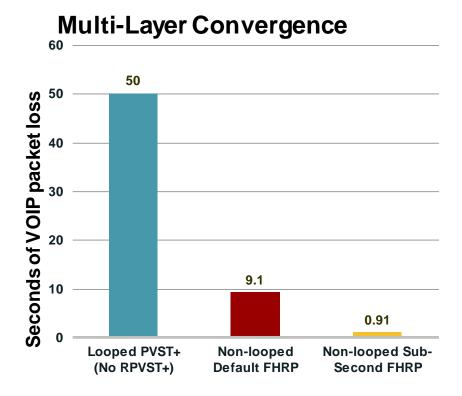
 Convergence is dependent on multiple factors

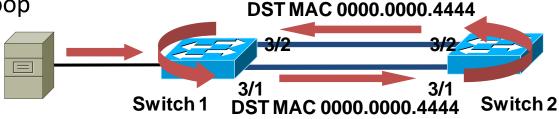
FHRP—900msec to 9 seconds

Spanning Tree—Upto 50 seconds

Poor load balancing—single uplink, asymmetric routing etc

STP, if it breaks badly, no inherent mechanism to stop
 the loop





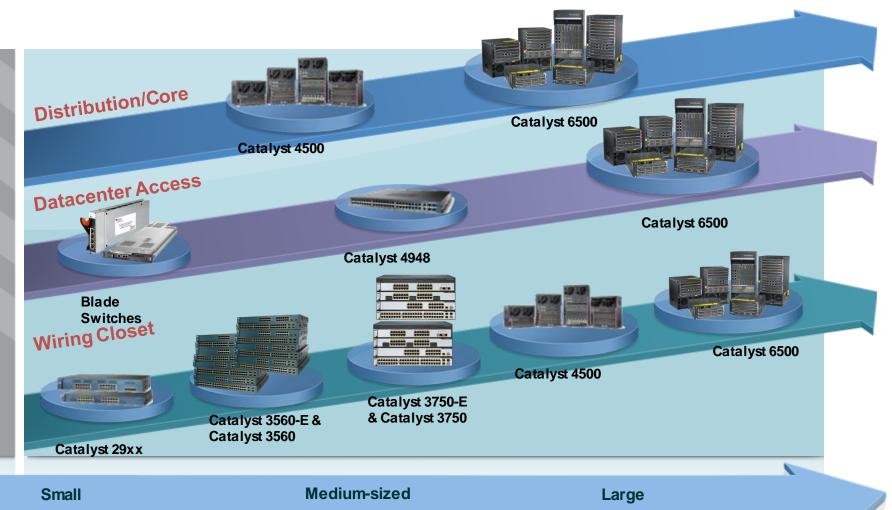
Latest Campus Networking Technologies



Cisco Expo 2009



Catalyst Switching Portfolio



Catalyst 6500 Linecards

Latest Core, Distribution, and Data Center Portfolio



WS-X6708-10GE



WS-X6704-10GE



WS-X6748-SFP



WS-X6724-SFP



WS-X6748-GE-TX

8 Port 10GE X2 (2x20G Fabric Enabled Card)

ER, LR, LX4, SR, CX4 optics; DFC included; Queuing: TX-1p7q4t, RX-8q8t; Jumbo frame support: 200MB/port Buffer

4 Port 10GE XENPAK (2x20G Fabric Enabled Card)

ER, LR, LX4, SR, CX4, ZR optics; Optional DFC; Queuing: TX-1p7q4t, RX-1q8t or 8q8t (w/DFC); Jumbo frame support

48 Port GE SFP (2x20G Fabric Enabled Card)

SX, LX, ZX, Tx, CWDM SFPs; Optional DFC; Queuing: TX-1p3q8t, RX-1q8t or 2q8t (w/DFC); Jumbo frame support

24 Port GE SFP (1x20G Fabric Enabled Card)

SX, LX, ZX, Tx, CWDM SFPs; Optional DFC; Queuing: TX-1p3q8t, RX-1q8t or 2q8t (w/DFC); Jumbo frame support

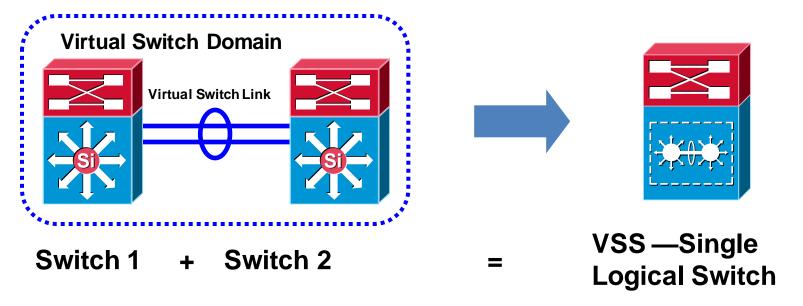
48 Port 10/100/1000 (2x20G Fabric Enabled Card)

Supports TDR; Optional Distributed Forwarding Card (DFC); Queuing: TX - 1p3q8t, RX - 1q8t; Jumbo frame support

Virtual Switch

Virtual Switching System 1440 (VSS)

- Virtual Switching System consists of two Cisco Catalyst 6500 Series defined as members of the same virtual switch domain
- Single control plane with dual active forwarding planes
- Design to increase forwarding capacity while increasing availability by eliminating STP loops
- Reduced operational complexity by simplifying configuration



Virtual Switching System

Hardware and Software Requirements

Software Support

Native and modular Cisco IOS are supported

Minimum IOS required is 12.2(33)SXH1, however current recommendation is 12.2(33)SXH2(a)

Supervisor—VS-S720-10G-3C/XL

PFC3C/XL contains new hardware support to forward traffic across multiple physical chassis and lookup enhancements

