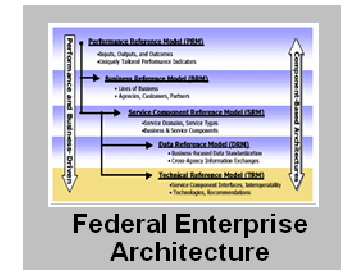
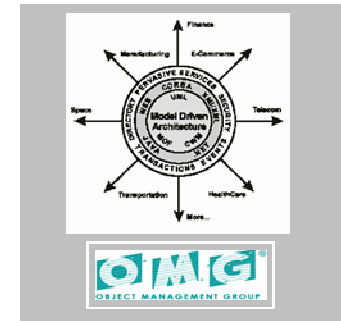
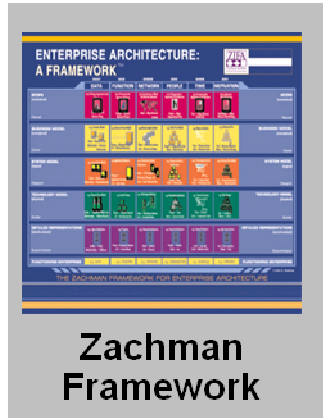
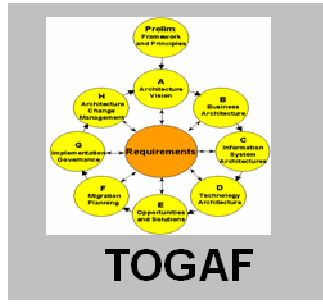




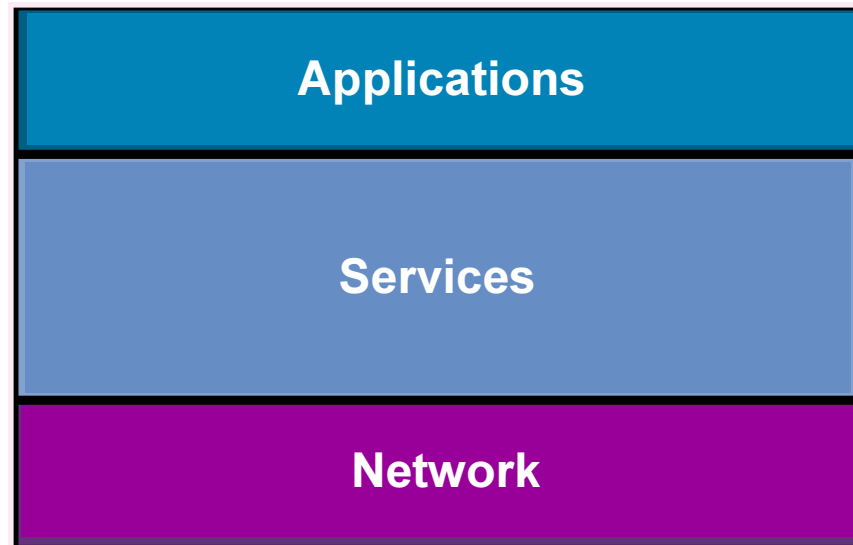
An Architectural Approach to Campus Networks



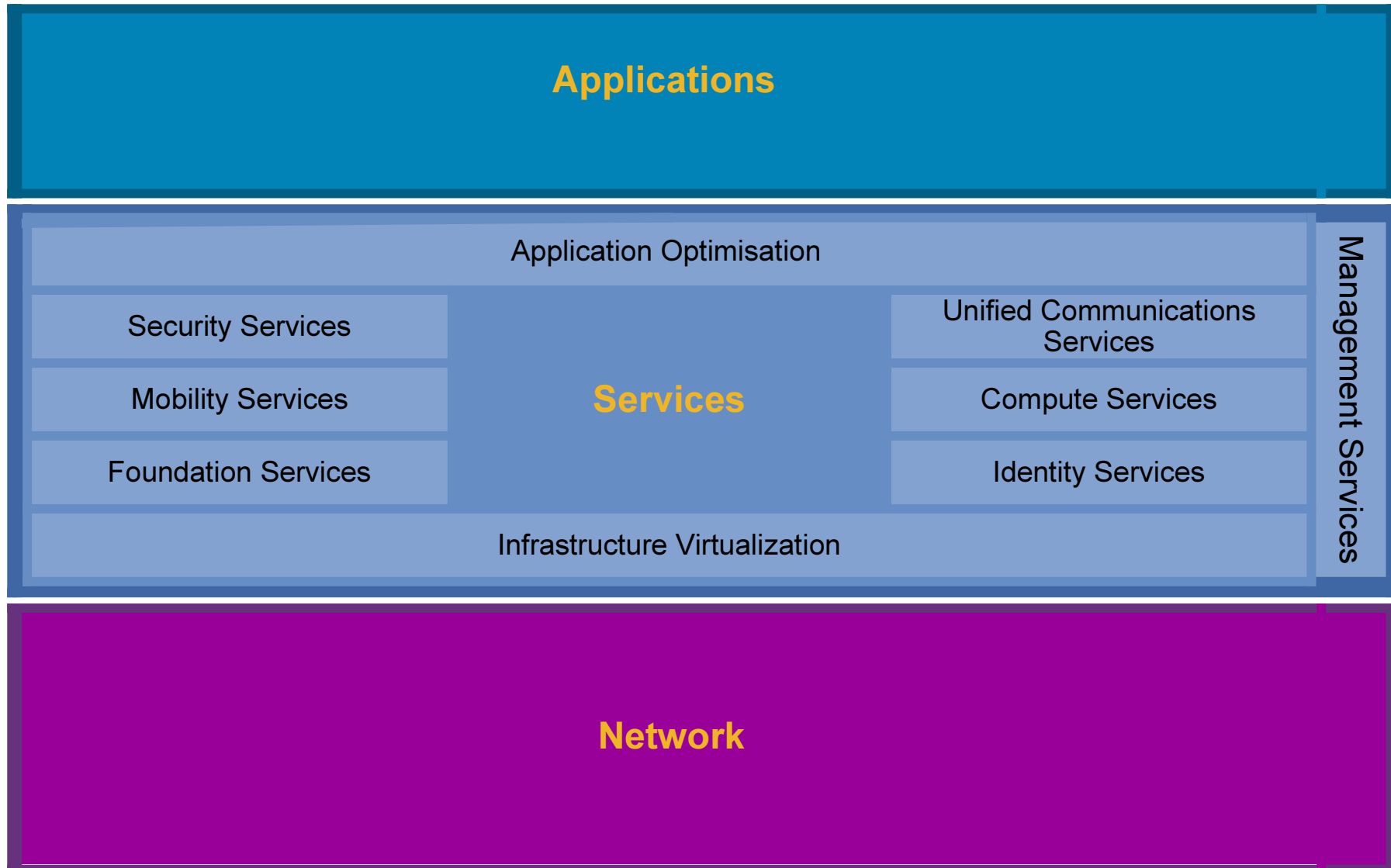
Enterprise Architectures & The Network



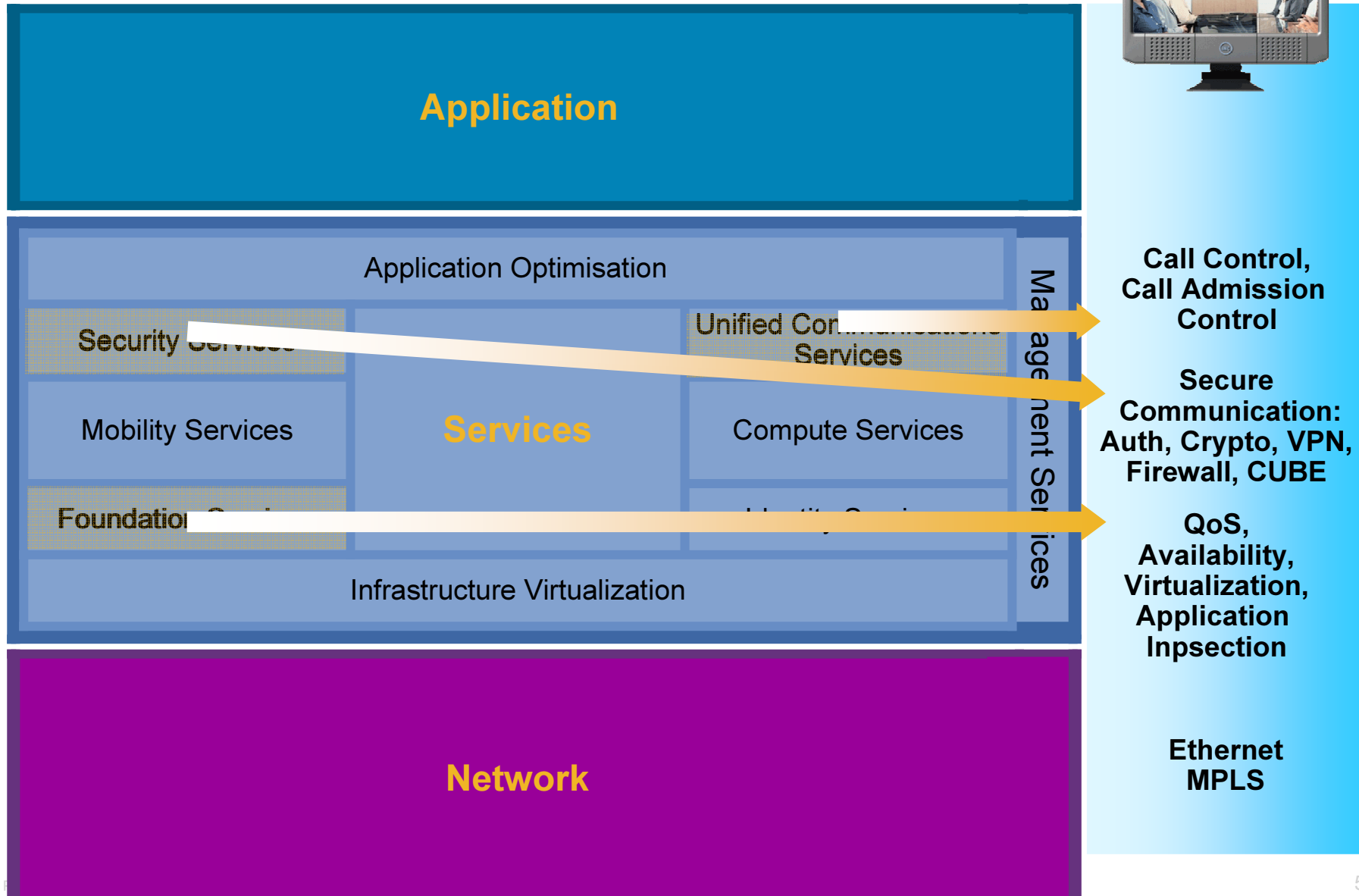
More Than Just Connectivity



An Architectural Framework



The Network as The Platform for ...



The Network as The Platform for ...



Application

Application Optimisation

Call Control,

**Services:
Interaction & Integration**

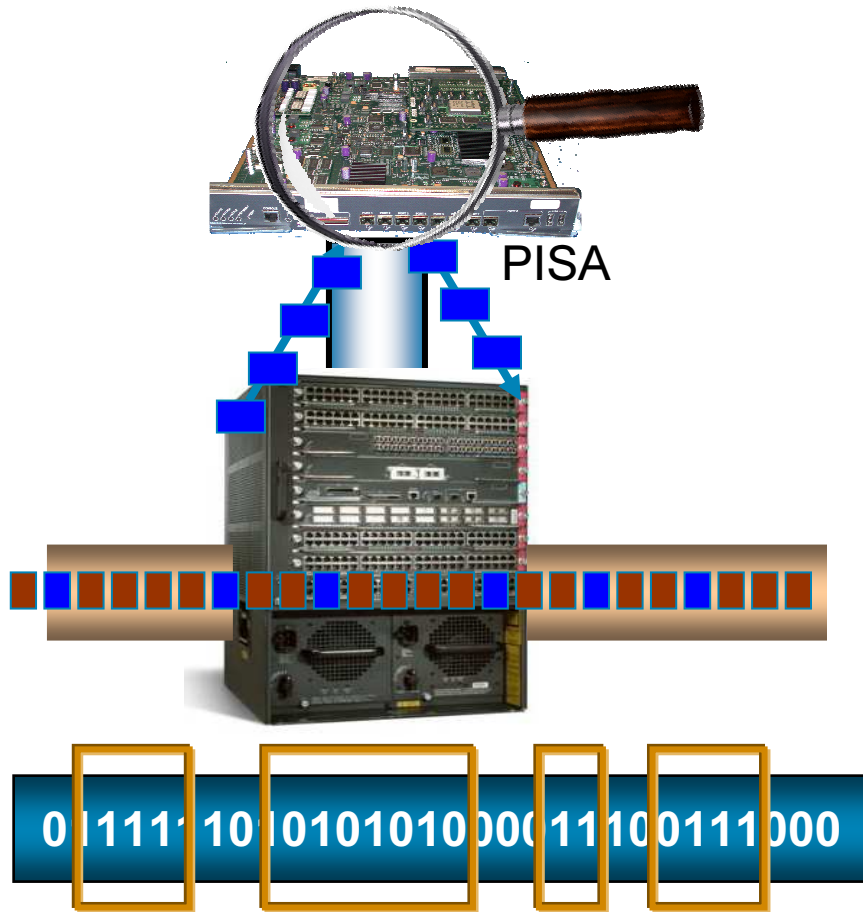
**Application
Inspection**

Network

**Ethernet
MPLS**

Architecture Example

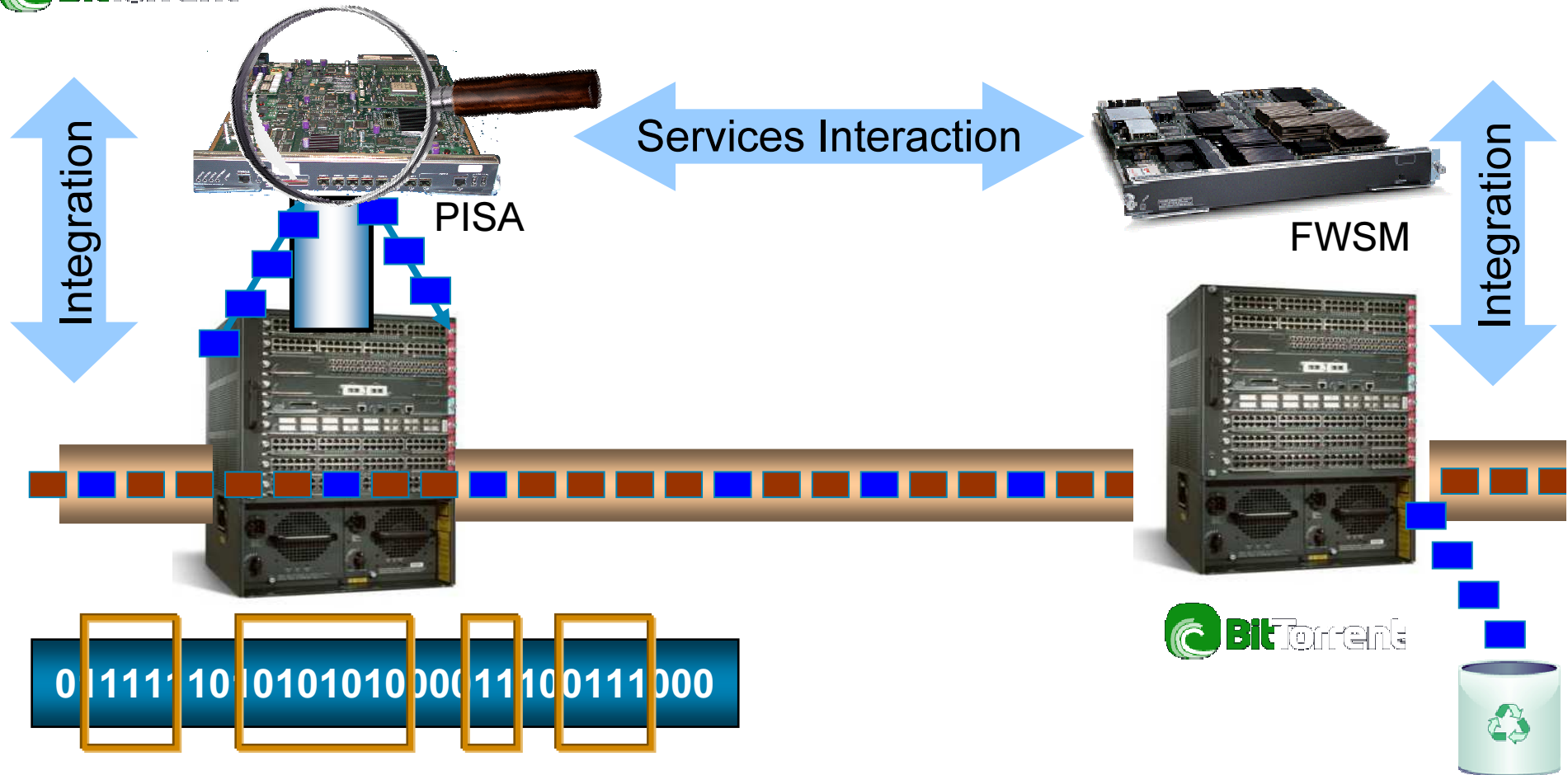
Distributed Application Intelligence



- Deep Packet Inspection for Application Recognition & Security
- Recognizes 100+ Applications
- Protection from misbehaving Applications
- Highly scalable distributed control of application flows
- Perimeter defense through flexible packet matching
- Programmatic Interface
- Policy enforcement and distributed policy enforcement

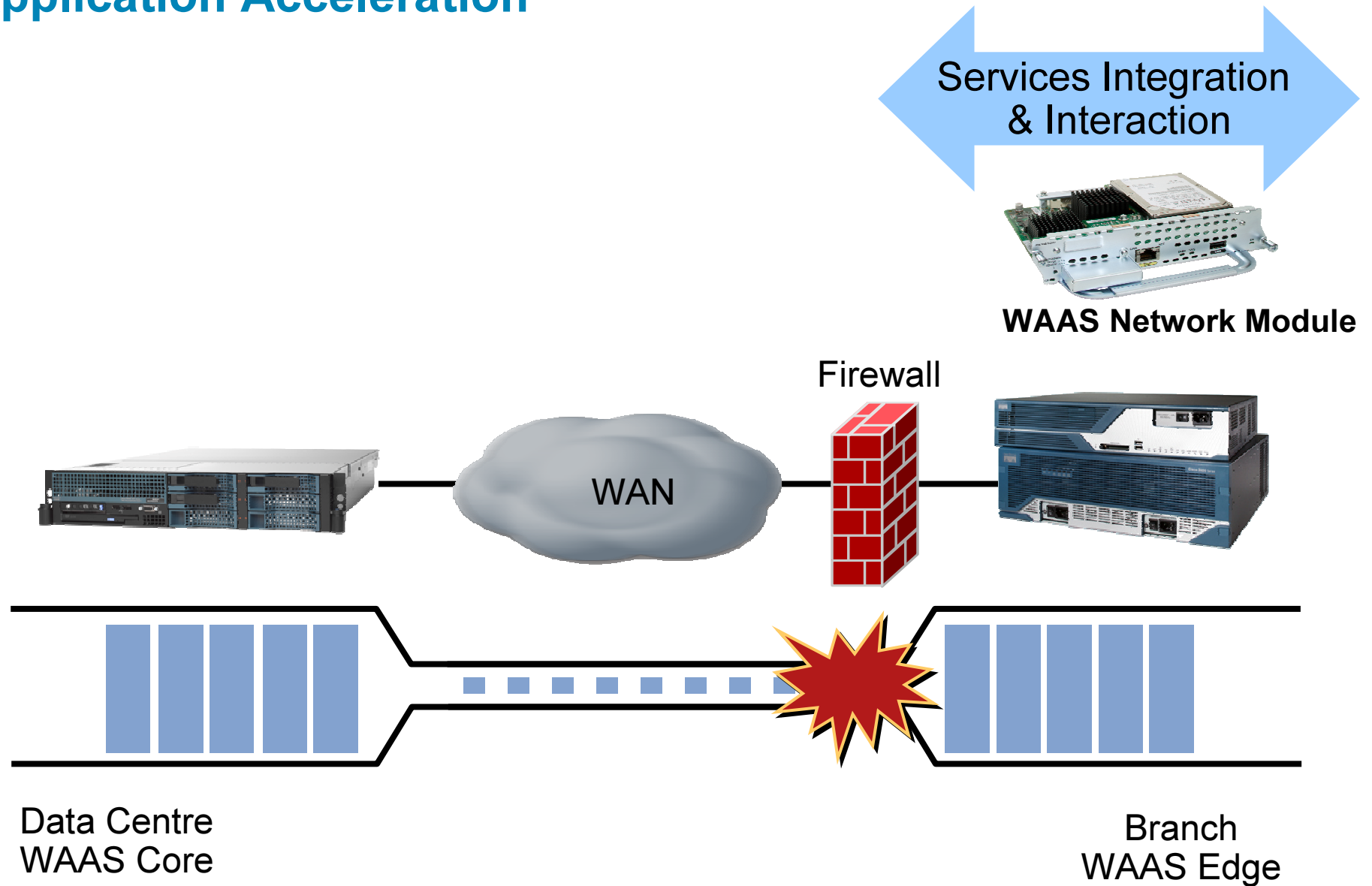
Architecture Example

Distributed Application Intelligence

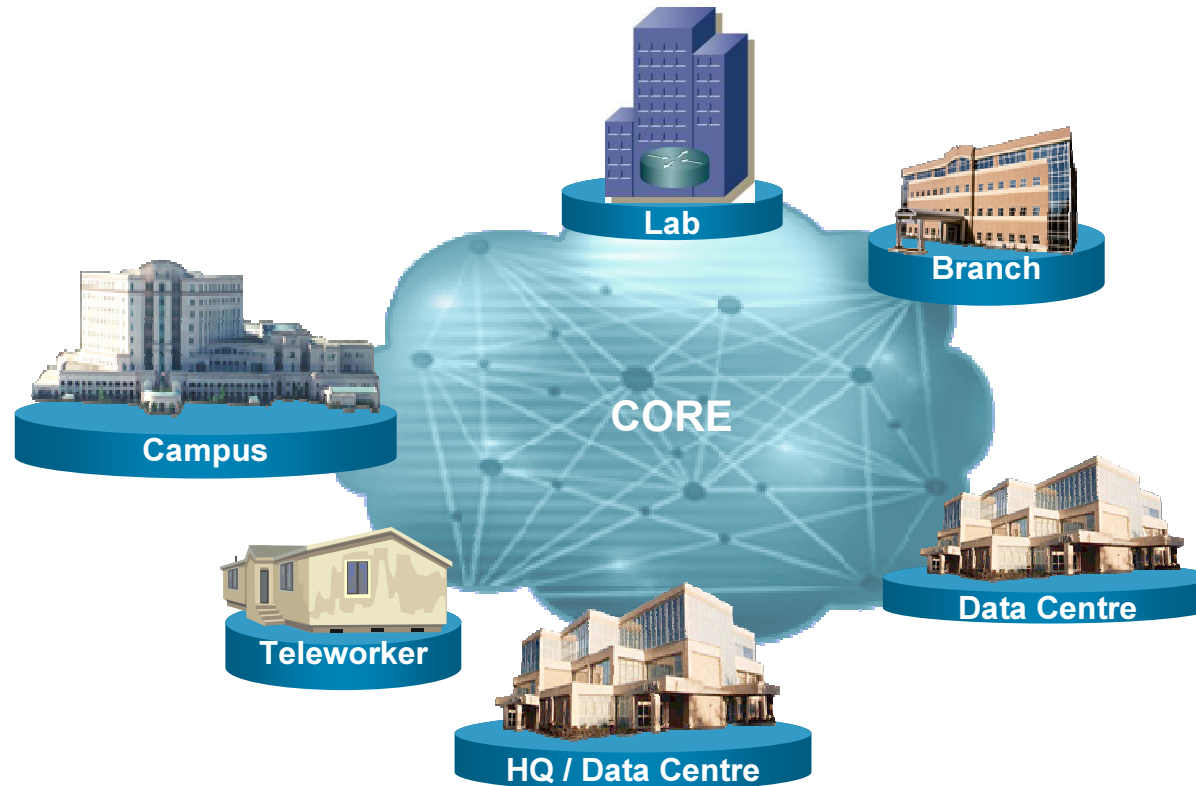


Architecture Example

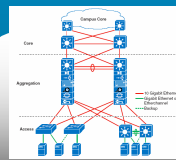
Application Acceleration



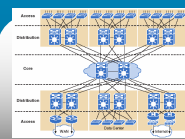
Places In The Network



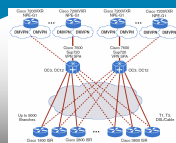
Cisco Validated Designs (CVD)