Virtual switch link

VS header encapsulation requires new port ASIC

VS-S720-10G-3C/XL Supervisor 10G port or WS-X6708-10G-3C/XL

10 Gigabit Ethernet only

• WS-X6716: 16-port 10Gbps line card

Catalyst 6500 16p 10GBASE-T line card





VS-S720-10G-3C/XL

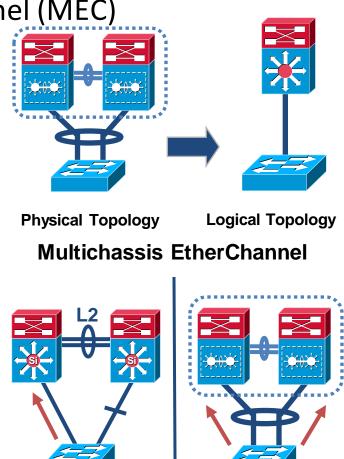


WS-X6708-10G-3C/XL

Virtual Switching System

Multichassis EtherChannel (MEC)

- MEC is an advanced EtherChannel technology extending link aggregation to two separate physical switches
- MEC enables the VSS appear as single logical device to devices connected to VSS, thus significantly simplifying campus topology
- Traditionally spanning VLANs over multiple closets would create STP looped topology, MEC with VSS eliminates these loops in the campus topology
- MEC replaces spanning tree as the means to provide link redundancy and thus doubling bandwidth available from access
- MEC is supported only with VSS



Non-MEC MEC
BW Capacity in Non-MEC and MEC Topology

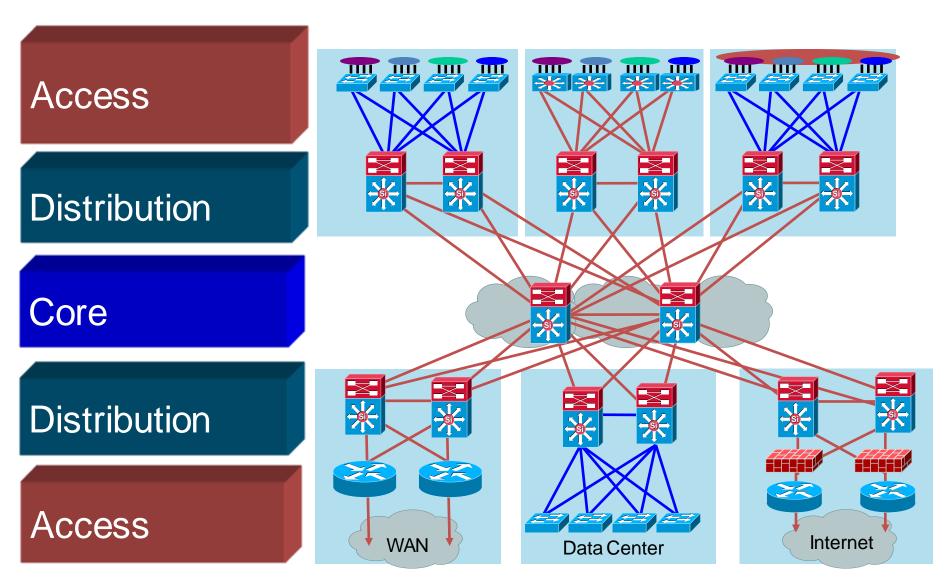
Cisco Nexus 7000

- Nexus Family: Increasingly seen in the enterprise campus network
- High density
 - 256 10G interfaces per system
 - 384 1G interfaces per systems
- High performance
 - 64 non-blocking 10G ports
 - 1.2Tbps system bandwidth at initial release
 - 80Gbps per slot
- Future proof
 - Initial fabric provides up to 4.1Tbps
 - Product family scaleable to 15+Tbps
- Nexus 7018: Second chassis in Nexus 7000 family
- Ultra-high density
 - 512 10G interfaces per system
 - 768 1G interfaces per system
- High performance
 - 128 non-blocking 10G ports
 - 2.5Tbps system bandwidth at initial release; 80G/slot
- Future proof
 - Initial fabric provides up to 7.8Tbps
 - Chassis scaleable to 17.6Tbps

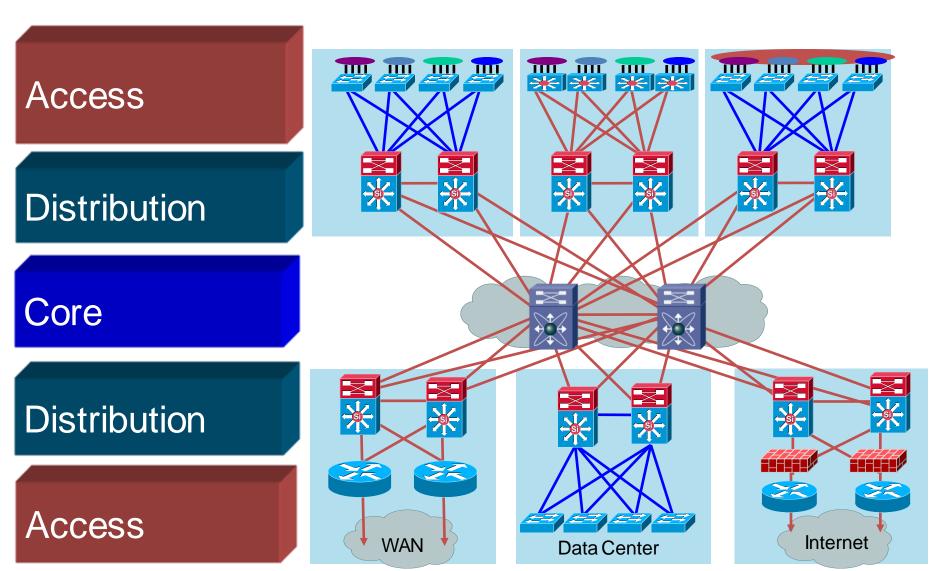




Recall: Campus Network Design



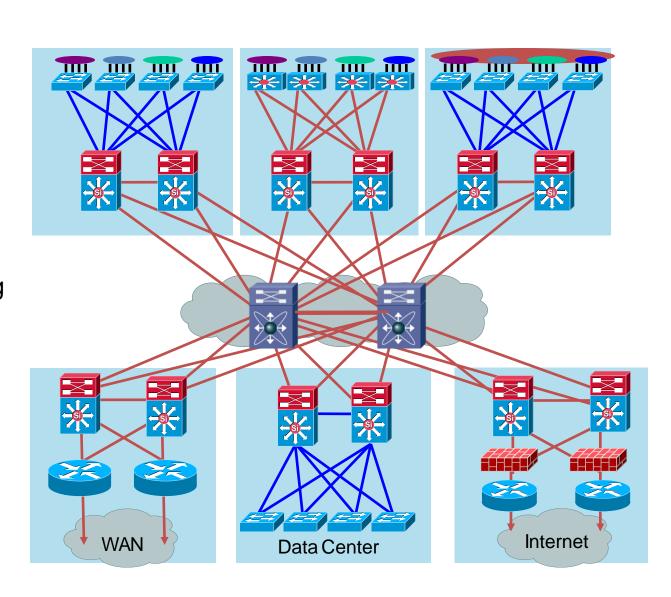
Campus Network Design with the Nexus 7000



Campus Network Design with the Nexus 7000

Drive for Nexus 7000 in the enterprise core:

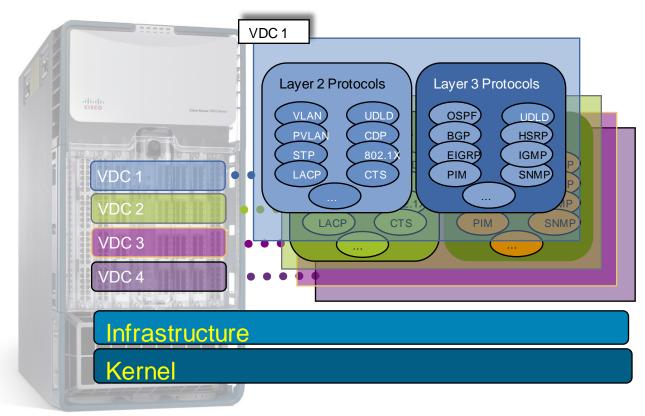
- Higher 10 Gbps port density
- Increased backbone switching and routing capacity
- Enhanced network resiliency with SSO/ISSU



Virtual Device Contexts

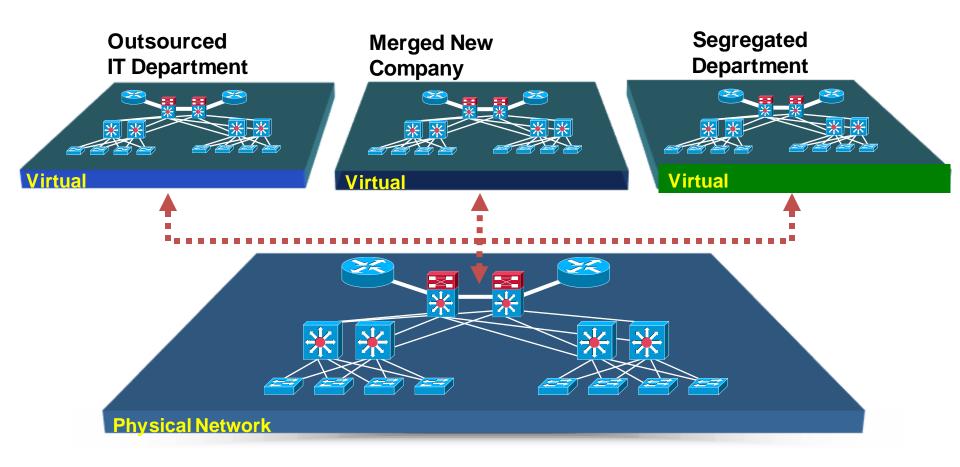
Virtual Device Contexts provides virtualization at the device level allowing multiple instances of the device to operate on the same physical switch at the same

time.



Network Segmentation

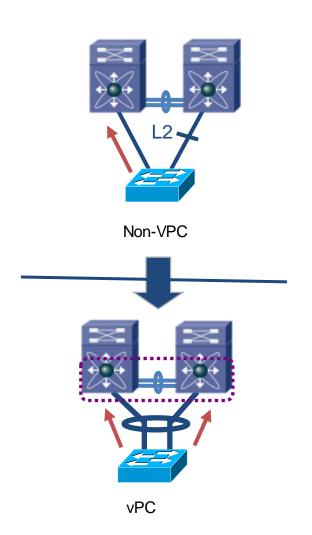
1 to Many: One network supports many virtual networks



Virtual Port-Channel (vPC)

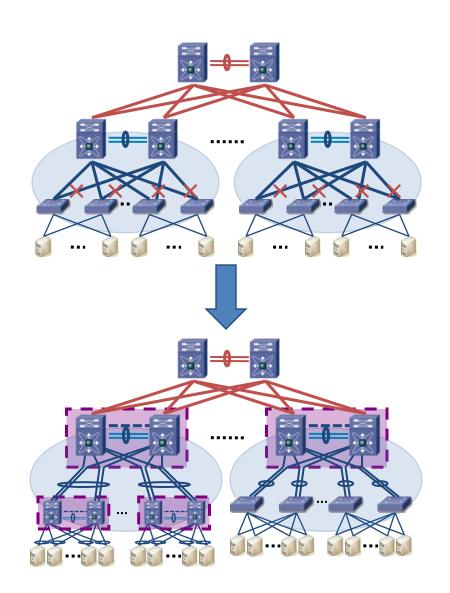
Feature Overview

- Allow a single device to use a port channel across two upstream switches
- Eliminate STP blocked ports
- Uses all available uplink bandwidth
- Dual-homed server operate in active-active mode
- Provide fast convergence upon link/device failure
- Reduce CAPEX and OPEX
- Available in NX-OS 4.1 with current and future hardware (10G card only)



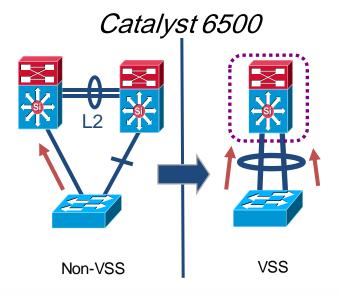
How does vPC help with STP?