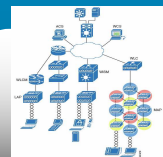
Data Centre Architecture



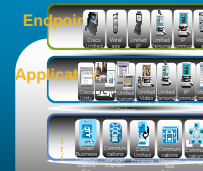
Campus Architecture



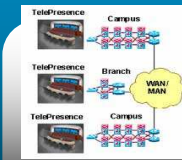
WAN/Branch Architecture



Secure Mobility Architecture



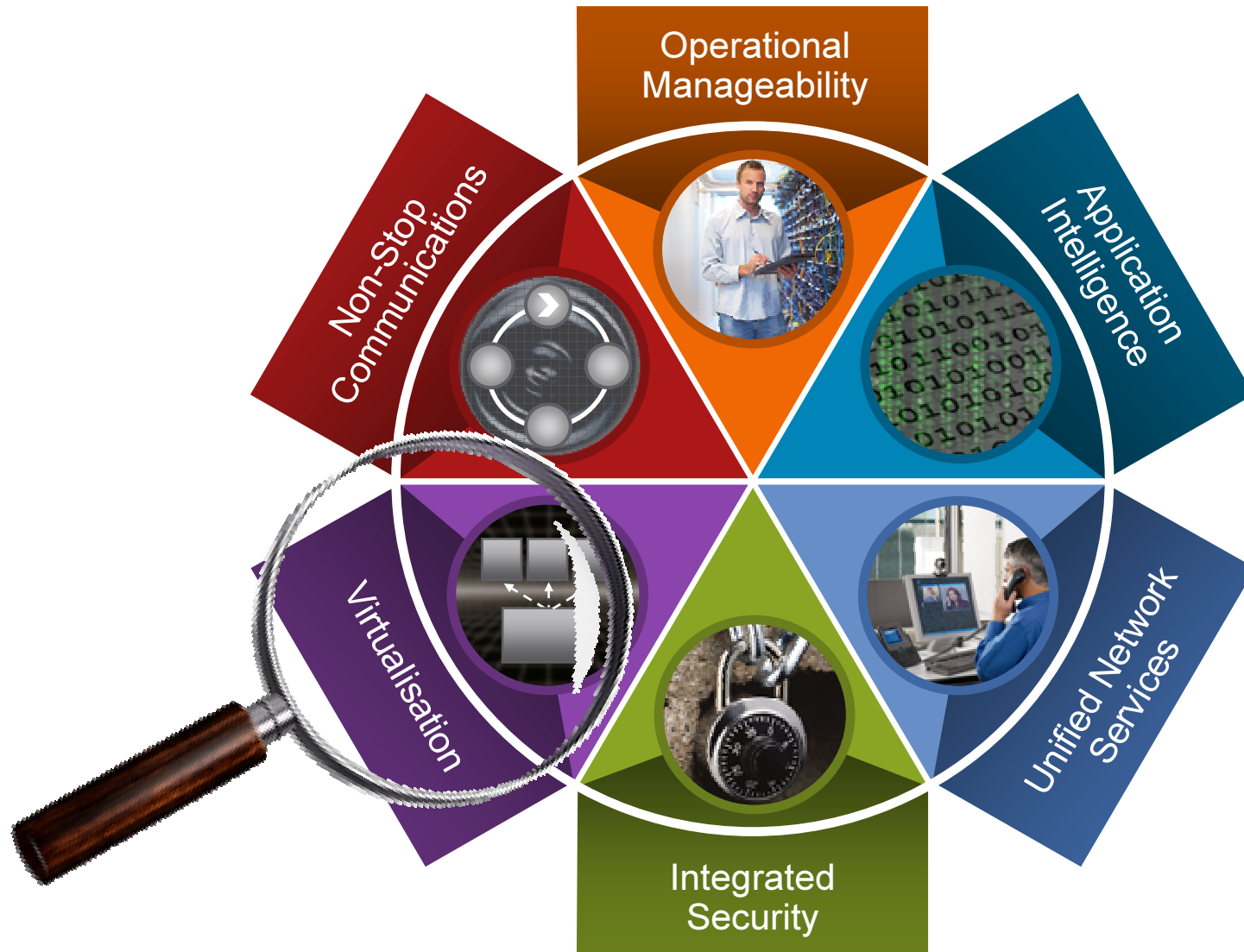
Unified Communications



Telepresence Architecture

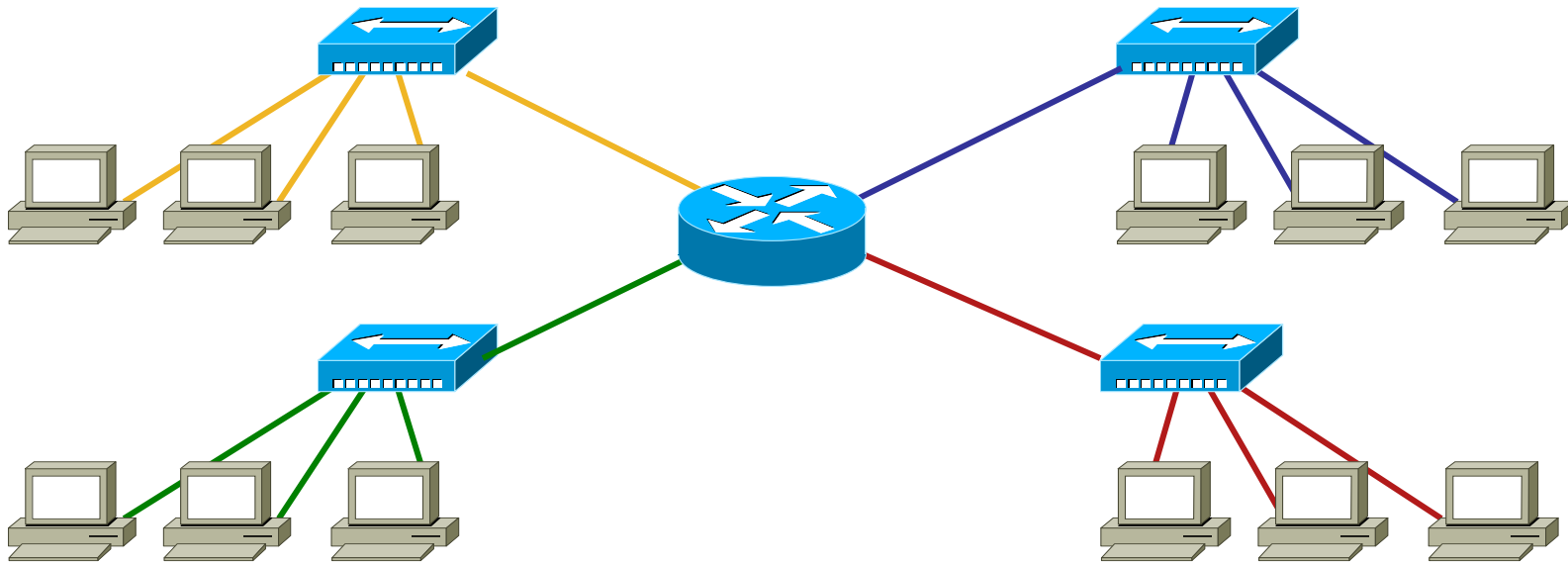
Architectural Approach

Campus Communications Fabric



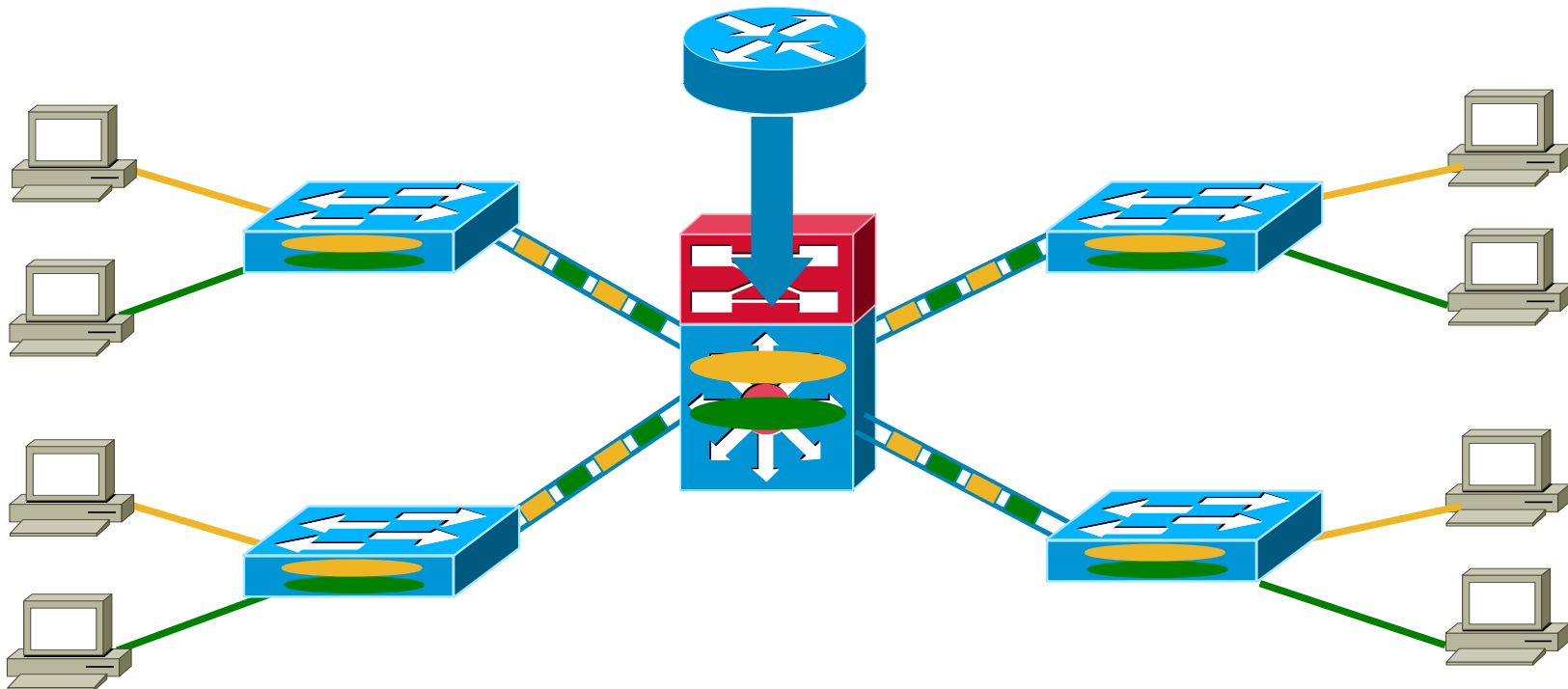
Network Virtualisation History

- **Physically separate workgroups of hub-based users**
- **As the hub islands grew, so did the requirements to connect disparate users with some form of security to police traffic between them – from this the router was born...**



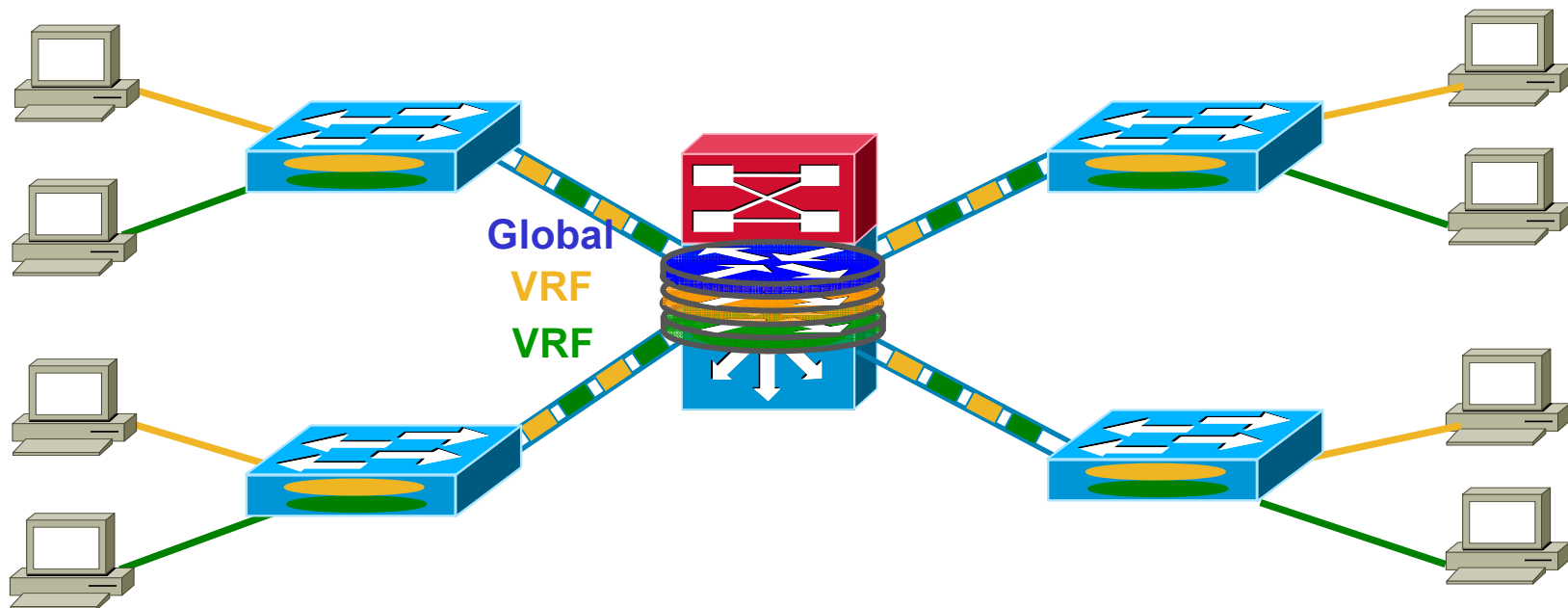
Network Virtualisation History

- **Switches with trunking provided the next evolutionary step.**
- **Trunks implemented to allow transport of VLANs across multiple devices – this provided Virtualisation at Layer 2...**
- **Router physically ‘Virtualised’ in switch**



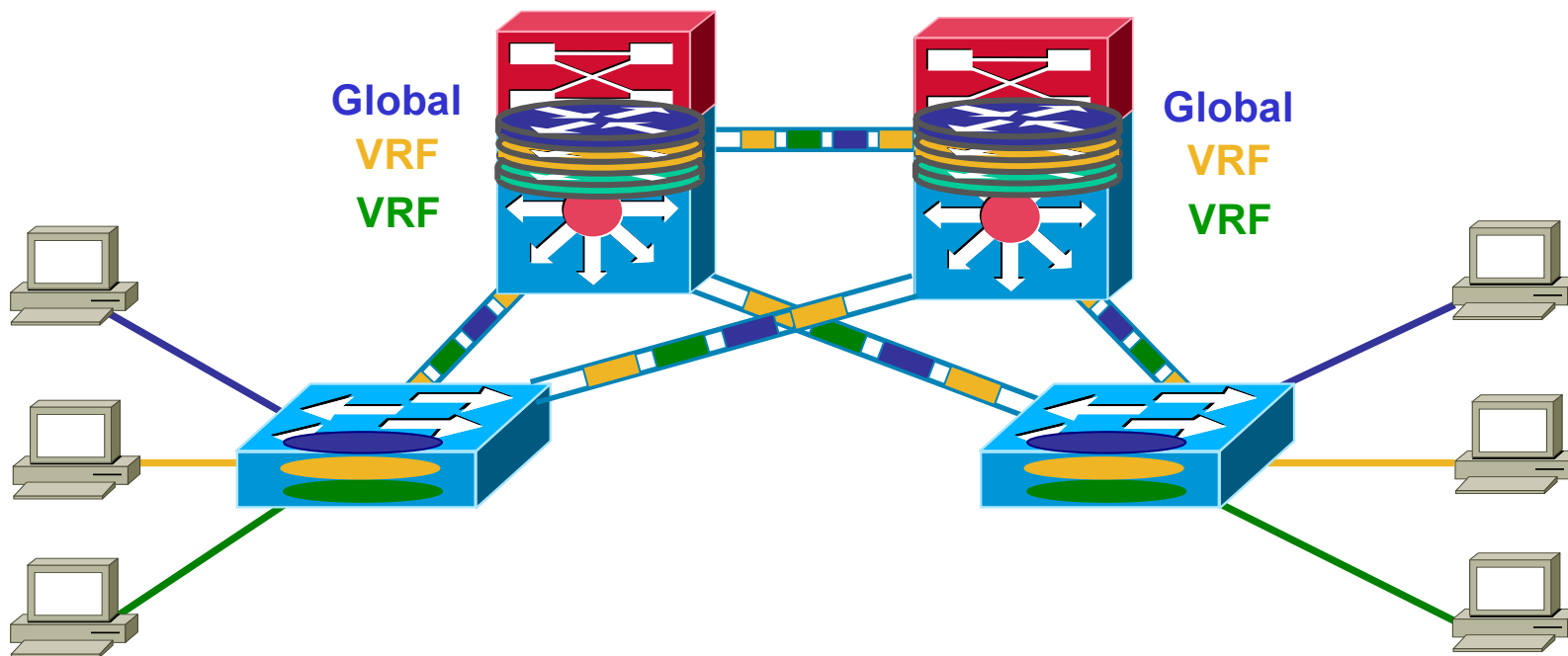
Network Virtualisation History

- Layer 3 Virtualisation appeared in the form of Virtual Routing and Forwarding (VRF) which allowed a Layer 3 device to virtualize routing tables giving each VRF each it's logical view of the network ...



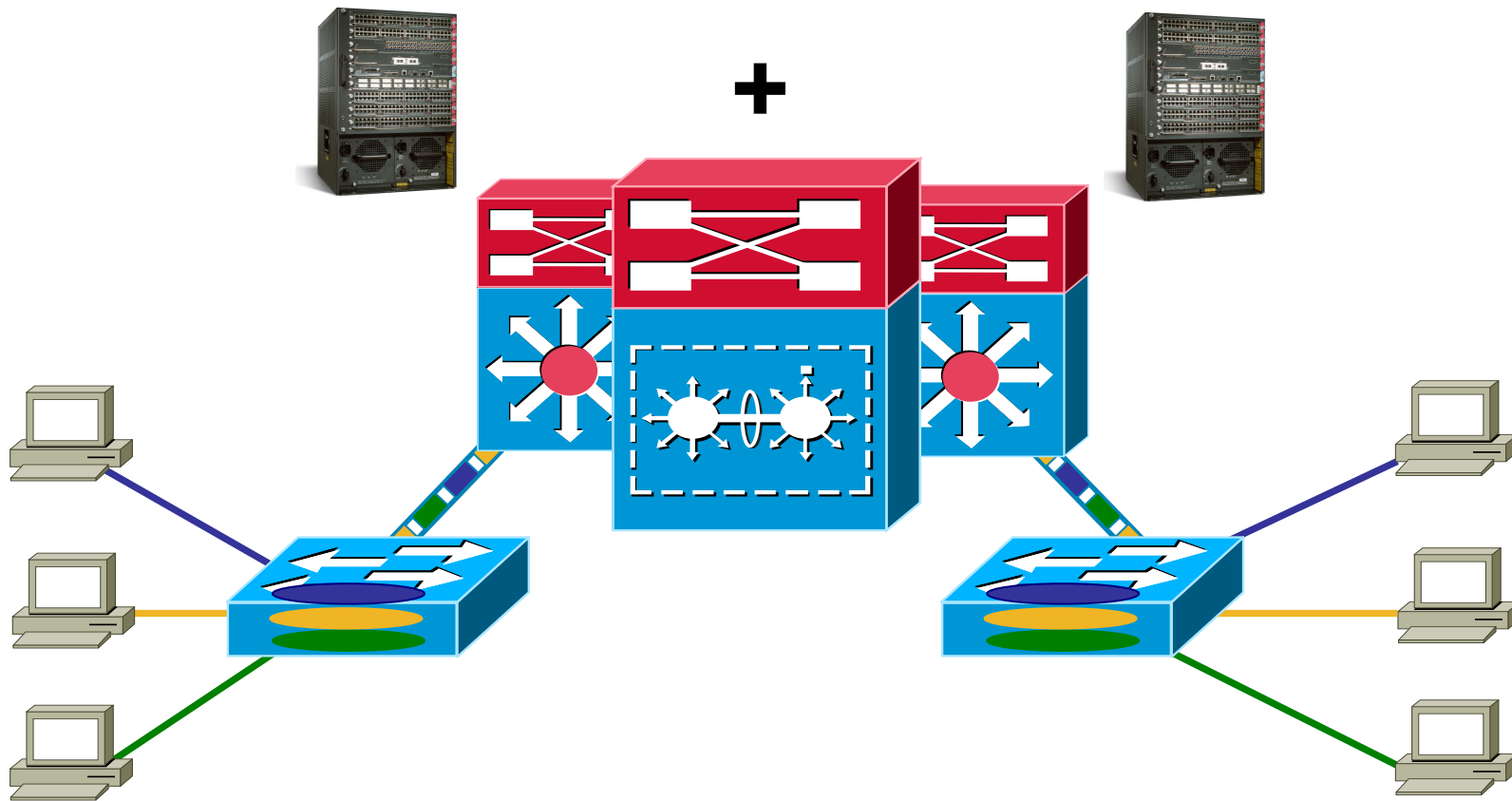
Network Virtualisation History

- **High Availability requirements meant dual homing and associated complexities.**



Network Virtualisation History

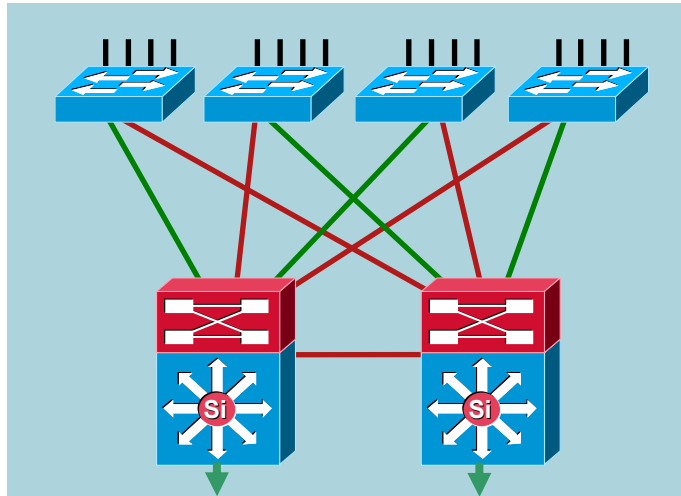
- **Virtual Switching is the latest addition to the Virtualisation story allowing 2 devices to operate as a single logical network device**