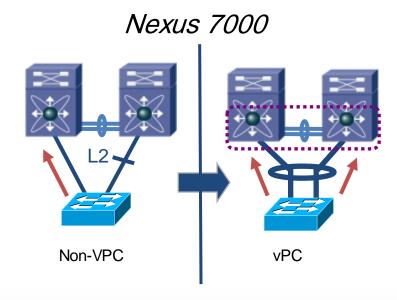
- Before vPC
 - STP blocks redundant uplinks
 - VLAN based load balancing
 - Re-convergence relies on STP
 - Protocol Failure →
- With vPC
 - No blocked uplinks
 - Lower oversubscription
 - EtherChannel load balancing (hash)
 - Convergence sub-second
 - Reduced STP logical port count



Virtual Port Channels

Multi-Chassis Etherchannel (MEC)





Virtual Switching System (VSS)

Virtual Port Channel (vPC)

- Both VSS-MEC and vPC are a Port-channeling concept extending link aggregation to two separate physical switches
- Allows the creation of resilient L2 topologies based on Link Aggregation.
- Eliminates the dependence on STP in the L2 access-distribution Layer
- Enable seamless VM Mobility, Server HA Clusters
- Scale Available Layer 2 Bandwidth
- Simplify Network Design

Routed Access Campus Design

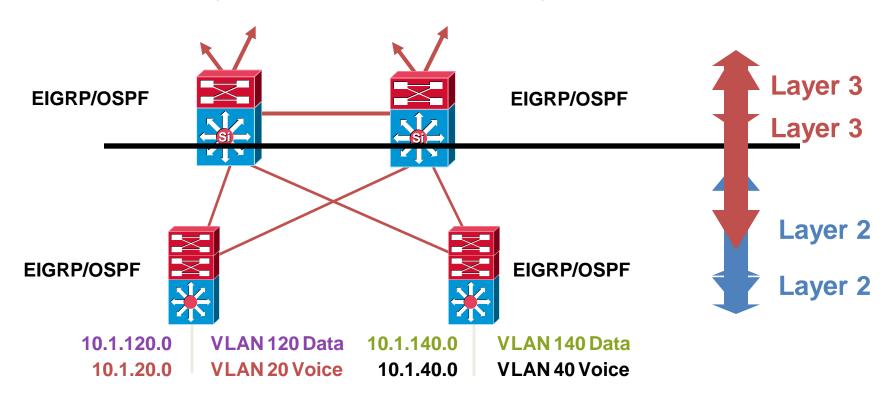


Cisco Expo 2009



Routed Access

Layer 3 Distribution with Layer 3 Access



- Move the Layer 2/3 demarcation to the network edge
- Upstream convergence times triggered by hardware detection of light lost from upstream neighbor
- Beneficial for the right environment

Routed Access

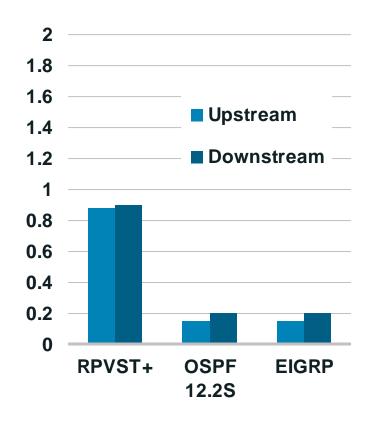
Advantages, Yes in the Right Environment

- EIGRP converges in < 200 msec
- OSPF with subsecond tuning converges in < 200 msec
- Ease of implementation, less to get right

No matching of STP/HSRP/ GLBP priority No L2/L3 multicast topology inconsistencies

- Single control plane and well known tool set traceroute, show ip route, show ip eigrp neighbor, etc.
- Convergence times dependent on GLBP/HSRP tuning

Both L2 and L3 Can Provide Subsecond Convergence



Routed Access Simplified Network Recovery

- Routed access network recovery is dependent on L3 reroute
- Time to restore upstream traffic flows is based on ECMP (Equal Cost Multi-Path) reroute

Time to detect link failure

Process the removal of the lost routes from the software CEF (Cisco® Express Forwarding)

Update the hardware CEF table

 Time to restore downstream flows is based on a full routing protocol reroute

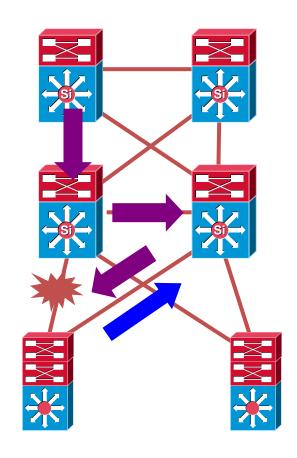
Time to detect link failure

Time to determine new route

Process the update of the software RIB and FIB

Update the hardware FIB

 Except uplink failure, all other fault recovery is ECMP-based i.e., consistent and predictable

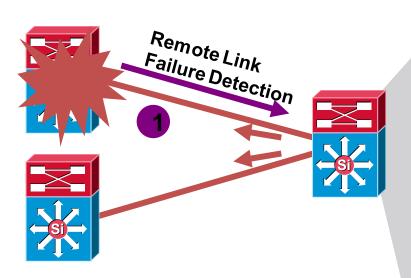


Upstream: ECMP Recovery

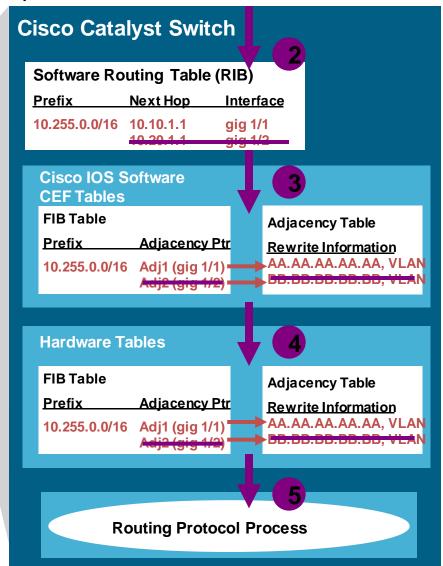
Downstream: Routing Protocol Recovery

Routed Access

Time to Recovery CEF Paths



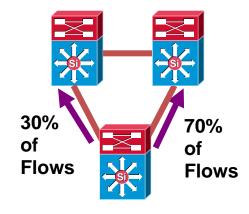
- Link Failure Detection
- Removal of the Entries in the Routing Table
- Update of The Software CEF
 Table to Reflect to Loss of the
 Next Hop Adjacencies
- Update of the Hardware Tables
- Routing Protocol Notification and Reconvergence



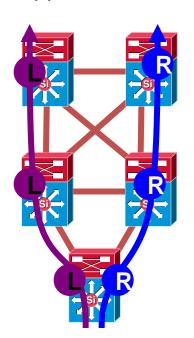
Equal Cost Multipath

Optimizing CEF Load-Sharing

- CEF load-sharing options influence the traffic utilization of Equal Cost Multi-Path (ECMP) links
- Criteria for optimizing CEF load-sharing options would depend on IP address plan and application flows (ports, location, etc.)







different possible side effects when using ECMP
 CEF polarization can only happen with multiple stages

CEF polarization and unequal load-sharing are two

• If each stage uses the same decision criteria some links can be totally unused

of CEF load balancing decisions

CEF Polarization

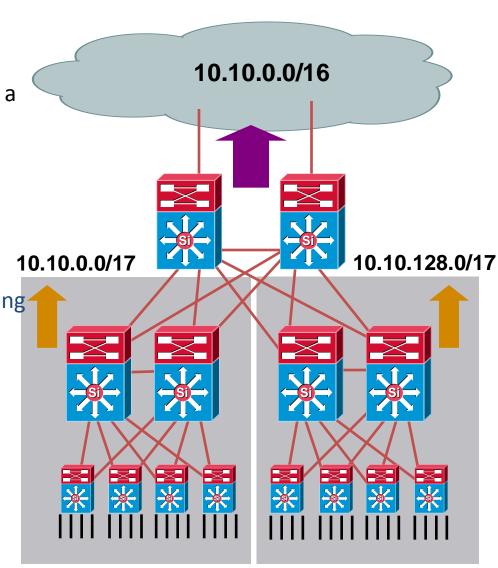
EIGRP Design Rules for Routed Access

Leverage the Tools Provided

 The greatest advantages of EIGRP are gained when the network has a structured addressing plan that allows for use of summarization and stub routers

EIGRP provides the ability
 to implement multiple tiers
 of summarization and route filtering

- Relatively painless to migrate to a L3 access with EIGRP if network addressing scheme permits
- Able to maintain a deterministic convergence time in very large L3 topology



EIGRP Design Rules for HA Campus

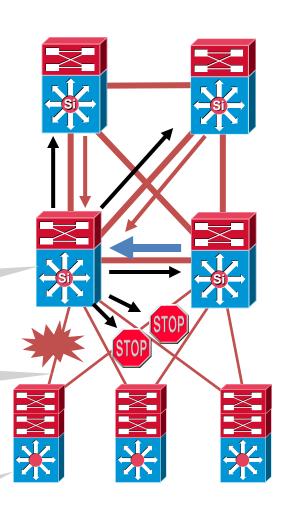
High-Speed Campus Convergence

- EIGRP convergence is largely dependent on query response times
- Minimize the number and time for query response to speed up convergence
- Summarize distribution block routes upstream to the core
- Configure all access switches as EIGRP stub routers
- Filter routes sent down to access switches

```
interface TenGigabitEthernet < Uplinks to the core >
  ip summary-address eigrp 100 10.120.0.0 255.255.0.0 5
<No summary routes on links between distributions>
```

```
router eigrp 100
network 10.0.0.0
distribute-list Default out <downstream links>
ip access-list standard Default
permit 0.0.0.0
```

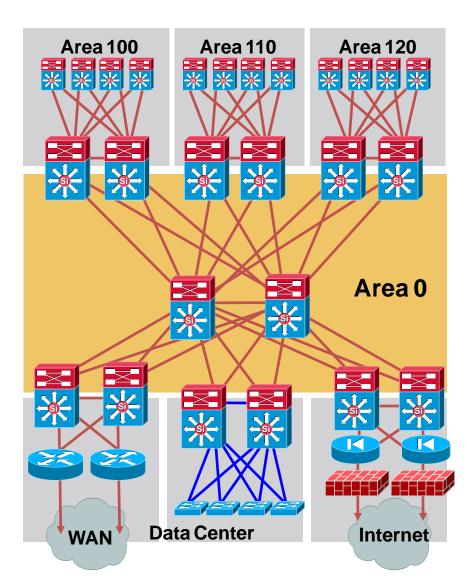
```
router eigrp 100
network 10.0.0.0
eigrp stub connected
```



OSPF Design Rules for Routed Access

Where Are the Areas?

- Area size/border is bounded by the same concerns in the campus as the WAN
- In campus the lower number of nodes and stability of local links could allow you to build larger areas
- Area design also based on address summarization
- Area boundaries should define buffers between fault domains
- Keep area 0 for core infrastructure and do not extend to the access routers



Utilize Totally Stubby Areas

Reduce SPF and LSA Load in Distribution Area

ABR for a regular area forwards

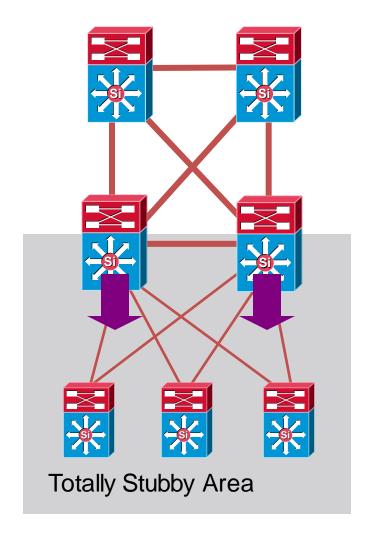
```
Summary LSAs (Type 3)
ASBR summary (Type 4)
Specific externals (Type 5)
```