High Availability Campus Design

Access

Distribution



Complex suite of protocols – Spanning Tree, HSRP, VRRP

Non-deterministic, Stateless Recovery From Failure

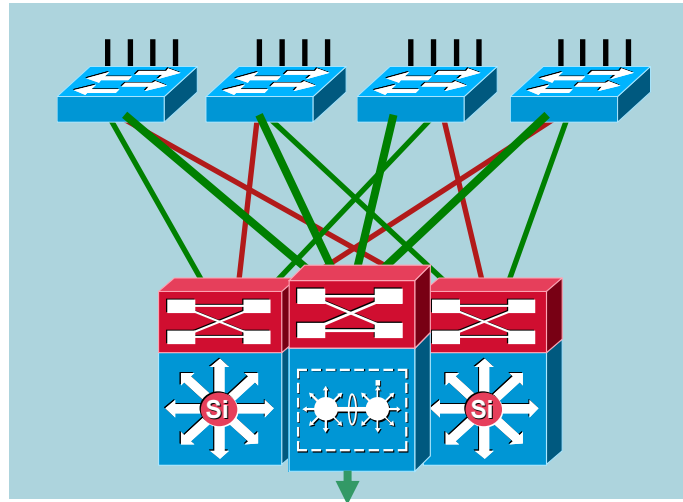
Inefficient Resource Utilization

Increased Management Burden

High Availability Campus Design

Access

Distribution



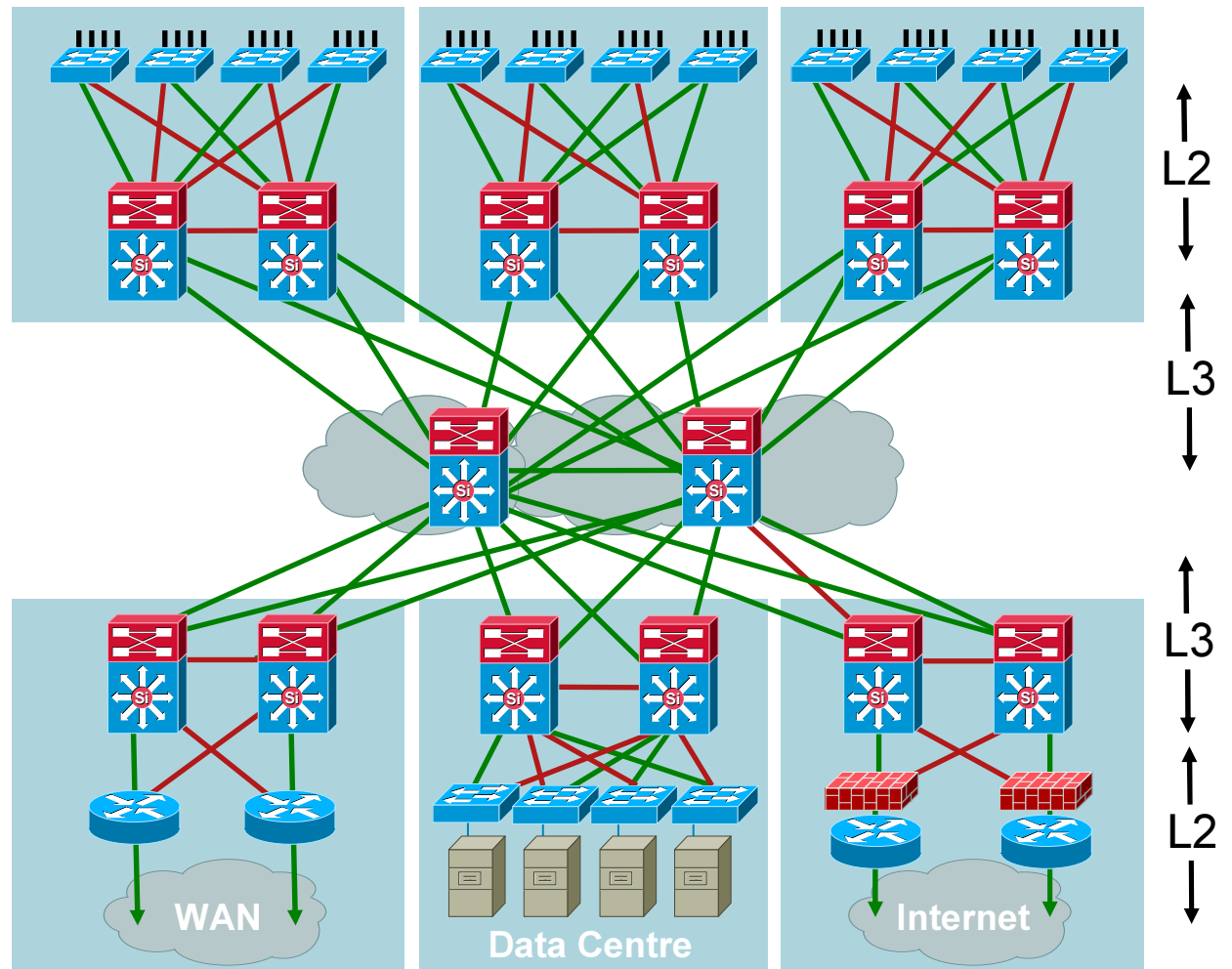
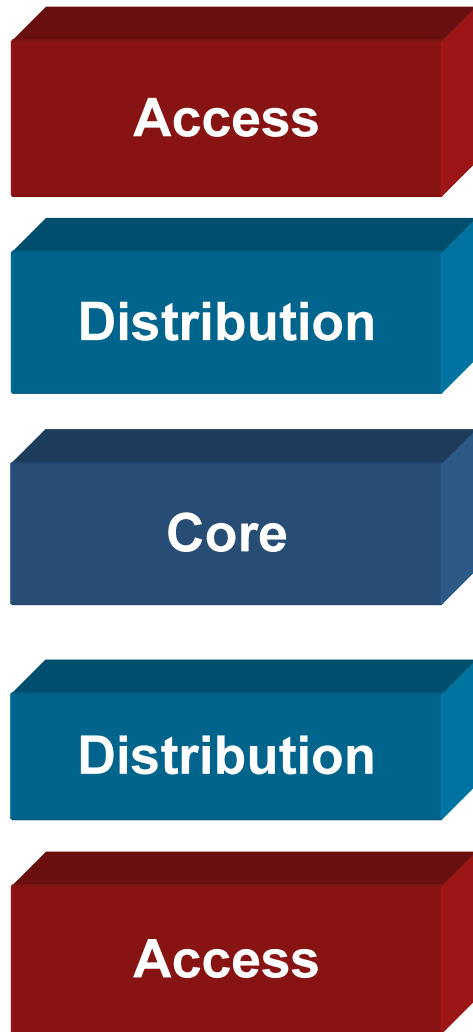
Eliminate complex suite of protocols

Stateful Deterministic Inter-chassis Recovery From Failure

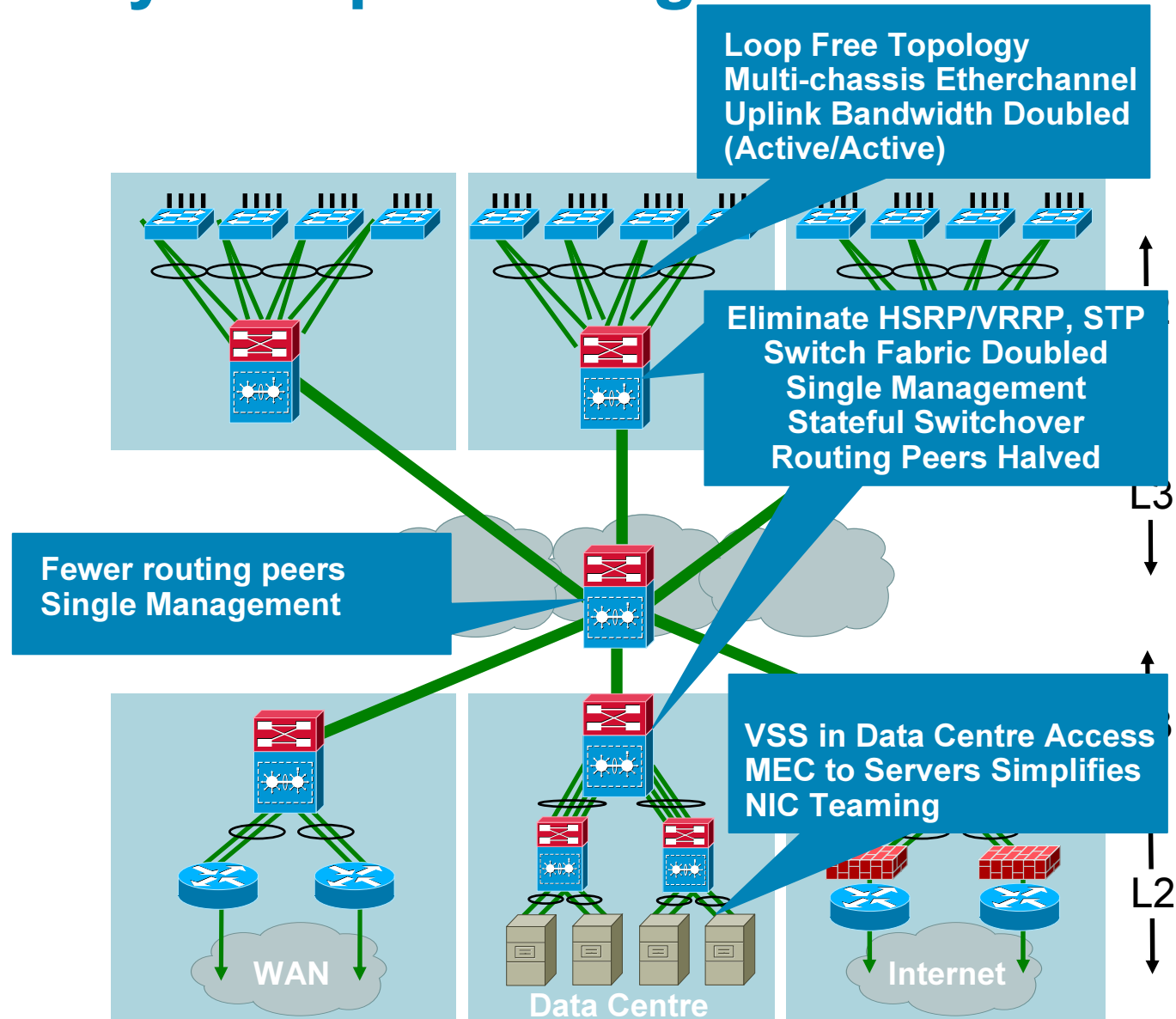
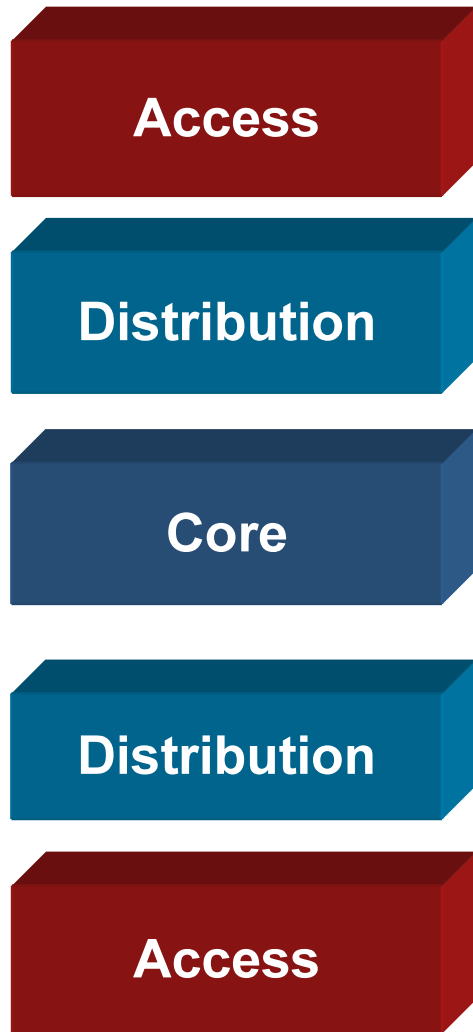
Double Fabric Bandwidth

Reduced Management Burden

High Availability Campus Design



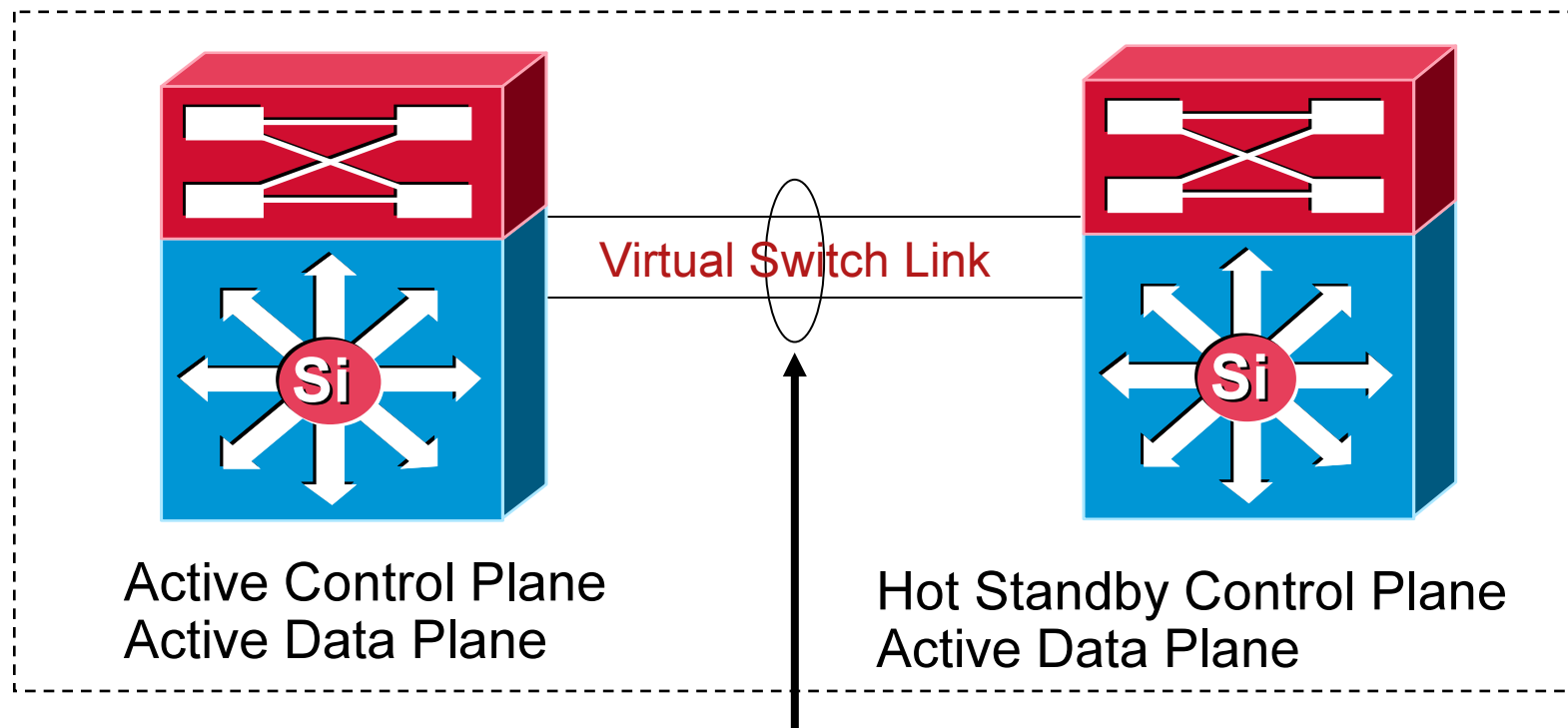
High Availability Campus Design



Virtual Switch System Concepts

Virtual Switch Domain

Defines 2 switches participating together as a single virtual switching system

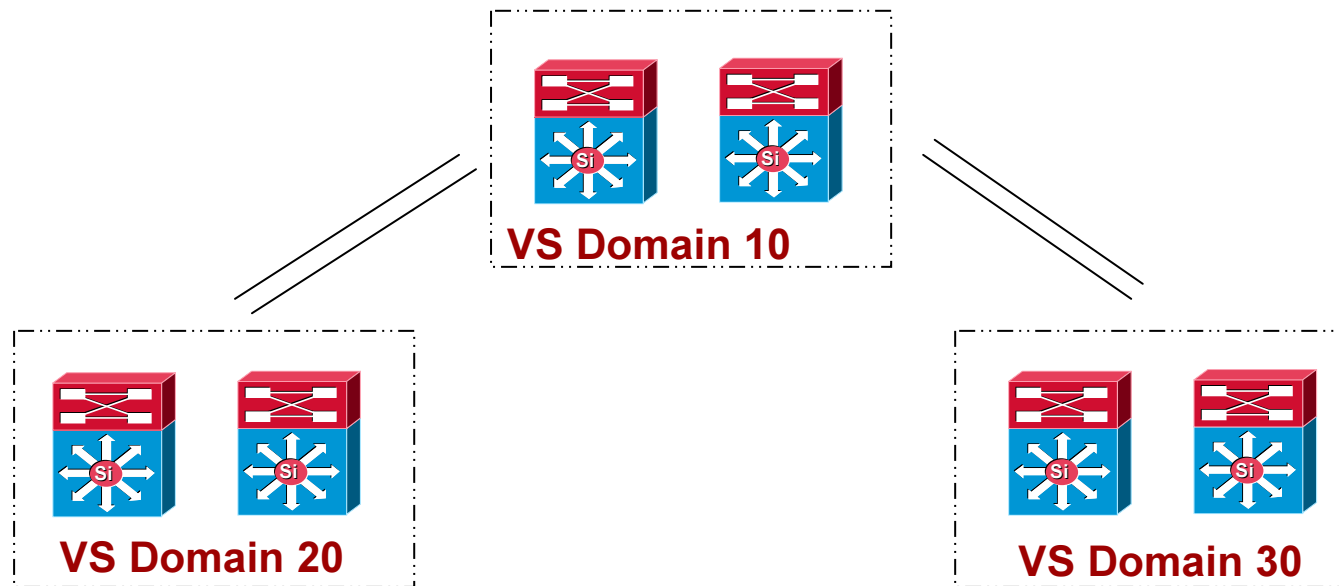


Special 10GbE link bundle joining the 2 switches allowing them to operate as a single device

Virtual Switch System Concepts

Virtual Switch Domain

A Virtual Switch Domain ID represents the logical grouping the 2 physical chassis within a VSS. It is possible to have multiple self-contained VS Domains throughout the network...



The configurable values for the domain ID are 1-255. It is always recommended to use a unique VS Domain ID for each VS Domain throughout the network...

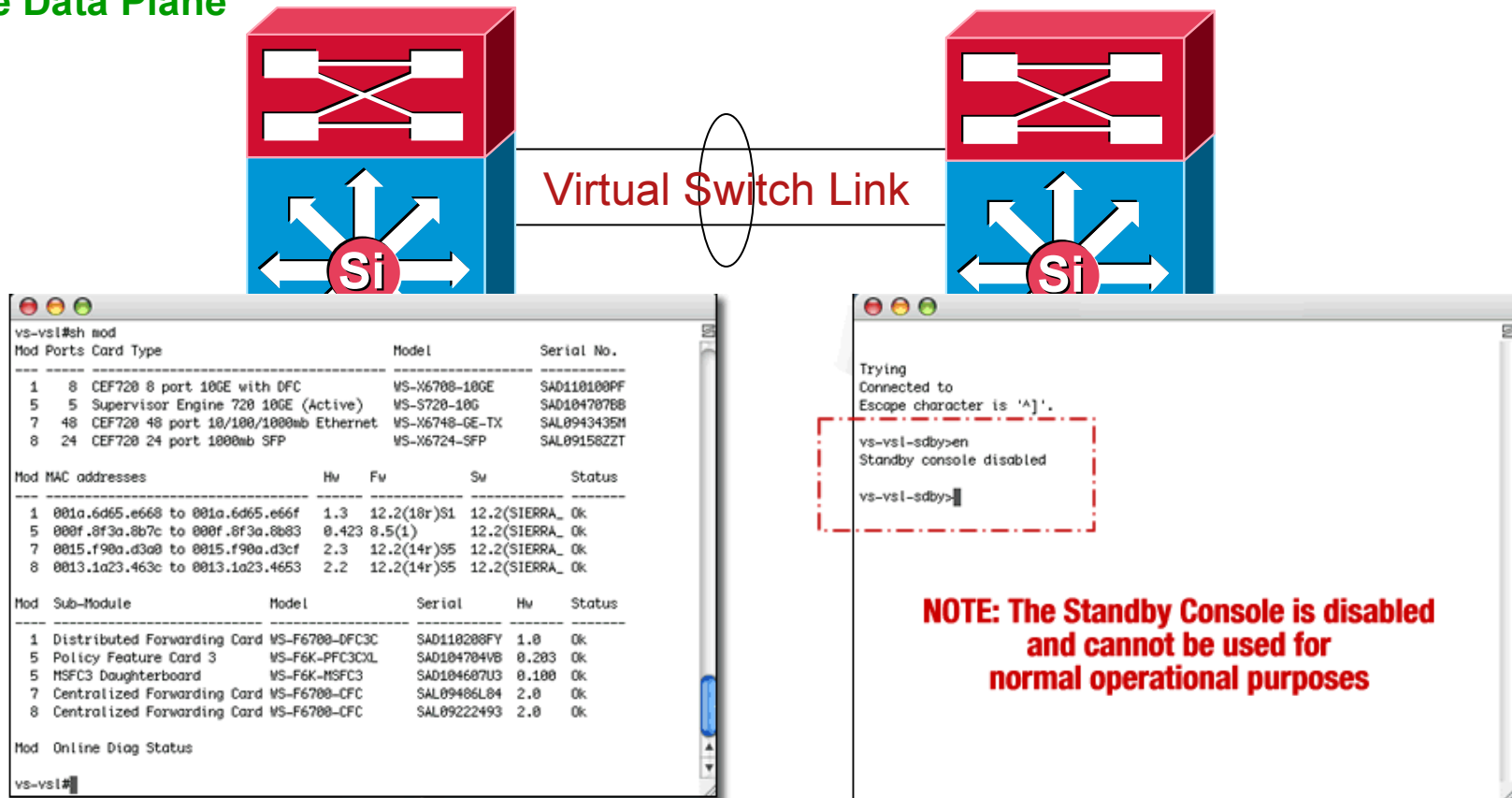
Virtual Switch System Concepts

Control Plane

Both data planes are active and forwarding, but only one control plane – hence there is only a single management point

Virtual Switch Active
Active Control Plane
Active Data Plane

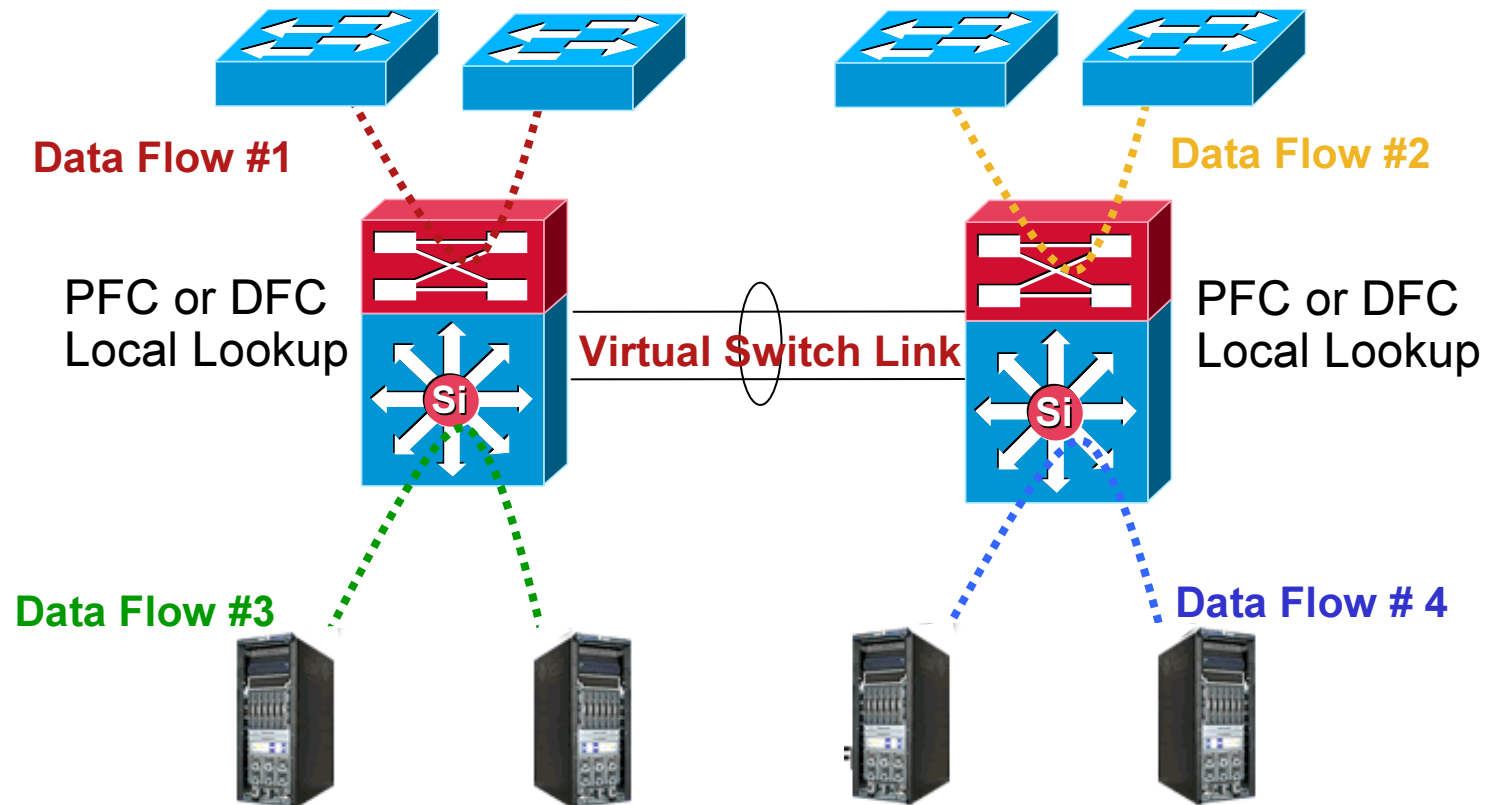
Virtual Switch Standby
Standby Control Plane
Active Data Plane



Virtual Switch System Concepts

Data Plane

Data Planes in both switches are active - each has a full copy of the forwarding tables and Security/QoS policies in hardware such that each can make a fully informed local forwarding decision...