Stub area ABR forwards

```
Summary LSAs (Type 3)
Summary default (0.0.0.0)
```

A totally stubby area ABR forwards
 Summary default (0.0.0.0)

```
router ospf 100
area 120 stub no-summary
area 120 range 10.120.0.0 255.255.0.0 cost 10
network 10.120.0.0 0.0.255.255 area 120
network 10.122.0.0 0.0.255.255 area 0
```

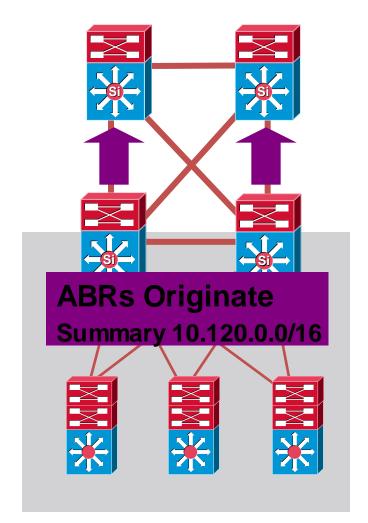


Summarization Distribution to Core

Reduce SPF and LSA Load in Area 0

- Summarize routes from the distribution block upstream into the core
- Minimize the number of LSAs and routes in the core
- Reduce the need for SPF calculations due to internal distribution block changes
- Incremental SPF (iSPF) is a mechanism to reduce the computational load of larger OSPF areas but is more applicable to WAN than campus environments

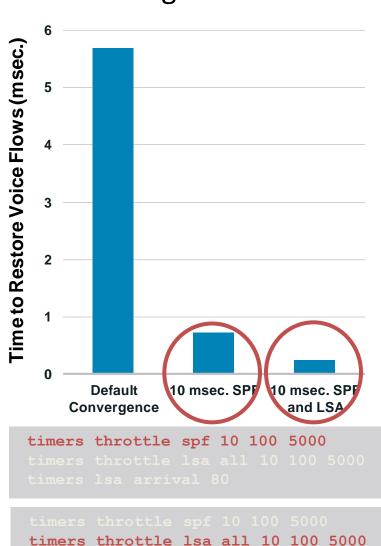
```
router ospf 100
area 120 stub no-summary
area 120 range 10.120.0.0 255.255.0.0 cost 10
network 10.120.0.0 0.0.255.255 area 120
network 10.122.0.0 0.0.255.255 area 0
```



OSPF Design Rules for HA Campus

OSPF SPF and LSA Throttling

- OSPF has an SPF throttling timer designed to dampen route recalculation
- 12.2S OSPF enhancements let us tune this timer to milliseconds; prior to 12.2S one second was the minimum
- After a failure, the router waits for the SPF timer to expire before recalculating a new route
- By default, there is a 500 ms delay before generating router and network LSAs; the wait is used to collect changes during a convergence event and minimize the number of LSAs sent
- Propagation of a new instance of the LSA is limited at the originator
- Acceptance of a new LSAs is limited by the receiver
- Make sure Isa-arrival < Isa-hold

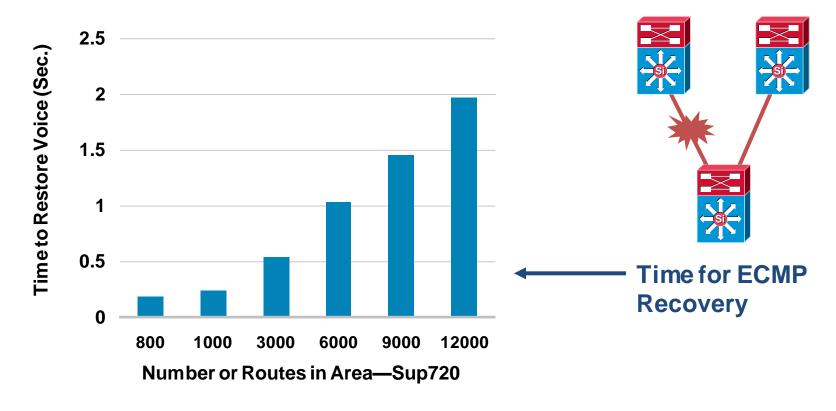


timers 1sa arrival 80

Convergence

ECMP Convergence Is Dependent on Number of Routes

- Time to update of switch HW FIB is linearly dependent on the number of entries (routes) to be updated
- Summarization will serve to decrease RP load as well as speed up convergence



Routed Access Design Considerations Design Requirements

- VLANs are localized to a single wiring closet switch
- IP addressing—do you have an address allocation plan to support a routed access design?
- Platform requirements

Requires a Cisco Catalyst® 3550 Series, Cisco Catalyst 3560 Series or above

Requires Cisco IOS® (native or hybrid)

Cisco Catalyst 6500 Series requires a Supervisor with an MSFC

Cisco Catalyst IOS feature set considerations

IP base feature set for EIGRP-Stub and PIM

IP services feature set for OSPF and PIM

Cisco Catalyst 3000 Series require 12.2(37)SE for PIM stub in IP base



Routed Access Design Considerations Design Motivations

Simplified Control Plane

No STP feature placement (root bridge, loopguard, ...)

No default gateway redundancy setup/tuning

No matching of STP/HSRP priority

No L2/L3 multicast topology inconsistencies

Ease of Troubleshooting (leverage well know toolset)

Show ip route

Traceroute

Ping and extended pings

Extensive protocol debugs

End to end consistent troubleshooting: Campus, WAN to Data Center

Failure differences

Routed topologies fails safe—i.e. loss of routing protocol neighbor disables path Layer 2 topologies fail open—i.e. loss of spanning tree BPDU's can open a blocked link and cause flooding

VSS Enabled
Access Layer
Campus Design



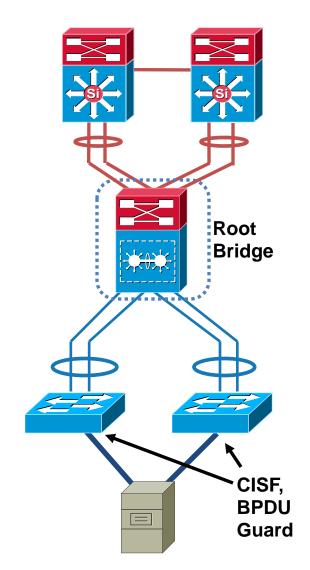
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VSS Enabled Campus Design

Control Plane Simplification

- VSS makes the network loop-free in normal topology; do not disable spanning tree to safeguard against possible loop introduced at the edge due to user error and daisy chaining
- Simplifies the topology allowing VLANs to span to increase flexibility in design options
- Ease of implementation, less to get it right
 No need for HSRP, GLBP, or VRRP
 No reliance on subsecond FHRP timers
 No asymmetric forwarding
- A single logical multicast router on the access subnets simplifies the multicast topology resulting in elimination of non-RPF traffic
- Redundant supervisors provide resiliency via SSO-enabled protocols; consistent recovery during the failover of nodes at the distribution



VSS Enabled Campus Design

Impact to the Campus Topology

Physical network topology does not change

Still have redundant chassis

Still have redundant links

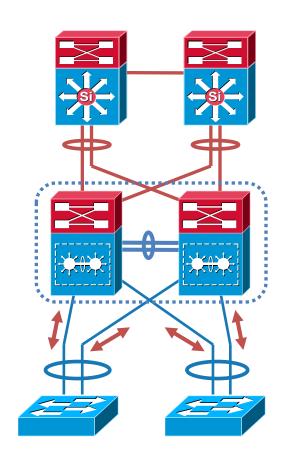
 Logical topology is simplified as we now have a single control plane

No unicast flooding Single configuration management

 Allows the design to replace traditional topology control plane with multichassis EtherChannel (MEC)

Enables loopfree topology, thus doubles the link capacity

Convergence and load balancing are based on EtherChannel



VSS Enabled Campus Design MEC Configuration

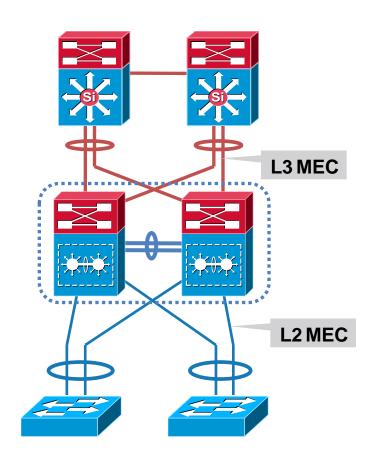
 MEC links on both switches are managed by PAgP or LACP running on the ACTIVE switch via internal control messages

All the rules and properties of EtherChannel applies to MEC such as negotiation, link characteristics (port-type, trunk), QoS, etc.

 Do not use "on" and "off" options with PAgP or LACP protocol negotiation

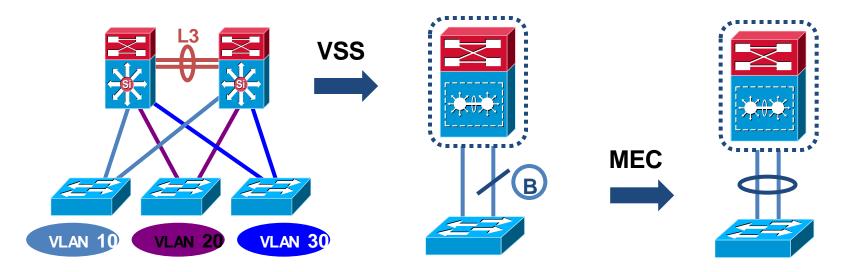
PAgP—Run Desirable-Desirable with MEC links LACP—Run Active-Active with MEC links

- L2 MEC enables loop free topology and doubles the uplink bandwidth as no links are blocked
- L3 MEC provides reduced neighbor counts, consistent load-sharing (L2 and L3) and reduced VSL link utilization for multicast flows



VSS Enabled Campus Design Multilayer Topology

- Optimized multilayer topology uses "V" shape design where VLANs do not span closets
- Deploying VSS in such topology without MEC reintroduces STP loops in the networks
- Use of MEC is recommended any time two L2 links from the same devices connected to VSS



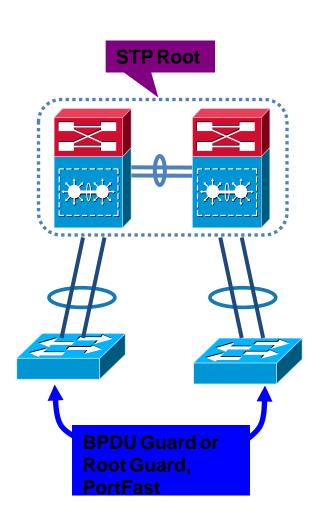
- Each access switch has unique VLANs
- No layer 2 loops
- No blocked links

Layer 2 Loop
Blocking One Link

MEC Creates Single Logical Link, No Loops, No Blocked Links

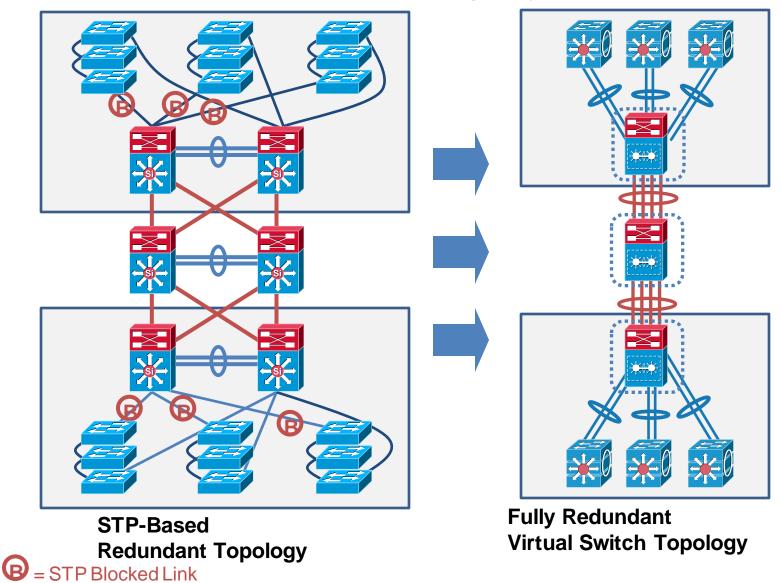
VSS Enabled Campus Design STP Optimization