Virtual Switch Architecture

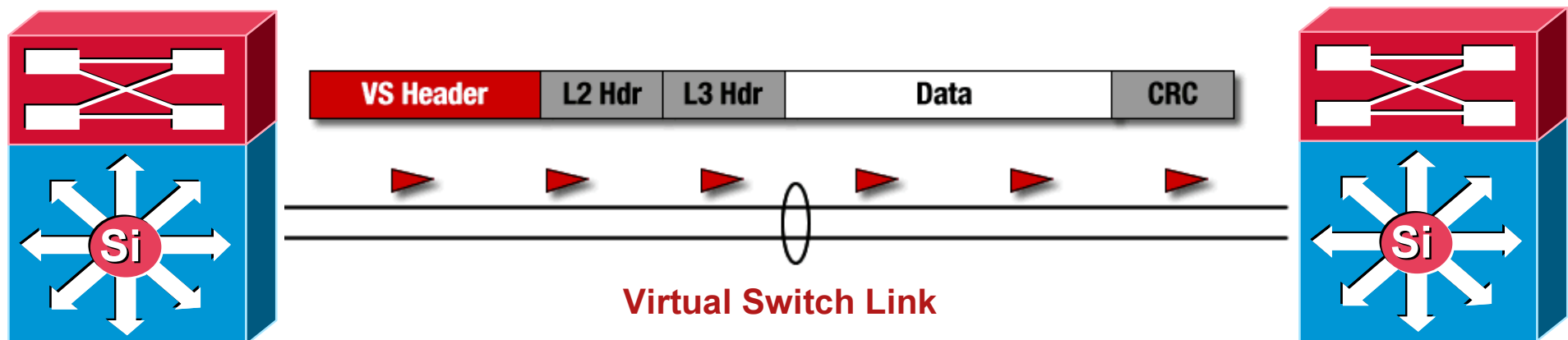
Virtual Switch Link

The Virtual Switch Link extends the out of band channel allowing the active control plane to manage the hardware in the second chassis...

A Virtual Switch Link Bundle can consist of up to 8 x 10GbE Links

All traffic traversing VSL is encapsulated in 32-byte header containing ingress & egress switch port indexes, CoS, VLAN ID & other important L2, L3 information

Control plane uses VSL for CPU to CPU communications while data plane uses VSL to extend internal switch fabric to external chassis

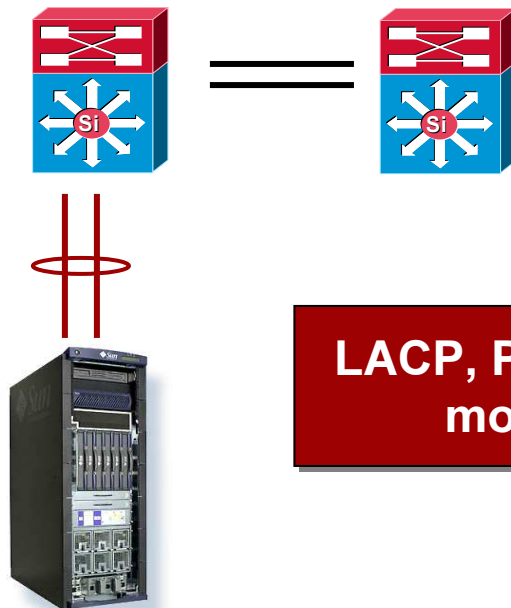


Virtual Switch Architecture

Multichassis EtherChannel (MEC)

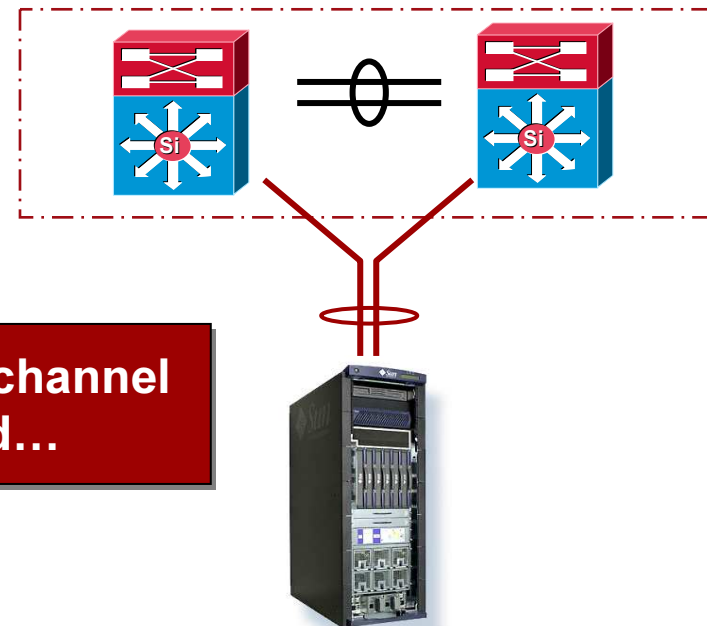
In a Virtual Switch environment, 2 physical switches form a single logical entity. Etherchannels can now also be extended across the 2 physical chassis

Traditional Switch Pair



Regular Etherchannel on single chassis

Virtual Switch

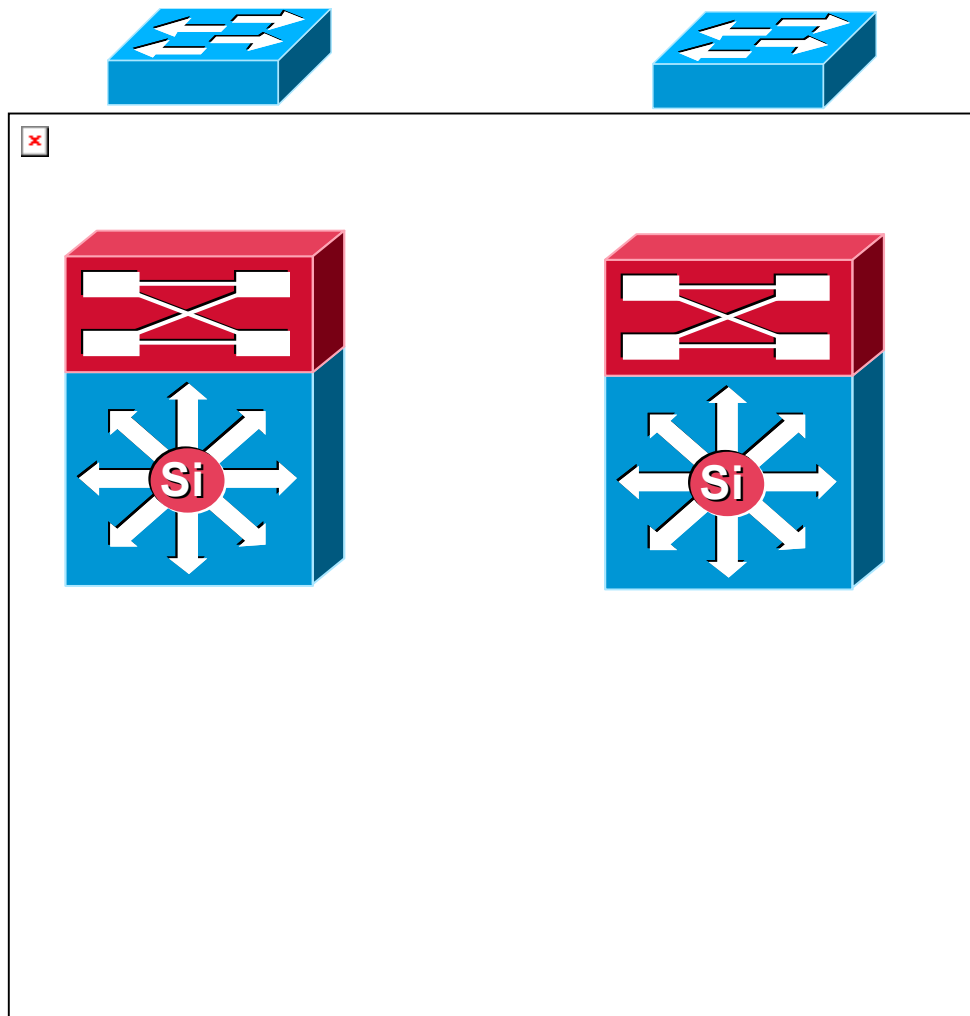


Multichassis EtherChannel across 2 VSL-enabled Chassis

LACP, PAGP or ON Etherchannel modes are supported...

Virtual Switch Architecture

Multi Chassis Etherchannel (MEC)



Multi Chassis Etherchannel (MEC)

Introduces a new deployment option for improving link resiliency

Allows an Etherchannel link to be terminated across 2 physical chassis

Up to 8 links support in a single multi-chassis Etherchannel

MEC is supported with both 802.3ad and Cisco PAGP

Attached host sees other end (Virtual Switch) as a single device

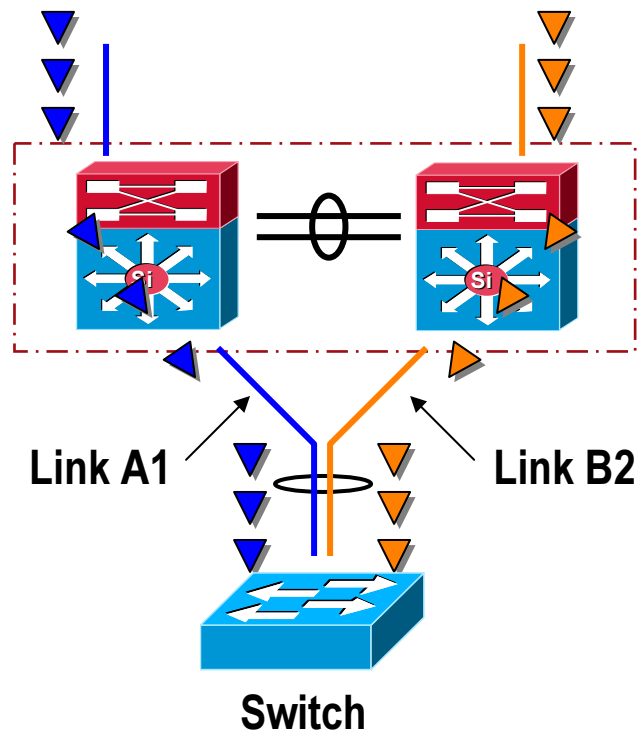
Etherchannel hash has been modified so in any bundle we always prefer local link over a link on other chassis

Virtual Switch Architecture

Etherchannel Hash for MEC

Deciding on which link of a Multi-chassis Etherchannel to use in a Virtual Switch is skewed in favour of 'switch-local' links in the bundle to avoid overloading the Virtual Switch Link (VSL) with unnecessary traffic loads...

Blue Traffic destined for the Switch will result in Link **A1** in the MEC link bundle being chosen as the destination path...

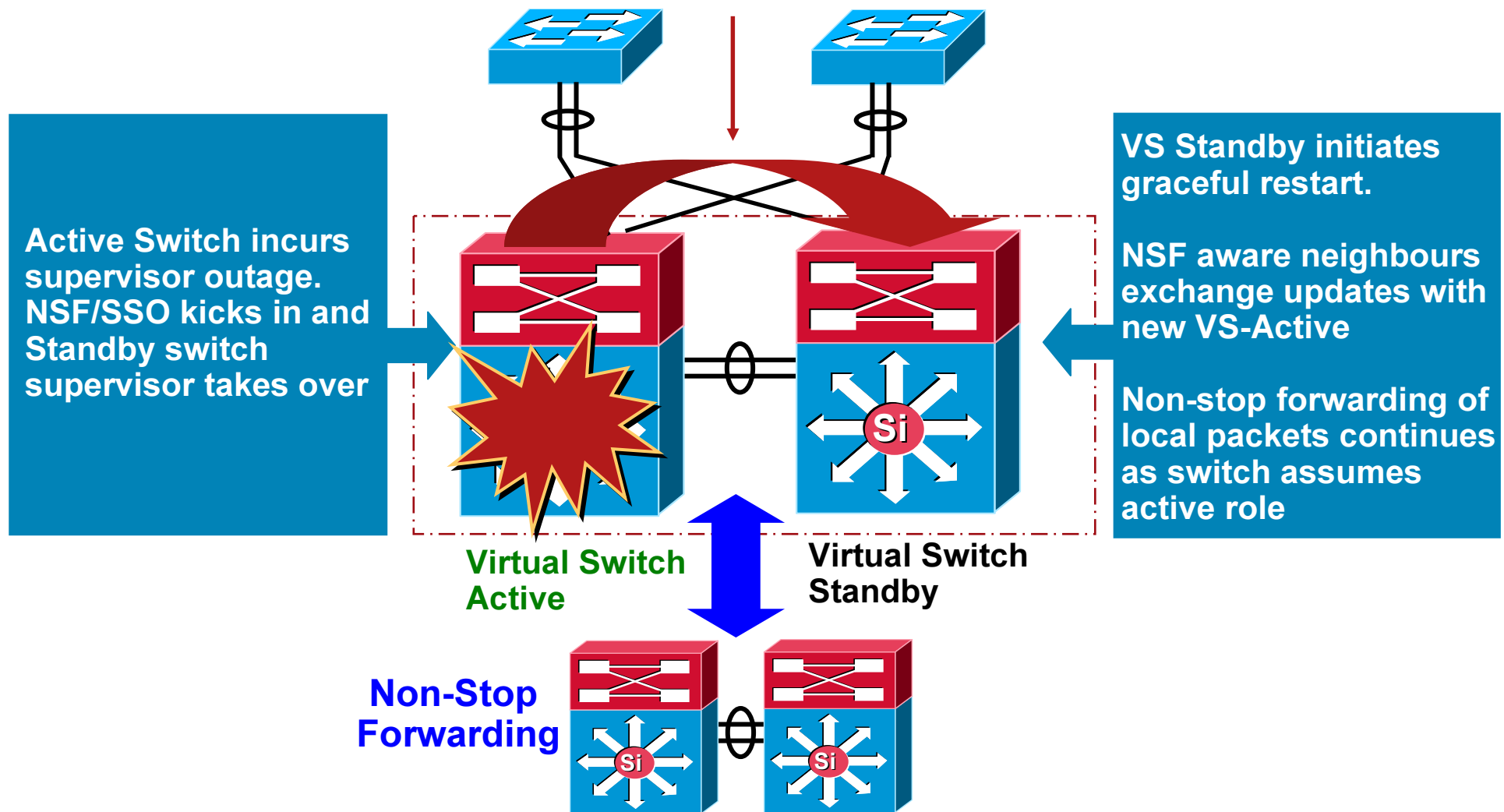


Orange Traffic destined for the Switch will result in Link **B2** in the MEC link bundle being chosen as the destination path...

Virtual Switch Architecture

Inter Chassis NSF/SSO

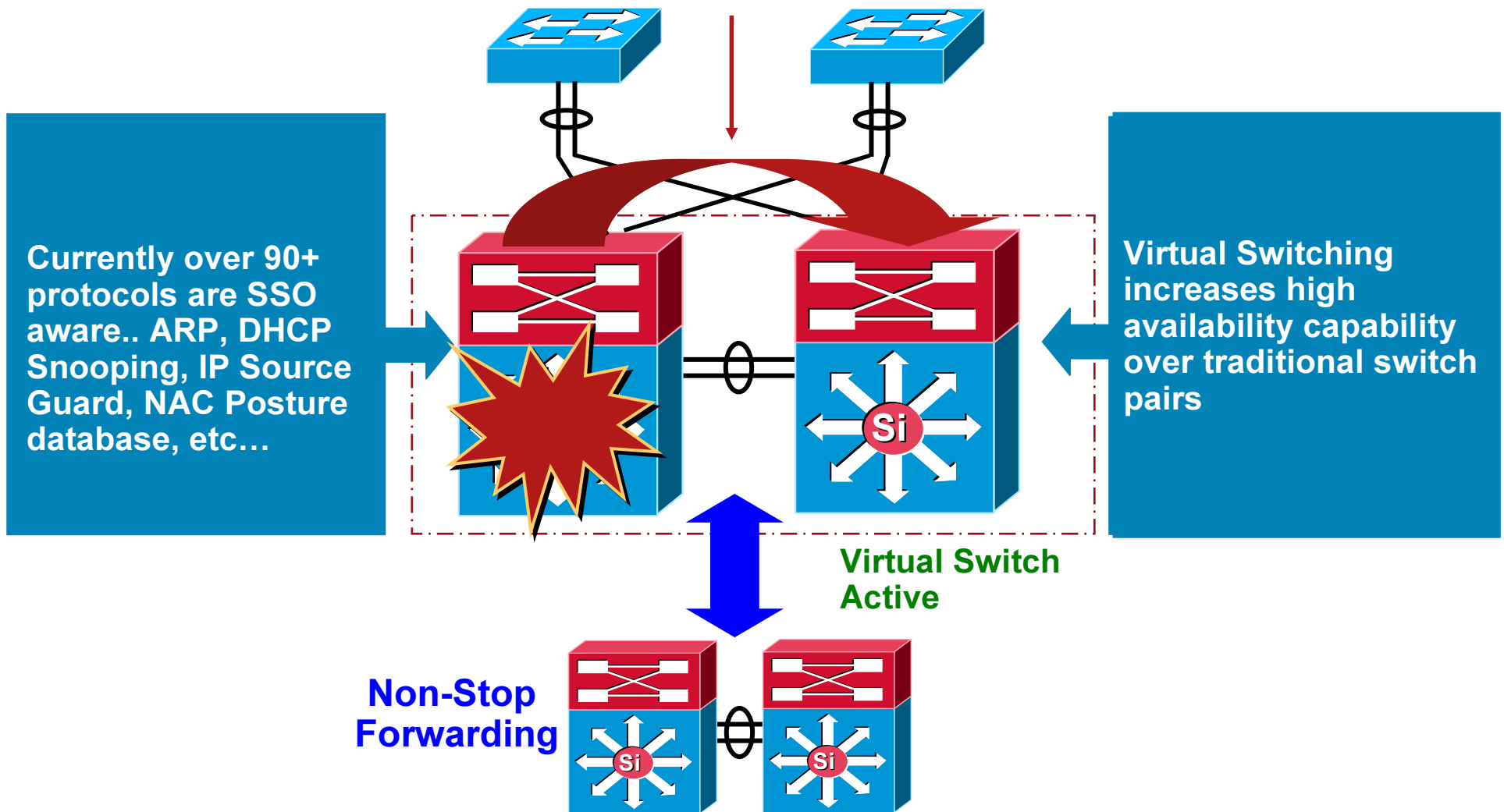
Stateful Switchover (SSO)



Virtual Switch System

Inter Chassis NSF/SSO

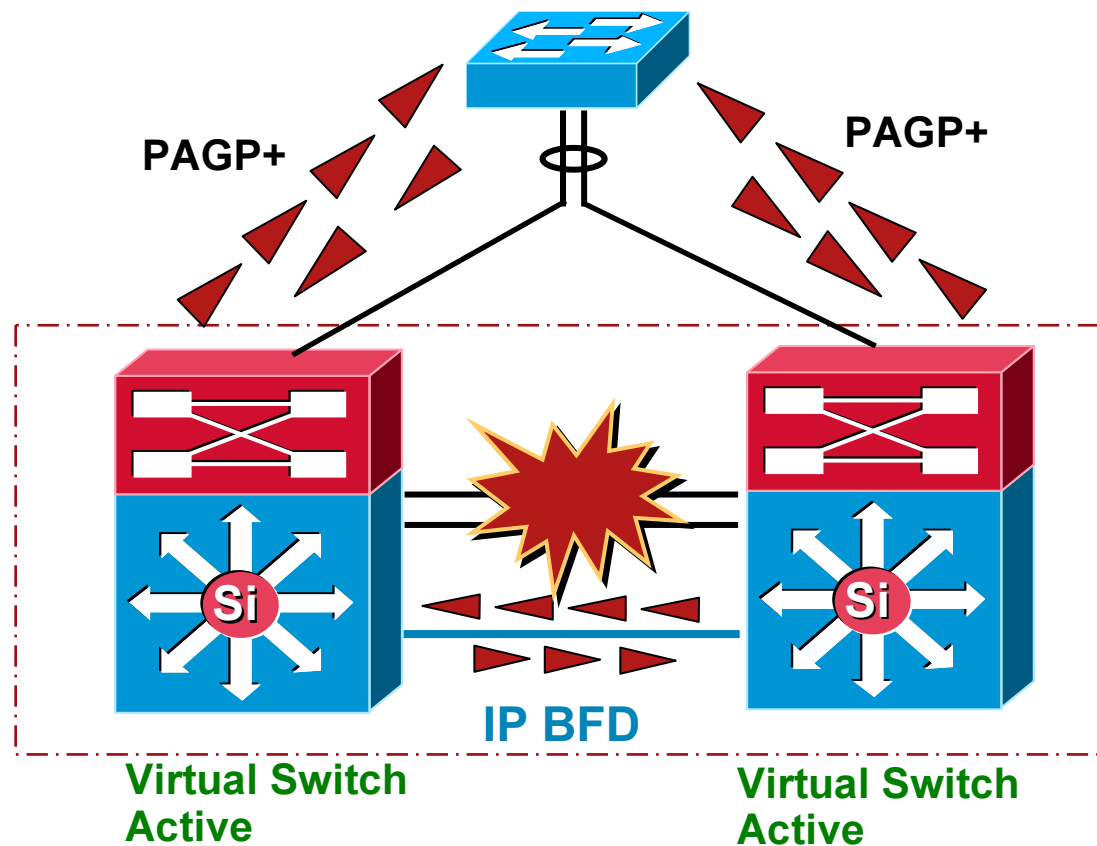
Stateful Switchover (SSO)



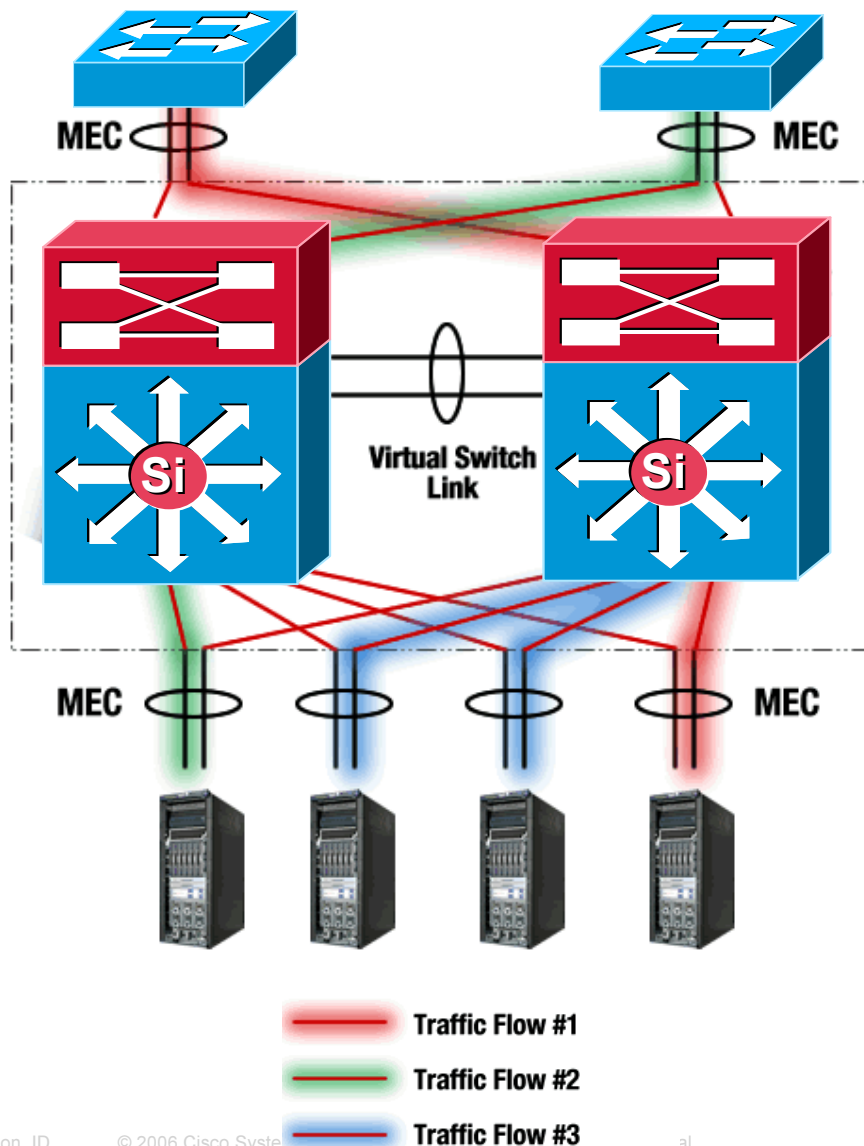
Virtual Switch System

Dual Active

Two mechanisms to recover if Virtual Switch Link Fails: PAGP +, IP-BFD
Either mechanism can be used to detect active peer.