- Make sure VSS remains root of all VLANs
- Do not use loop guard as it will disable the entire MEC channel on fault detection
- Use root guard at the edge port to protect external switch introducing superior BPDUs, e.g., temporary connectivity
- BPDU guard and root guard are mutually exclusive
- PortFast and BPDU guard is still necessary at the edge switch to prevent accidental loop introduce either due to user error or topology change



VSS Enabled Campus Design

End-to-End VSS Design Option

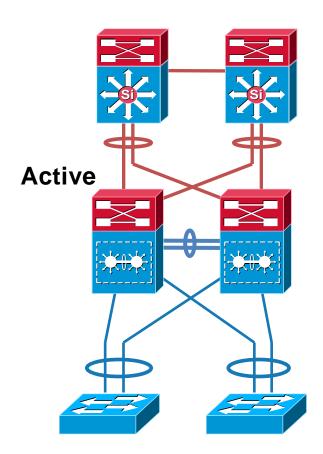


VSS Enabled Campus Design Summary

- VSS enables highly available campus with sub-second convergence without the complexity of managing dual node at distribution layer
- Eliminates FHRP configuration
- Must use L2 MEC to create loop free topology, STP should remained enabled
- Use of L3 MEC significantly improves convergence for multicast traffic
- Enabled NSF in adjacent routed devices for better convergence

Use default Hello and Hold timers for EIGRP & OSPF

 Use STP tool kits guidance applicable to loop free "V" shape design



Summary



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Next Generation Campus Design

Evolving the Campus Foundation Architecture

	Multitier Access	Routed Access	Virtual Switch
Access Distribution Control Plane Protocols	Spanning Tree (PVST+, Rapid- PVST+ or MST)	EIGRP or OSPF	PAgP, LACP
Spanning Tree Required	STP Required for Network Redundancy and to Prevent L2 Loops	No	No
Network Recovery Mechanisms	Spanning Tree and FHRP (HSRP, GLBP, VRRP)	EIGRP or OSPF	Multichassis EtherChannel (MEC)
VLAN Spanning Wiring Closets	Supported (Not Desirable Design)	No	Supported

Next Generation Campus Design

Evolving the Campus Foundation Architecture

	Multitier Access	Routed Access	Virtual Switch
Layer 2/3 Demarcation	Distribution	Access	Distribution (Could Be Access)
First Hop Redundancy Protocol	HSRP, GLBP, VRRP Required	Not Required	Not Required
Load Balancing	Per Subnet or Host	Per Flow—ECMP	Per Flow—MEC
Convergence	900 msec—50 Seconds (Dependent on STP Topology and FHRP Tuning	50–600 msec	50–600 msec

Next Generation Campus Design

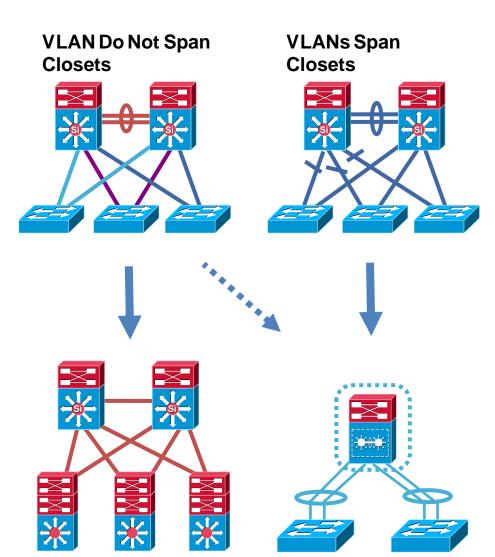
Evolving the Campus Foundation Architecture

- Traditional Layer 2 designs can evolve to VSS
- Evolving architectures provide:

Simplified control plane: remove dependence on STP

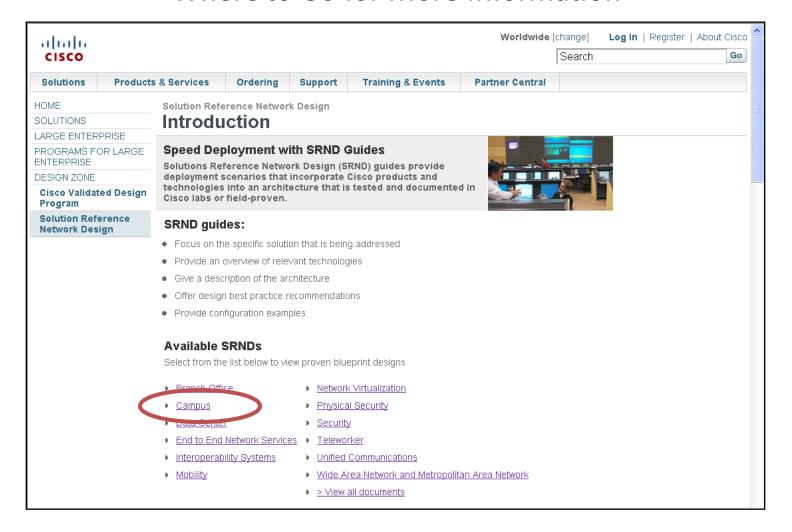
Increased capacity: provide flow-based load balancing High availability: 200 msec or better recovery

 Flexibility to provide for the right implementation for each network requirement



Campus Design Guidance

Where to Go for More Information



http://www.cisco.com/go/srnd/ and http://www.cisco.com/go/cvd

Q & A



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Yaman Hakmi



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