Virtual Switch System Deployment Considerations



Deployment Considerations

Connected nodes should **ALWAYS** be dual homed

Etherchannel hash modified so local link always preferred over remote link in same bundle

ECMP has been modified to choose local link over remote link in same bundle

Up to 128 Etherchannel link bundles in virtual switch domain

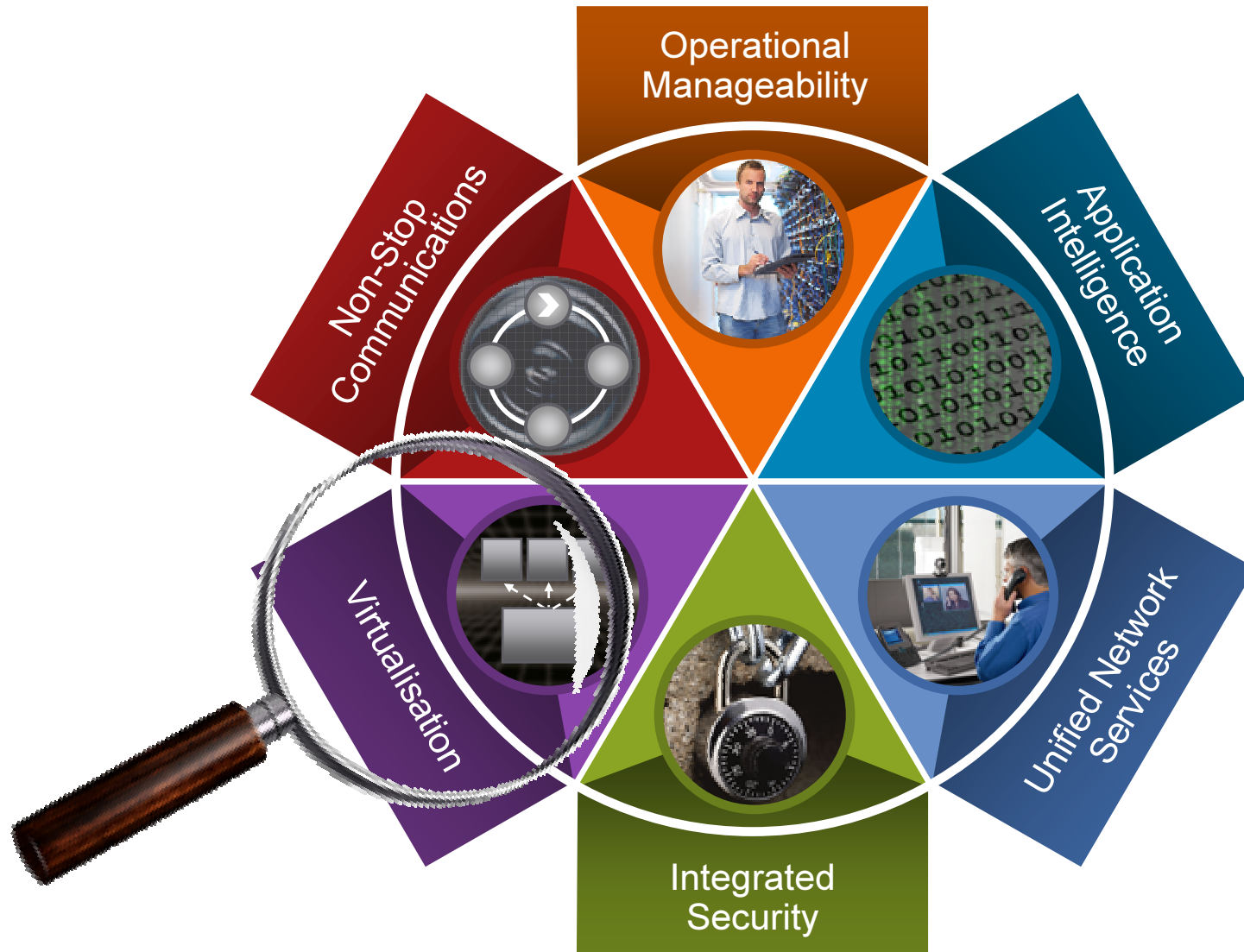
Minimal traffic crosses VSL

Size number of links in VSL to meet local redundancy and traffic requirements

Hardware/Software Dependencies

Architectural Approach

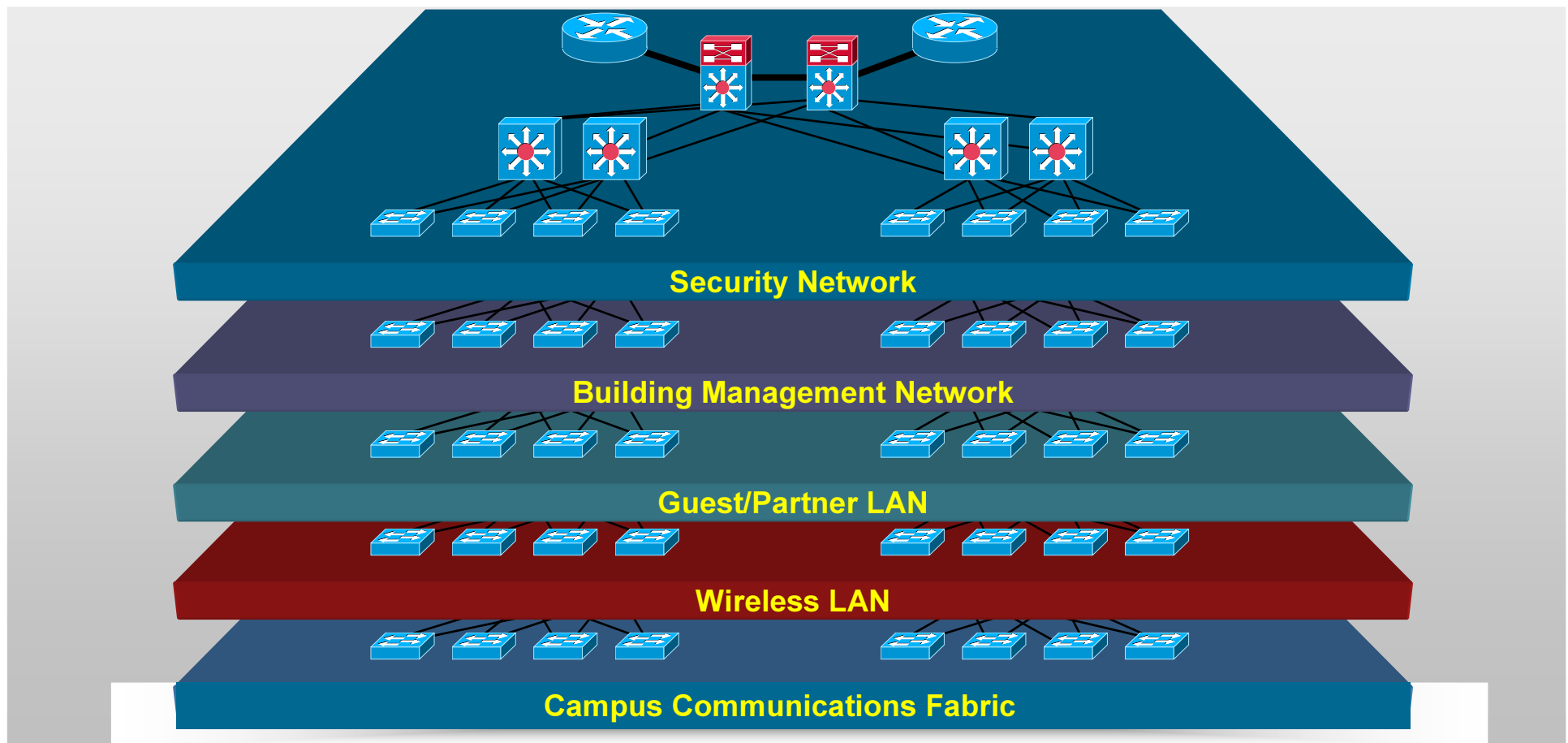
Campus Communications Fabric



Network Virtualisation

Many to One

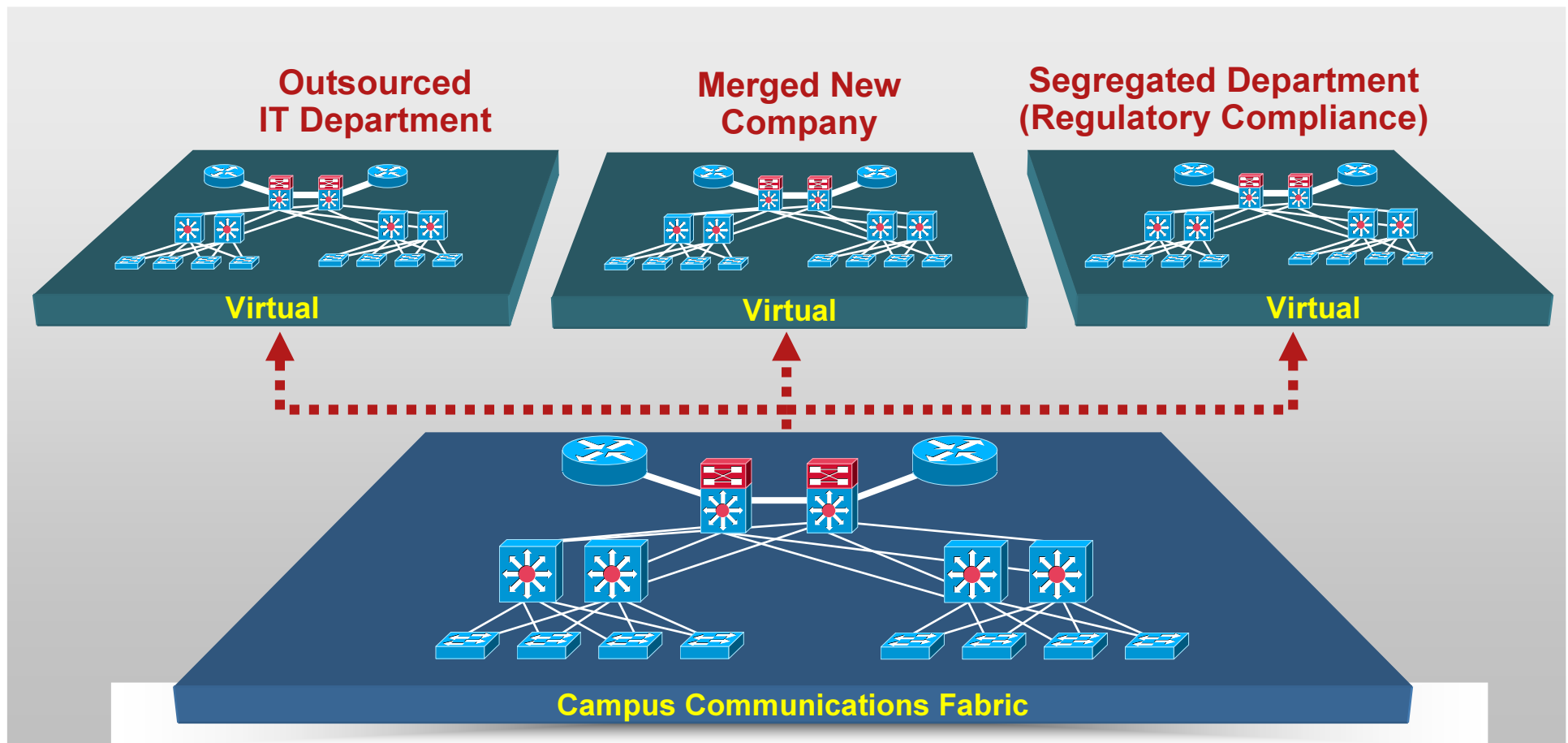
One network consolidates many physical networks



Network Virtualisation

One to Many

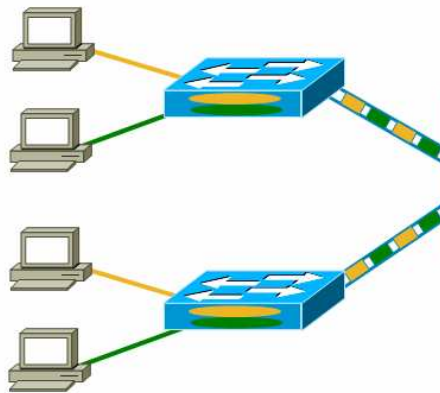
One network supports many virtual networks



History Repeated

Network Virtualization History

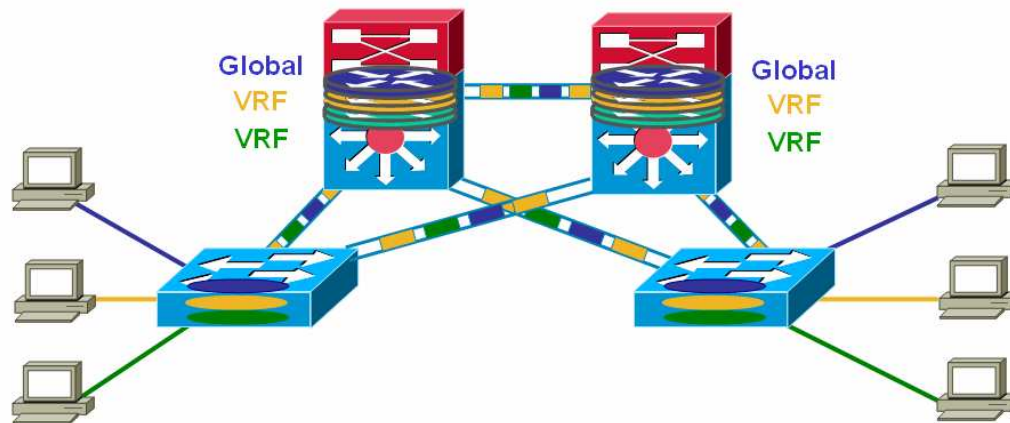
- The Layer 3 Switch with the evolutionary step – to achieve connectivity throughout the network – implemented to allow the devices – this provided v



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Network Virtualisation History

- High Availability requirements meant dual homing and associated complexities.

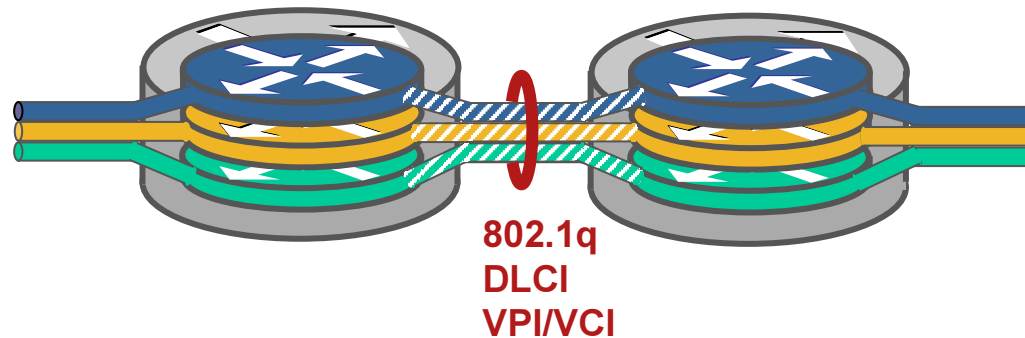


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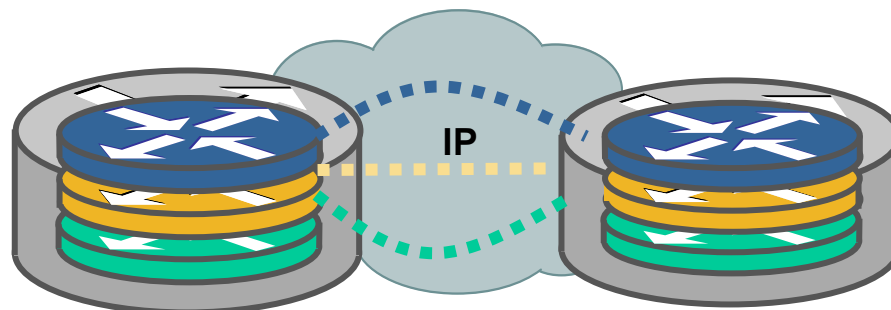
Data Path Virtualisation Techniques

Single Hop Data Path Virtualisation



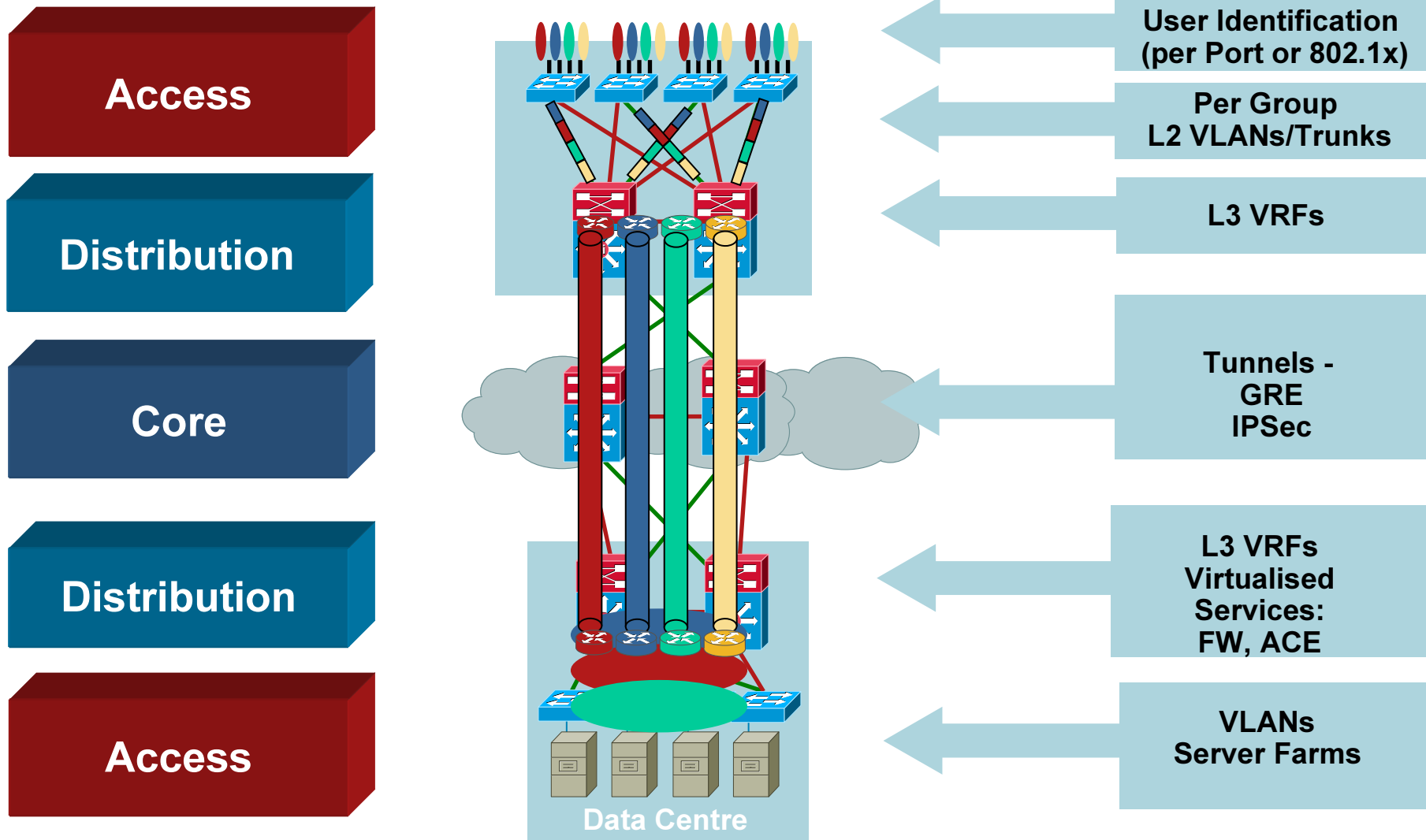
- Tags
 - 802.1q
 - Others (DSCP etc.)
 - More about Tags Later !
- Virtual circuits
 - ATM
 - Frame Relay
 - AToM L2 Circuits

Multi-Hop Data Path Virtualisation



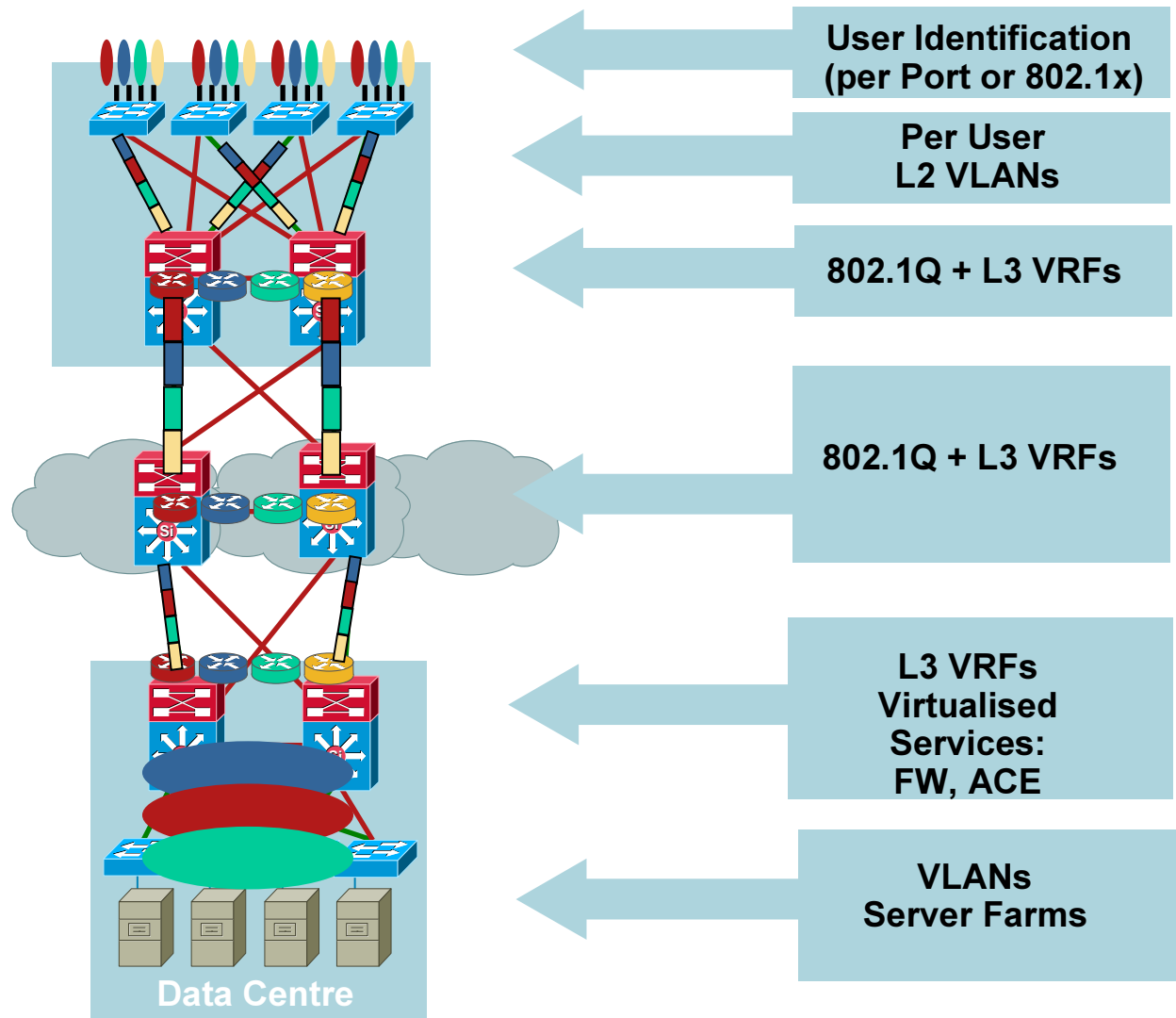
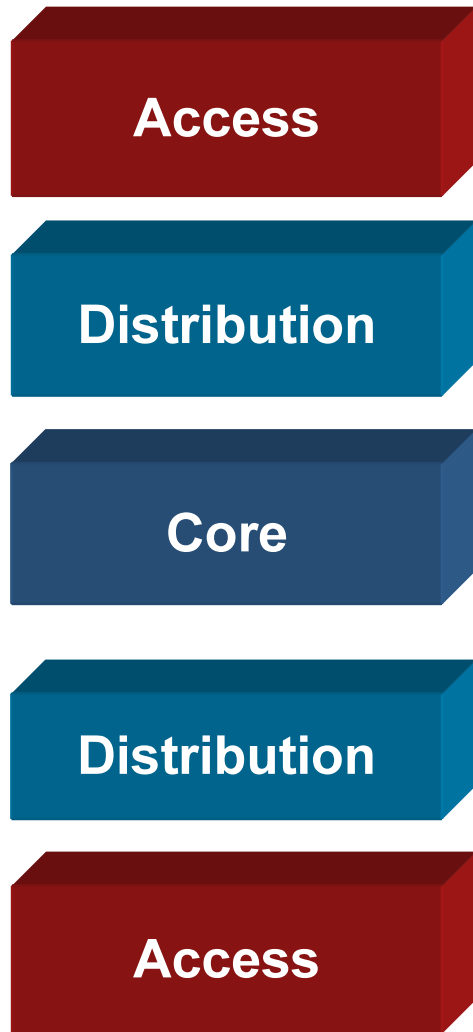
- Tunnels
 - GRE/mGRE
 - L2TPv3
 - Label Switched Paths—LSP (MPLS)
- Future Tag Transports
 - More about Tag Transports Later !

End to End Virtualised Groups Overlay Tunnels



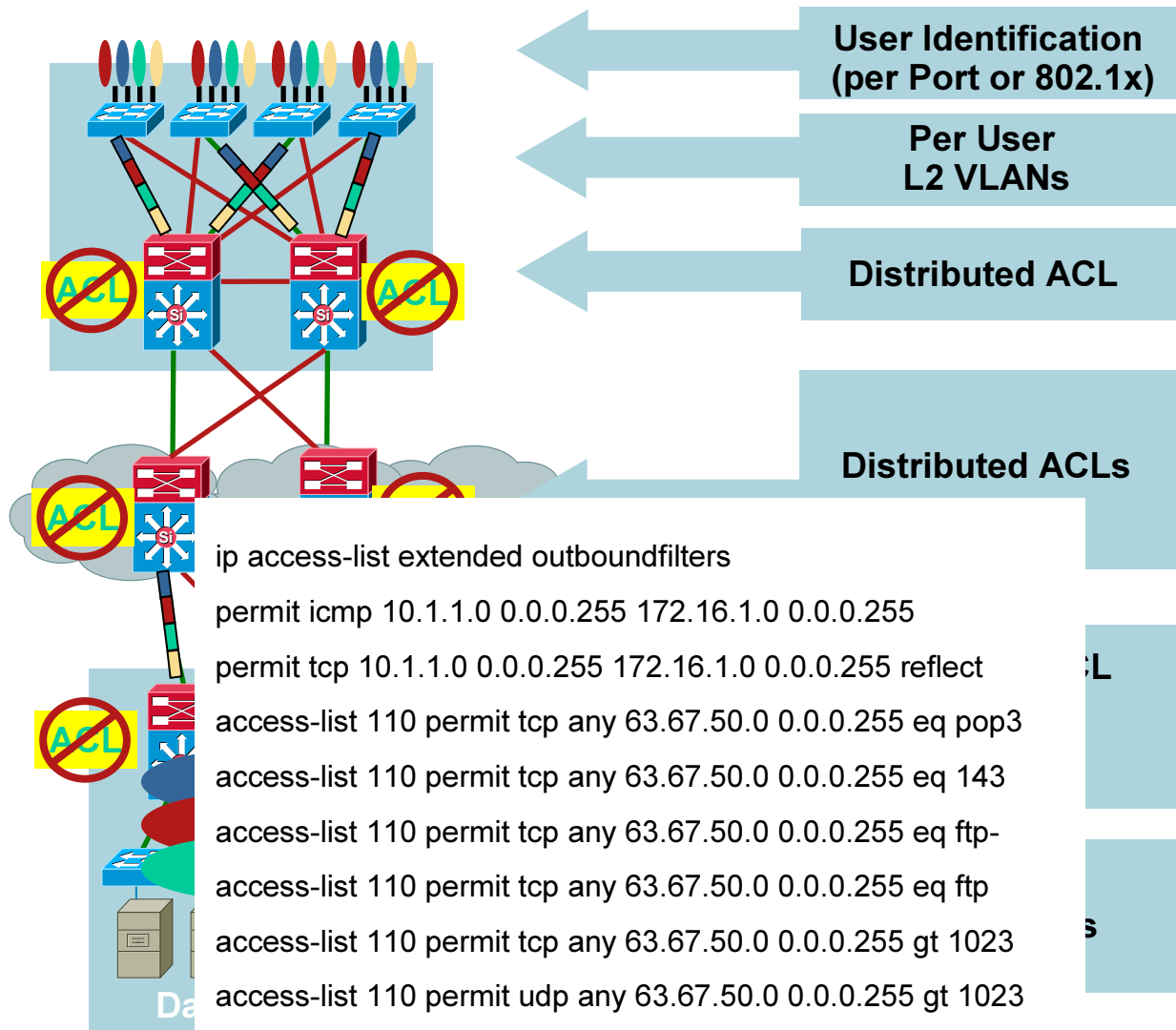
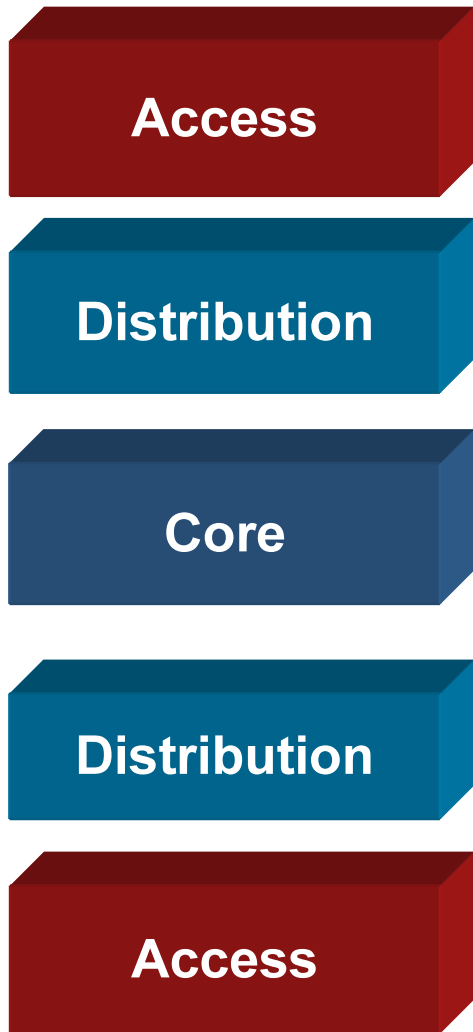
End to End Virtualised Groups

Hop By Hop



End to End Virtualised Groups

Distributed Access Control Lists

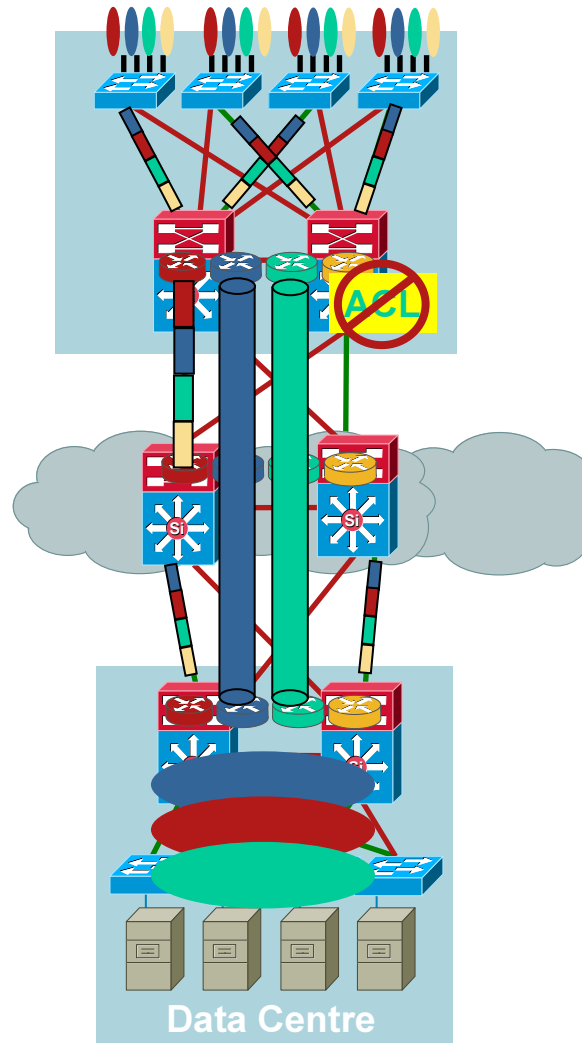
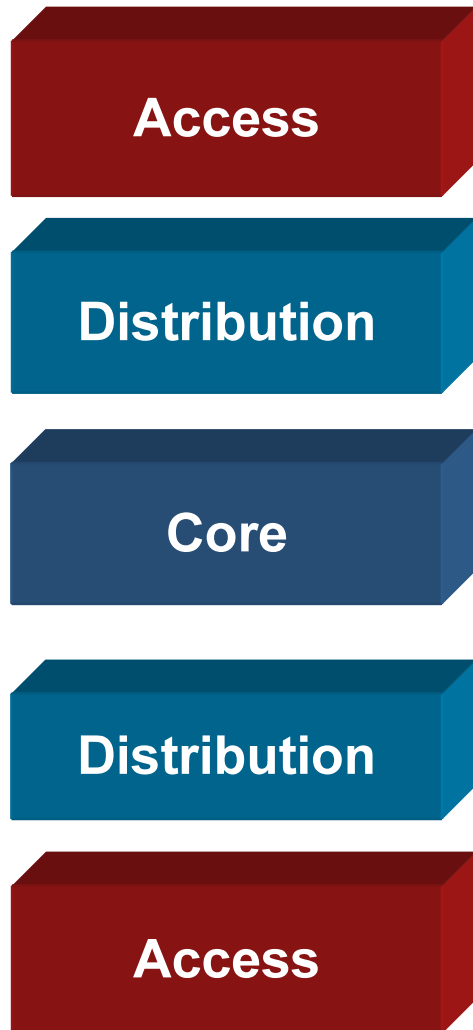


```

ip access-list extended outboundfilters
permit icmp 10.1.1.0 0.0.0.255 172.16.1.0 0.0.0.255
permit tcp 10.1.1.0 0.0.0.255 172.16.1.0 0.0.0.255 reflect
access-list 110 permit tcp any 63.67.50.0 0.0.0.255 eq pop3
access-list 110 permit tcp any 63.67.50.0 0.0.0.255 eq 143
access-list 110 permit tcp any 63.67.50.0 0.0.0.255 eq ftp-
access-list 110 permit tcp any 63.67.50.0 0.0.0.255 eq ftp
access-list 110 permit tcp any 63.67.50.0 0.0.0.255 gt 1023
access-list 110 permit udp any 63.67.50.0 0.0.0.255 gt 1023
access-list 110 deny ip 63.67.50.0 0.0.0.255 any
    
```

End to End Virtualised Groups

Static Topology Bound



All VLANs must appear on all access switches for mobility
Prone to configuration error
Not scalable manageable

VLANs may have to be transported manually across L3 boundaries per VRF

Overlay tunnels obfuscate traffic and bypass policy

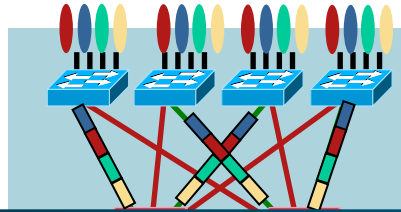
Any changes to policy requires changes to entire network

May suit specific requirements

End to End Virtualised Groups

Static Topology Bound

Access

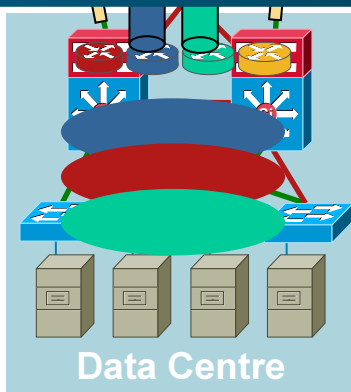


All VLANs must appear on all access switches for mobility
Prone to configuration error
Not scalable manageable

VLANs may have to be

Static Topology-Bound Separation & Policy Enforcement

Distribution

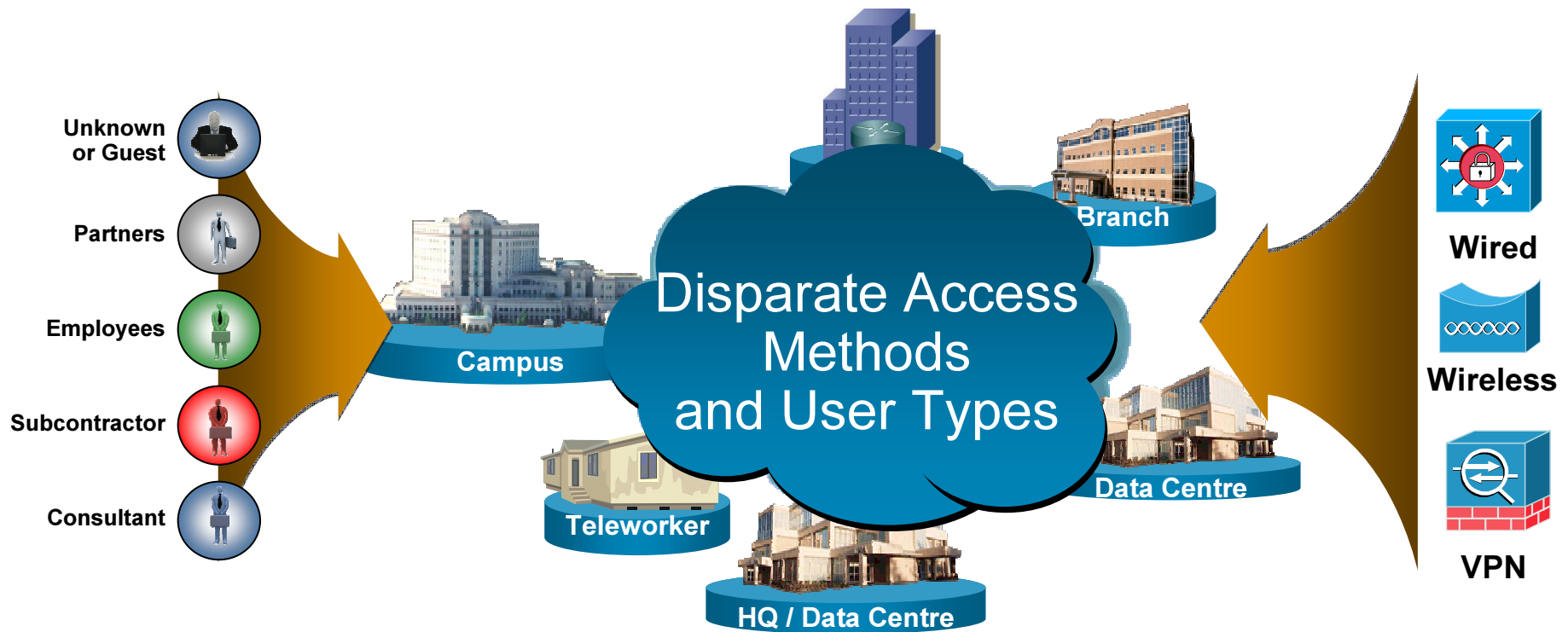


Any changes to policy requires changes to entire network

Access

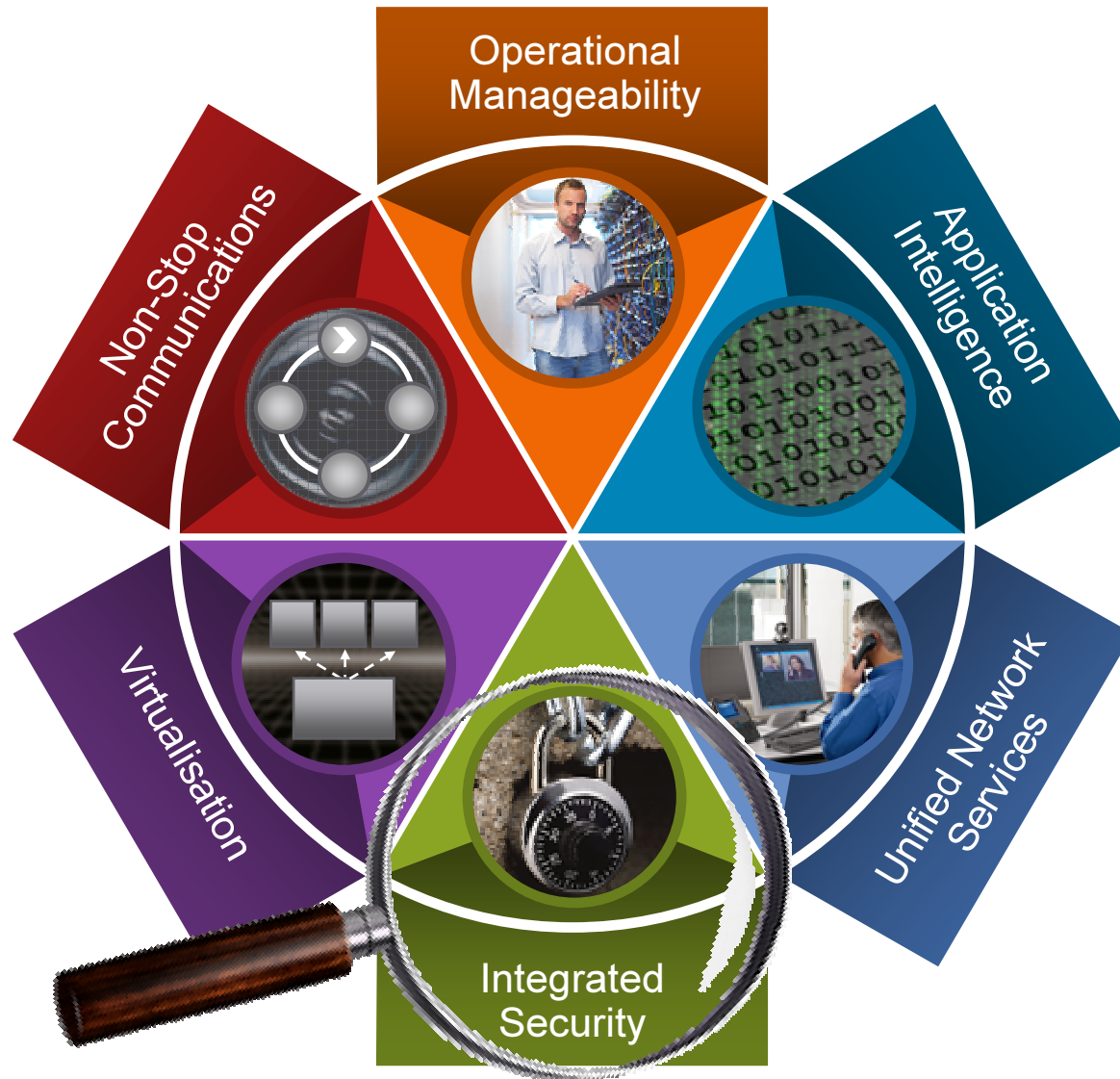
May suit specific requirements

Architectural Approach to Policy Enforcement



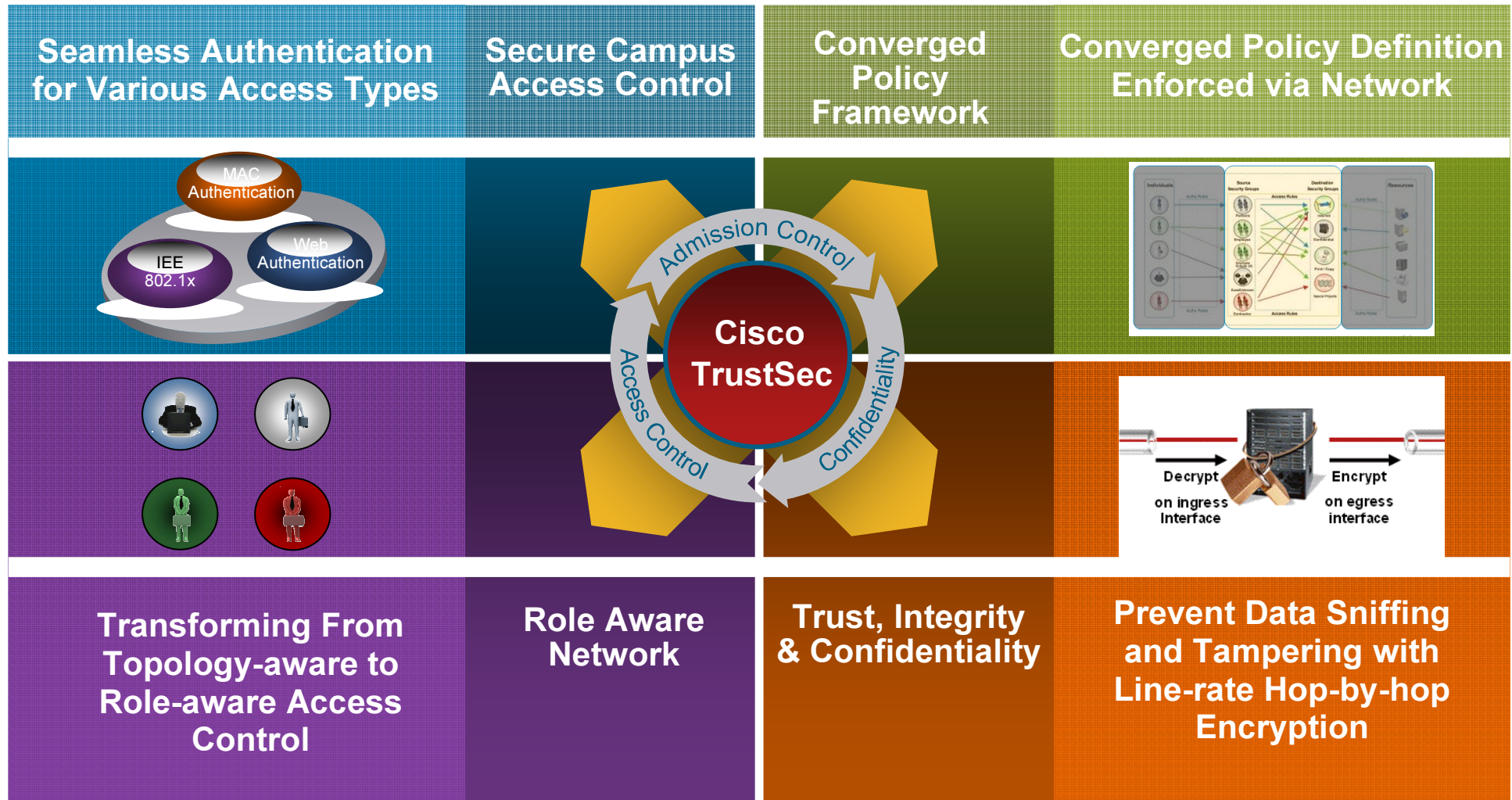
Architectural Approach

Campus Communications Fabric



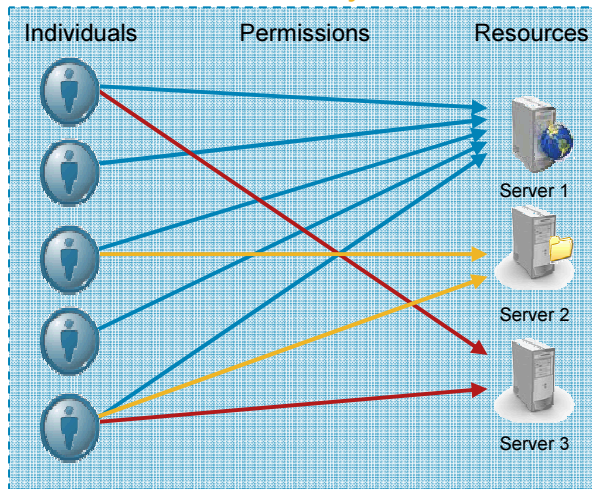
Cisco TrustSec (CTS) Vision

An Architectural Approach



Policy Enforcement Challenges

Traditional Discretionary Access Control



Access List for S1

```
access-list 101 permit tcp S1/32 D1/32 eq http
access-list 101 permit tcp S1/32 D1/32 eq https
access-list 101 permit tcp S1/32 D2/32 eq ftp
access-list 101 permit tcp S1/32 D2/32 eq http
access-list 101 permit tcp S1/32 D2/32 eq https
access-list 101 permit tcp S1/32 D2/32 eq ftp
access-list 101 permit udp S1/32 D1/32 gt 1023
access-list 101 permit udp S1/32 D2/32 gt 1023
```

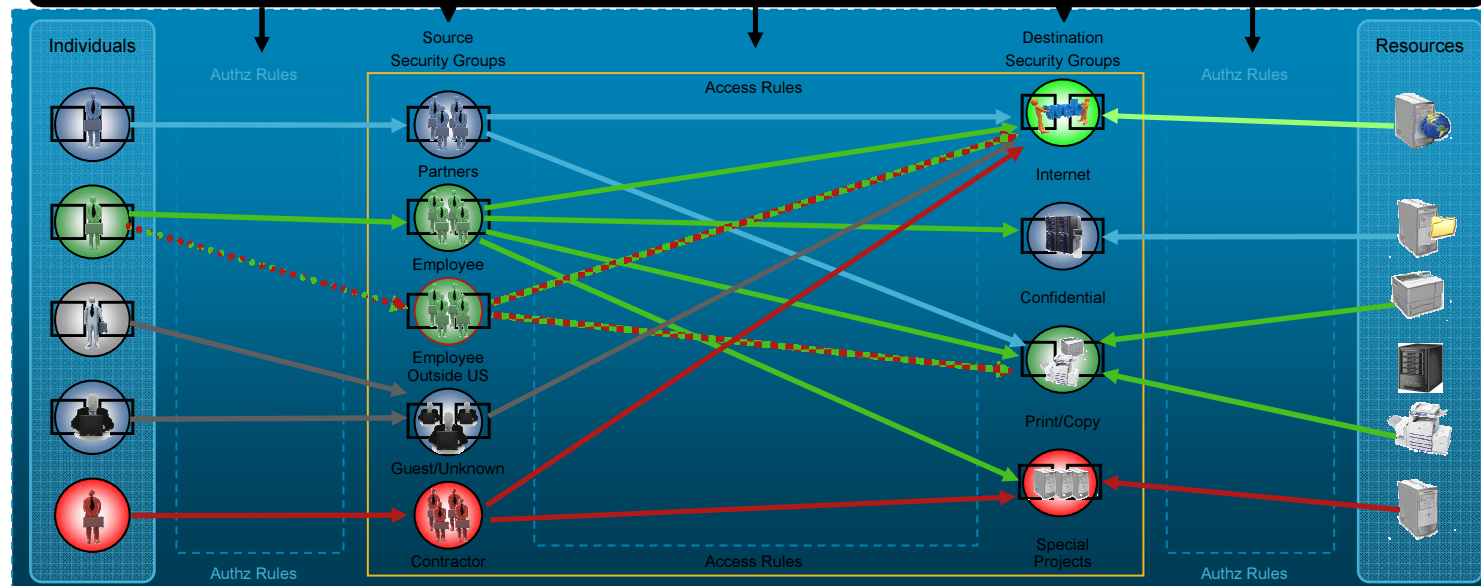
Challenges

- **Leads to ACE explosion**
(# of sources) X (# of Destinations)
X (# of permissions) = # ACEs
- **IP-address based ACLs are challenging**
Changes in addressing schemes
Use of DHCP
Proliferation of Wireless LAN devices
- **Assumes relatively static placement of users/resources**

Cisco TrustSec User Authorization and Access Control

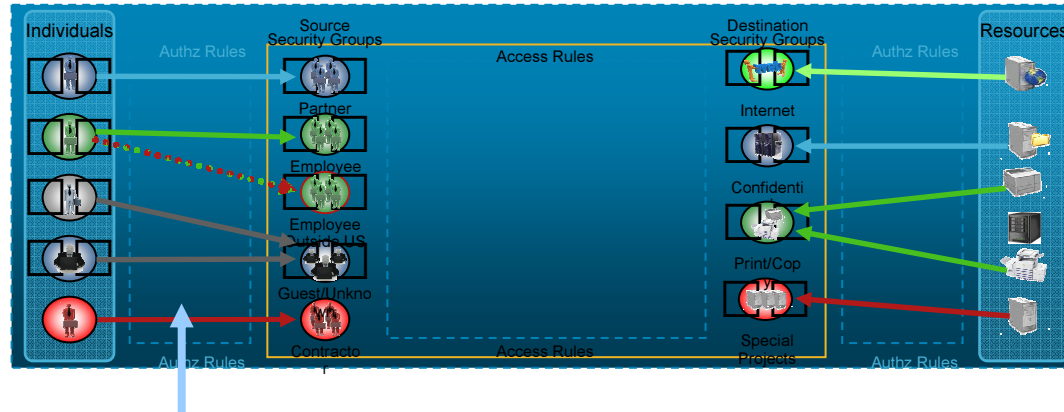
- Define Security Groups
- Users and Resources Sessions are Authorized via flexible ABAC model
- Access Control Policies are created without regards to Network Topology (No IP Addresses or subnets)
- Access Control Policies are mapped between source and destination Security Groups via a Matrix
- At runtime user's traffic carries the Security Group Tag (SGT) in every packet
- These SGTs are filtered (i.e., SGACLs) processed at wirespeed on egress devices

SG Tags & ACLs



Cisco TrustSec User Authorization Policy

Attribute Based Access Control



Users and Resources are grouped into logical *topology independent* security groups

- Ingress Session Authorization Example

Authorization Rule : if ((identity group = Teller) & (location = Sydney) & (access-type = LAN) apply SGT-4

Authorization Rule : if ((identity group = Teller) & (location = Sydney) & (access-type = WLAN) apply SGT-9

Authorization Rule : if ((identity group = Teller) & (location = Sydney) & (access-type = VPN) apply SGT-22

Authorization Rule : if ((identity group = Teller) & (location = Russia) & (access-type = LAN) apply SGT-50

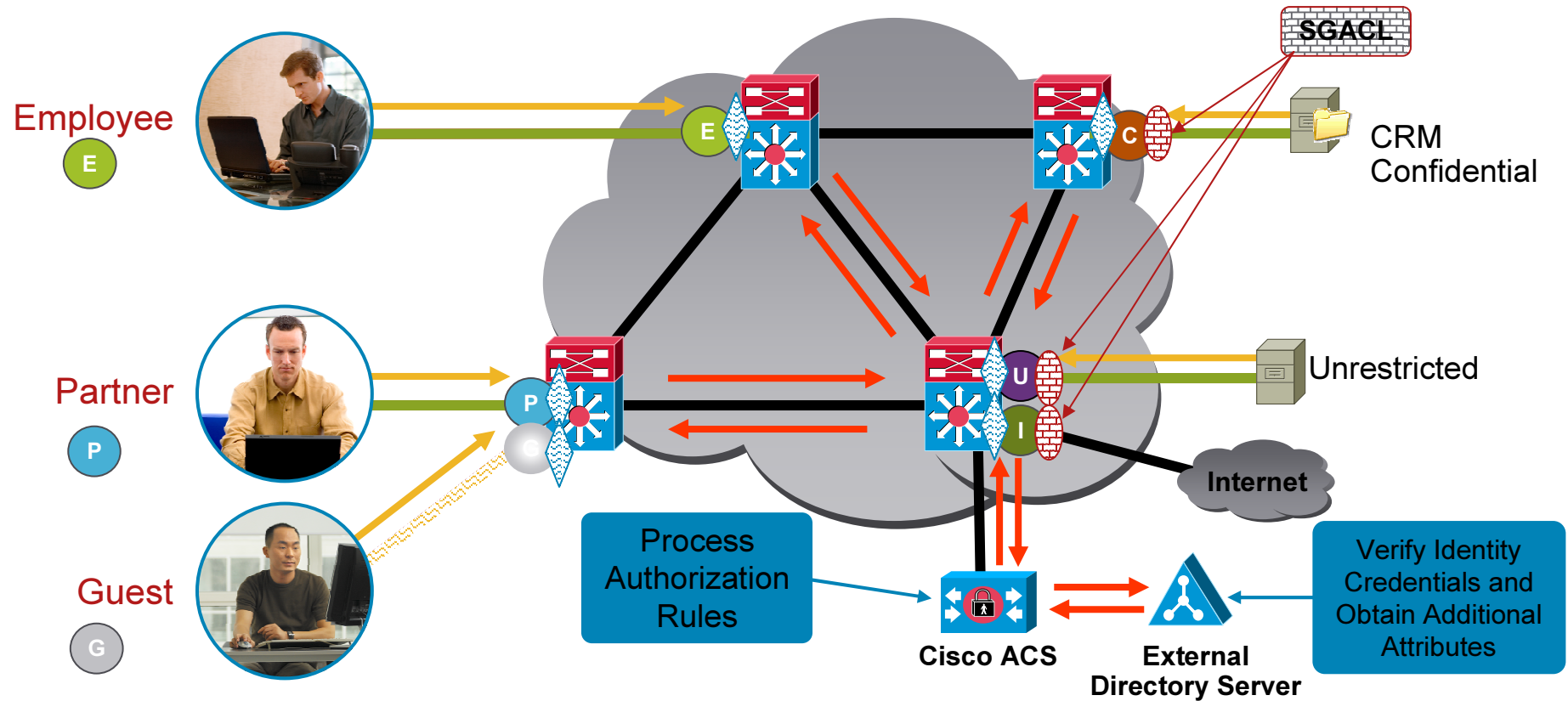
Authorization Rule : if ((identity group = Teller & Manager) & (location = Russia) & (access-type = LAN) apply SGT-51

Authorization Rule: if ((identity group = LoanApp) & (location == DC-Sydney) then apply SGT-1

Access Control is then defined by ACL between Source & Destination Groups

RBACL1 : s(SGT-4) d(SGT-1) permission-list A

Policy Enforcement Throughout the Network: Role Based Access Control Set-up

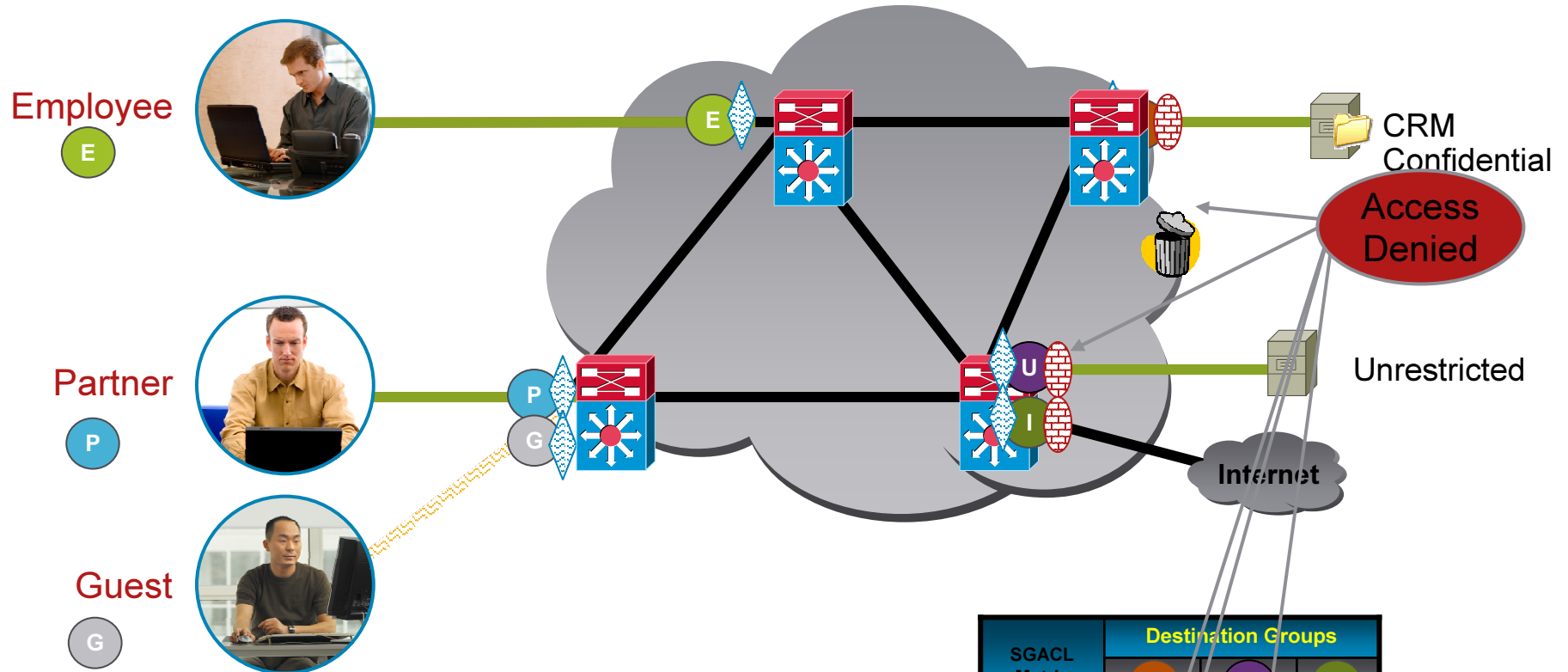


Legend

Link/Port Status		Security Group Classifications			
	Unauthenticated		Employee Group		Confidential Group
	Failed Authentication		Partner Group		Unrestricted Group
	Authenticated		Guest Group		Internet Group
	Shutdown		Ingress Tagging		Egress Filtering

1. Authentication Request
2. Radius and AD Authc/Authz
3. Group Membership Dynamically Assigned
4. SGACL Dynamically Applied
5. Links Up

Policy Enforcement Throughout the Network: Role Based Access Control Deployment



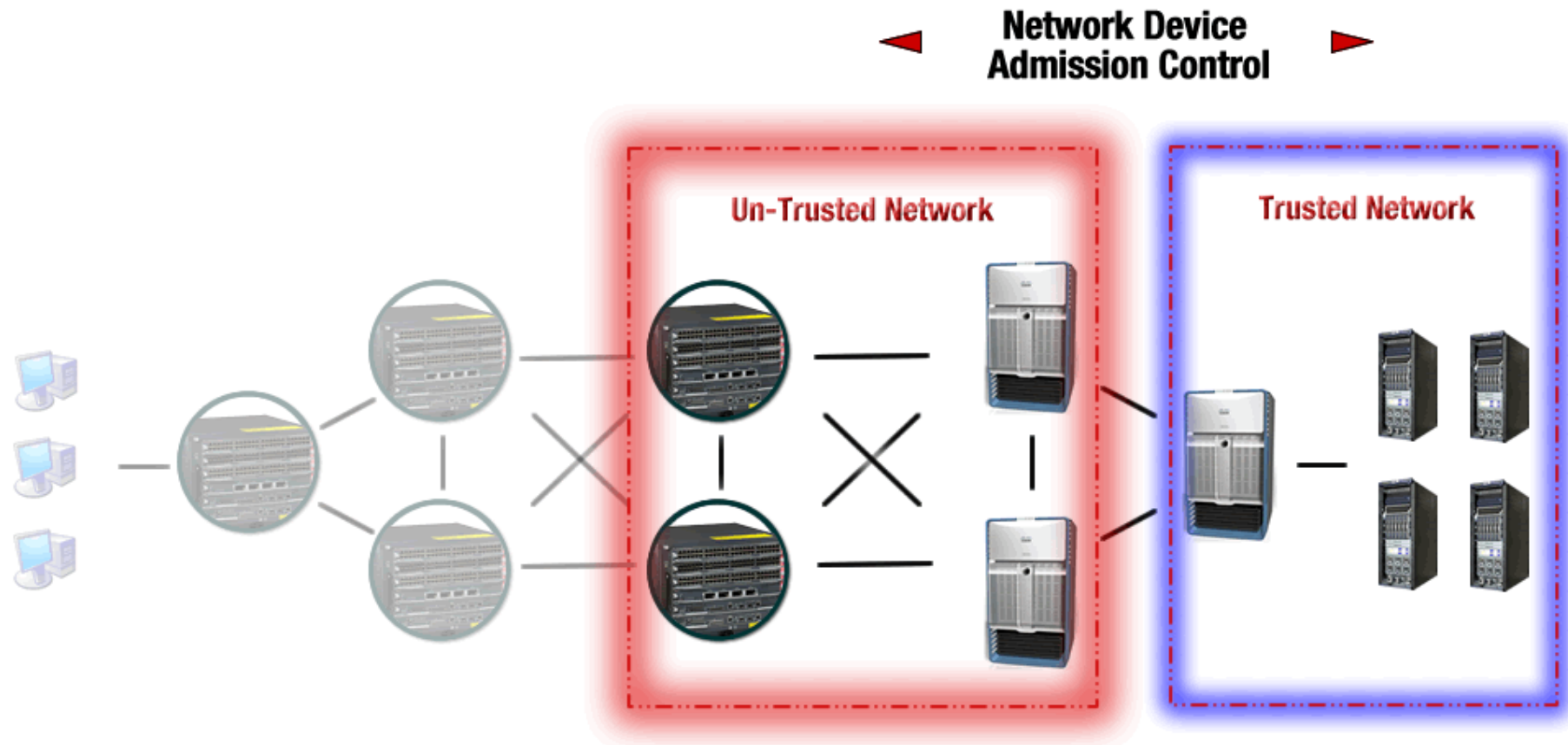
Legend

Link/Port Status		Security Group Classifications					
	Unauthenticated		Ingress Tagging		Employee Group		Confidential Group
	Failed Authentication		Egress Filtering		Partner Group		Unrestricted Group
	Authenticated				Guest Group		Internet Group
	Shutdown						

SGACL Matrix	Destination Groups		
	C	U	I
Source Groups			
E			
P			
G			

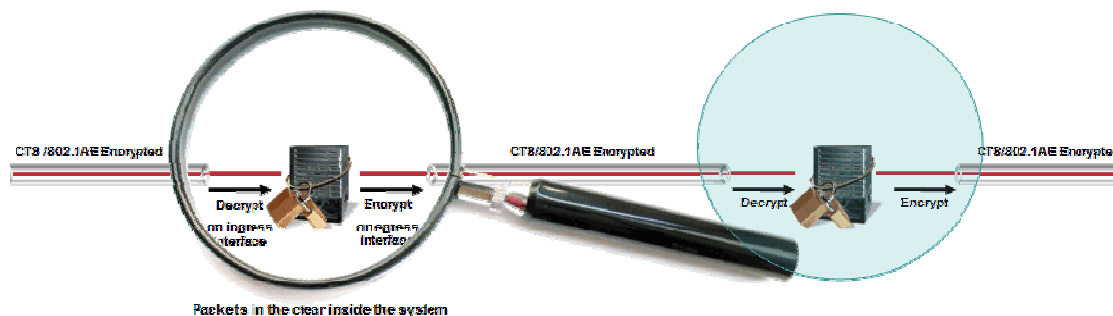
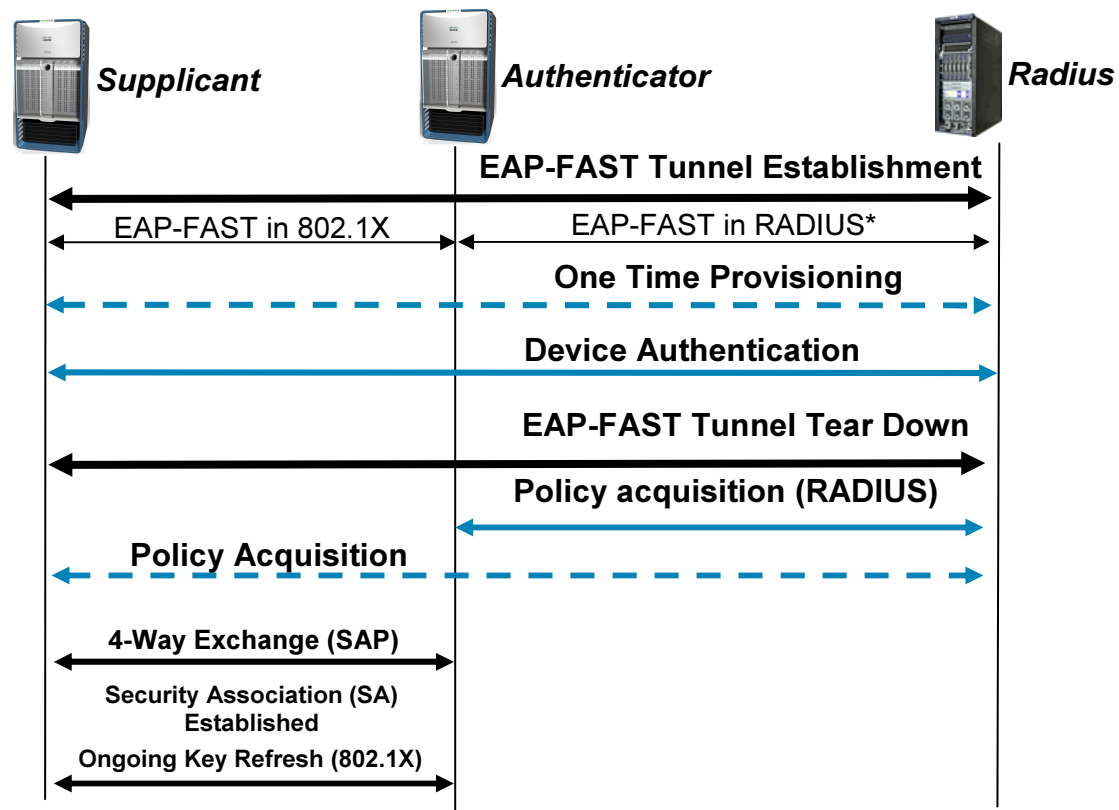
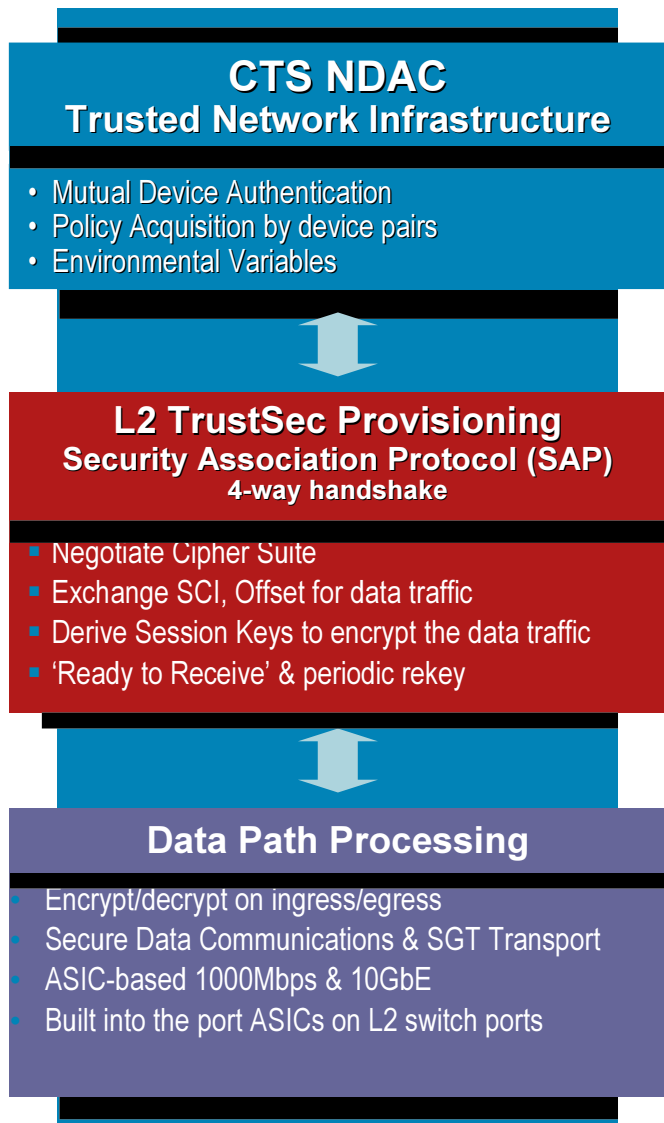
Network Device Admission Control (NDAC)

Authenticate network devices to trusted node via 802.1X prior to any network connectivity becoming available. Polices are then acquired as well as any applicable session keys or cipher suites...



TrustSec

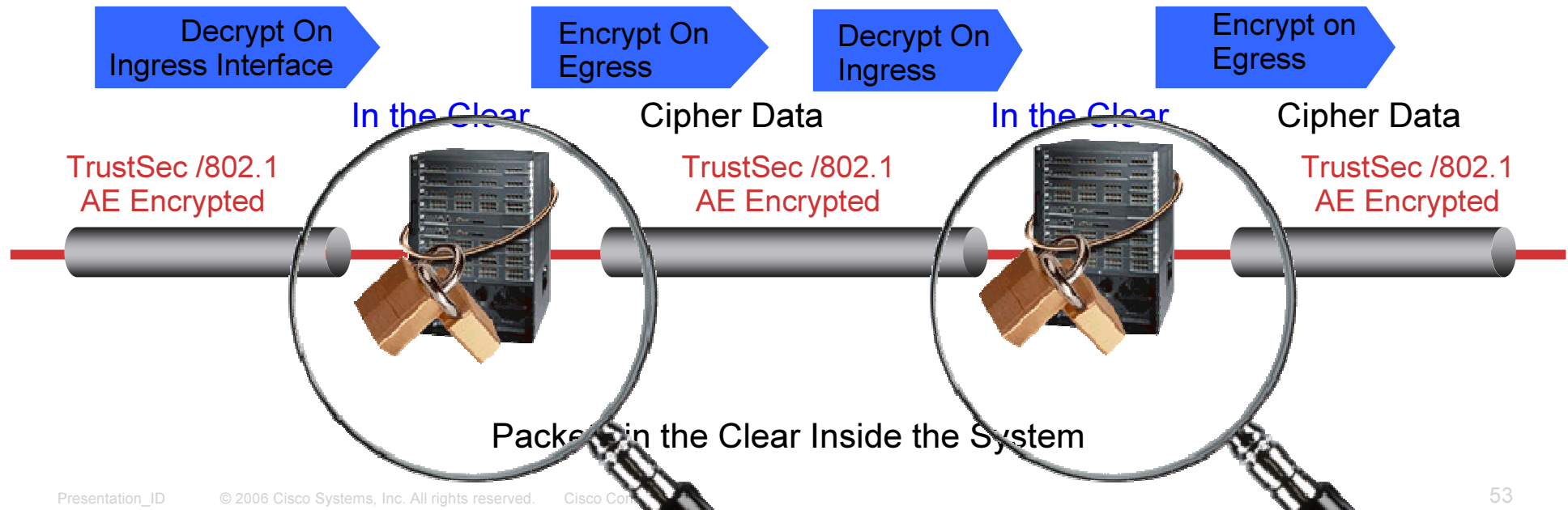
An Architectural Approach



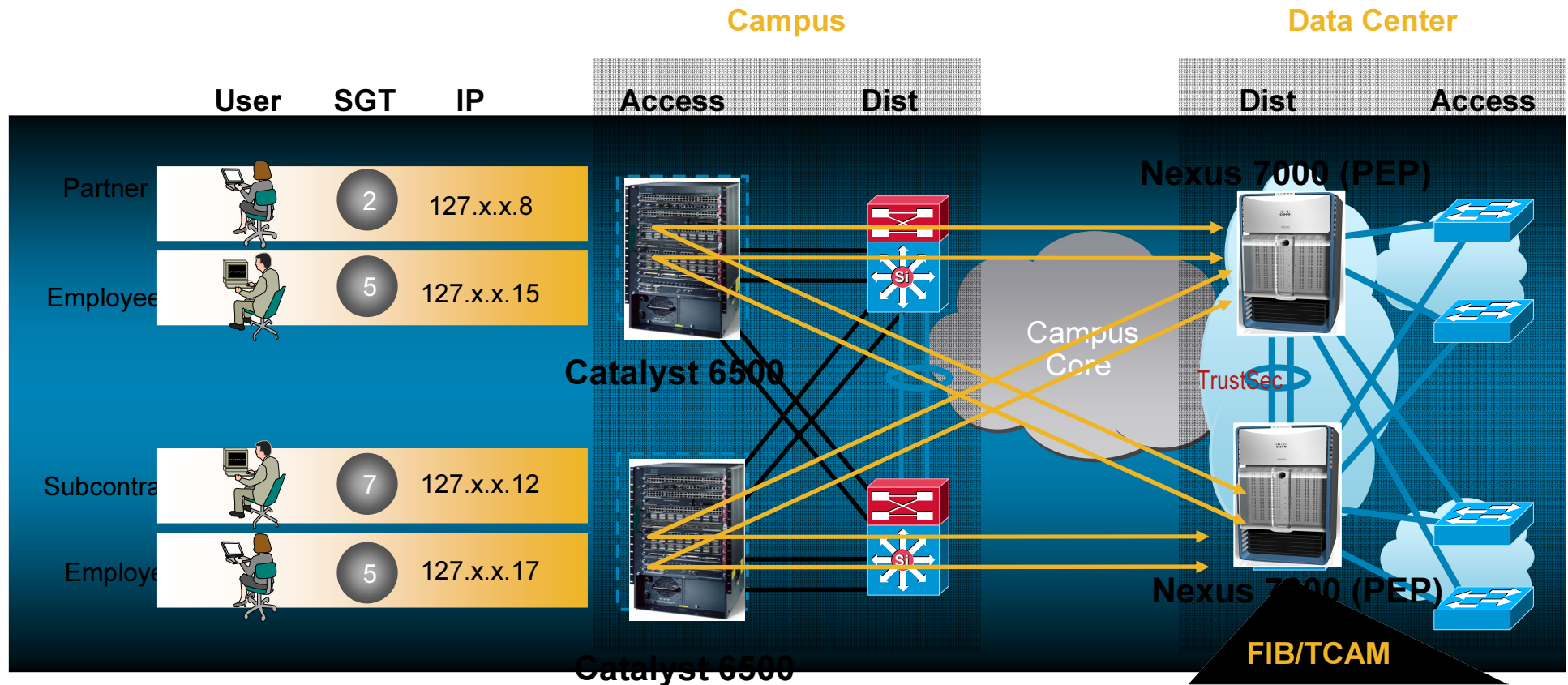
Link Layer Encryption

An Architectural Approach

- Hop-by-Hop packet confidentiality and integrity via IEEE 802.1AE
- Packets are encrypted on egress, Packets are decrypted on ingress
- Packets are in the clear in the device allowing the network to continue to perform all the packet inspection features currently used
- Can be incrementally deployed depending on link vulnerability



Security Group Tag Exchange Protocol (SXP) Migration for Access-Layer

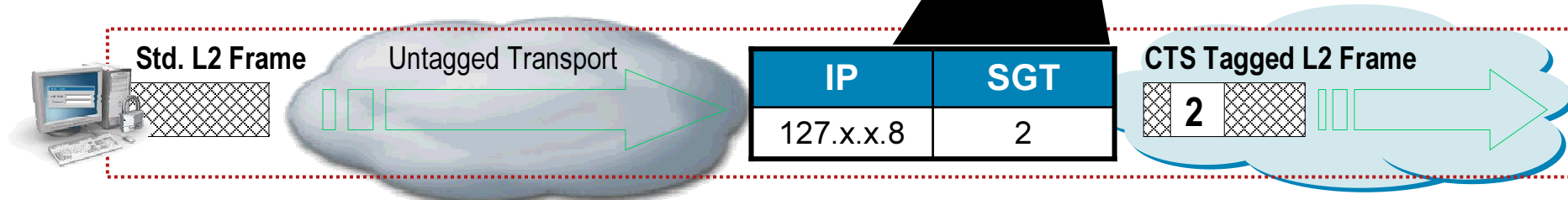
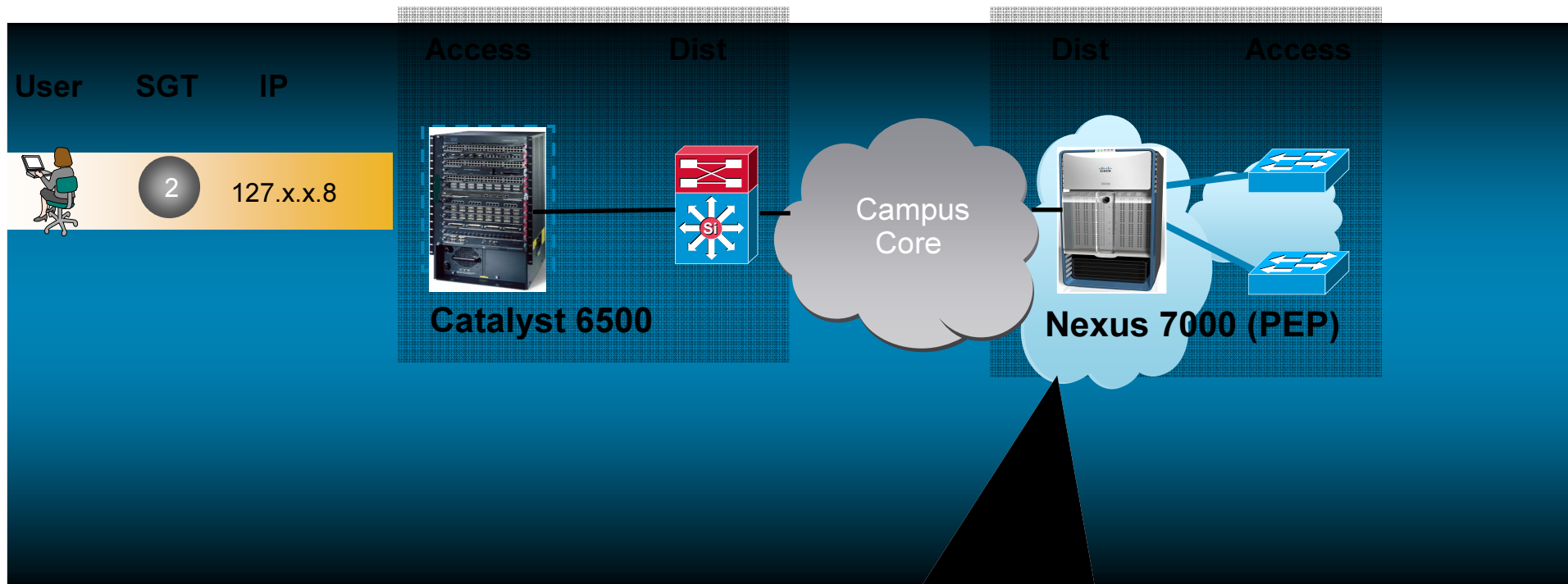


Legend

- SXP Flow
- TrustSec SW Only (SGT/IP Relay only)
- TrustSec HW Capable (SGT Tagging, Fwd and Filtering)

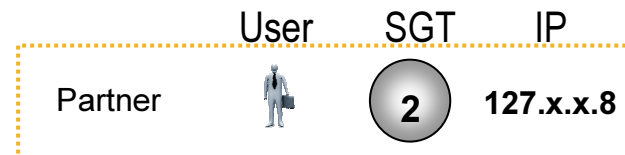
IP	SGT
127.x.x.8	2
127.x.x.15	5
127.x.x.12	7
127.x.x.17	5

Security Group Tag Exchange Protocol (SXP) Migration for Access-Layer



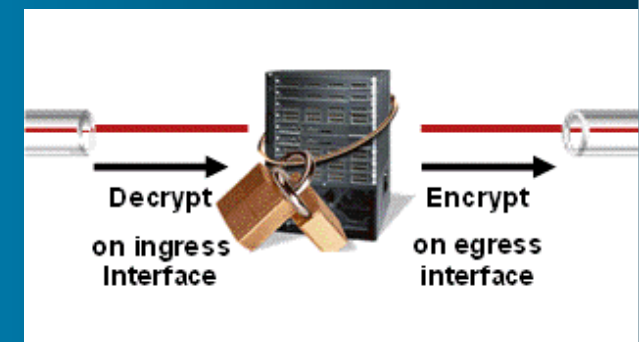
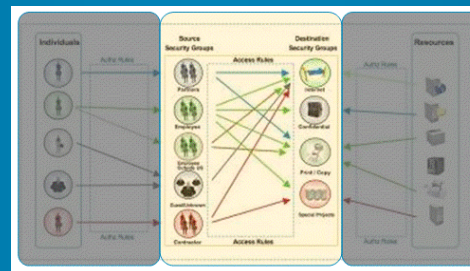
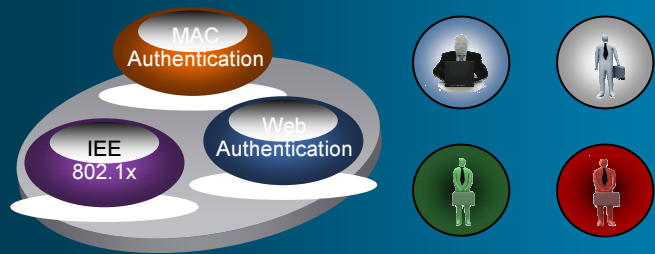
Legend

- CTS SW Only (SGT/IP Relay only)
- CTS HW Capable (SGT Tagging, Fwd & Filtering)



Cisco TrustSec Summary

- Identity = Who are you?
- NAC = Are you healthy?
- CTS = What can you do on this trusted and secure network ?
- Next generation architectural framework for identity, policy and security in DC and Campus
- Innovation & Technology Leadership



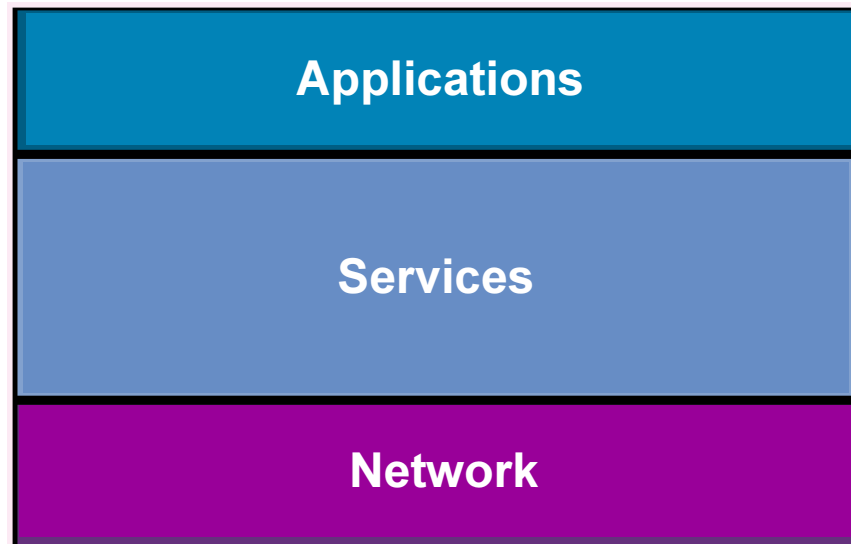
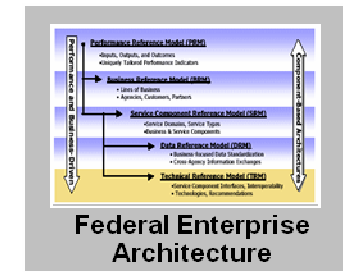
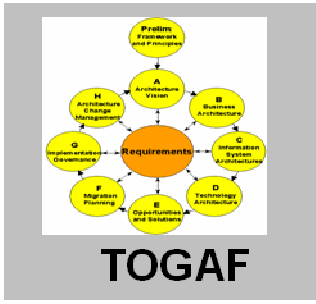
Summary

“Nothing is as dangerous in architecture as dealing with separated problems. If we split life into separated problems we split the possibilities to make good building art”

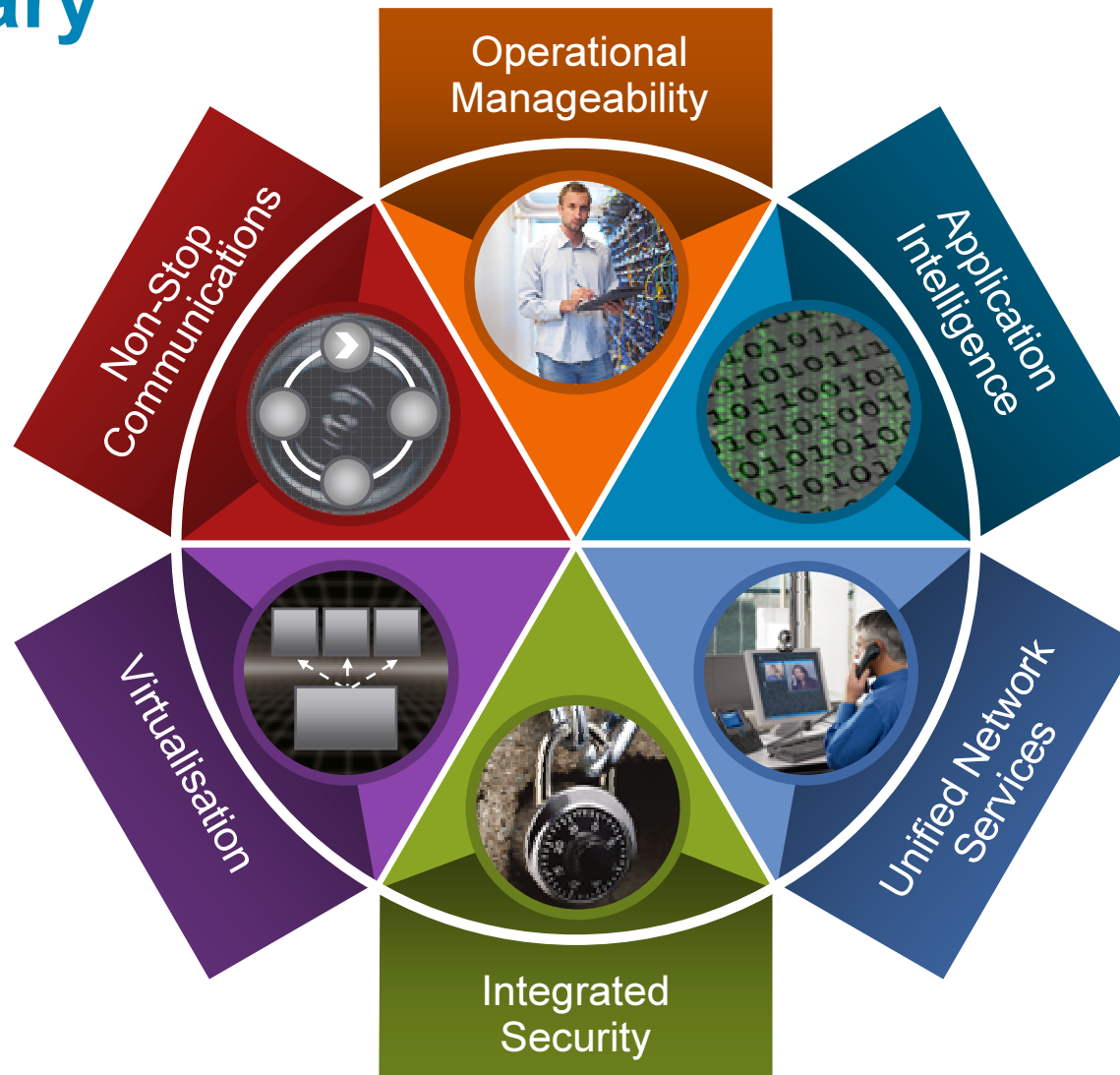
- Alvar Aalto 1898-1976

Finnish Modernist Architect/Designer

Summary



Summary



Campus Communication Fabric: <http://www.cisco.com/go/ccf>
Cisco TrustSec: <http://www.cisco.com/go/switchsecurity>

Thank You !

**Please Complete Your
Evaluation Form**

